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Practical Fault Injection Analysis of Lattice-based NIST PQC Standards - Kyber and Dilithium

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*NIST PQC Seminars,
5th May 2023*



Notice

- **Main Motive:**
 - Flavor of reported fault attacks on Kyber and Dilithium (and countermeasures)
 - Algorithmic properties and Implementation Choices facilitate efficient FIA
 - Towards fault resistant implementations of Kyber and Dilithium
- Most fault attacks demonstrated on bare-metal PQC software running on ARM Cortex-M4 processor: Clock/Voltage Glitching, EMFI
- Talk includes published works from journals, conferences, and IACR ePrint Archive.
- Talk includes works of other researchers (cited appropriately)
- For easier explanation, we ‘simplify’ concepts, and contains lots of illustrations!!

Outline

- ❑ **FIA on Kyber:**
 - ❑ **FIA on Key Generation and Encapsulation**
 - ❑ **FIA on Decapsulation**

- ❑ **FIA on Dilithium**
 - ❑ **FIA on Signing**
 - ❑ **FIA on Verification**

- ❑ **Conclusion**

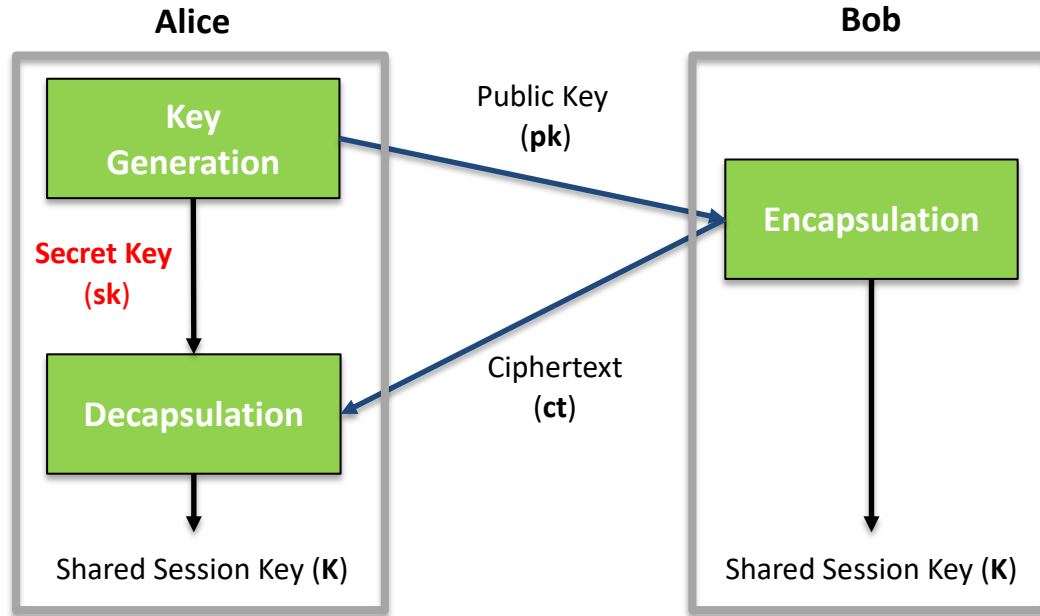
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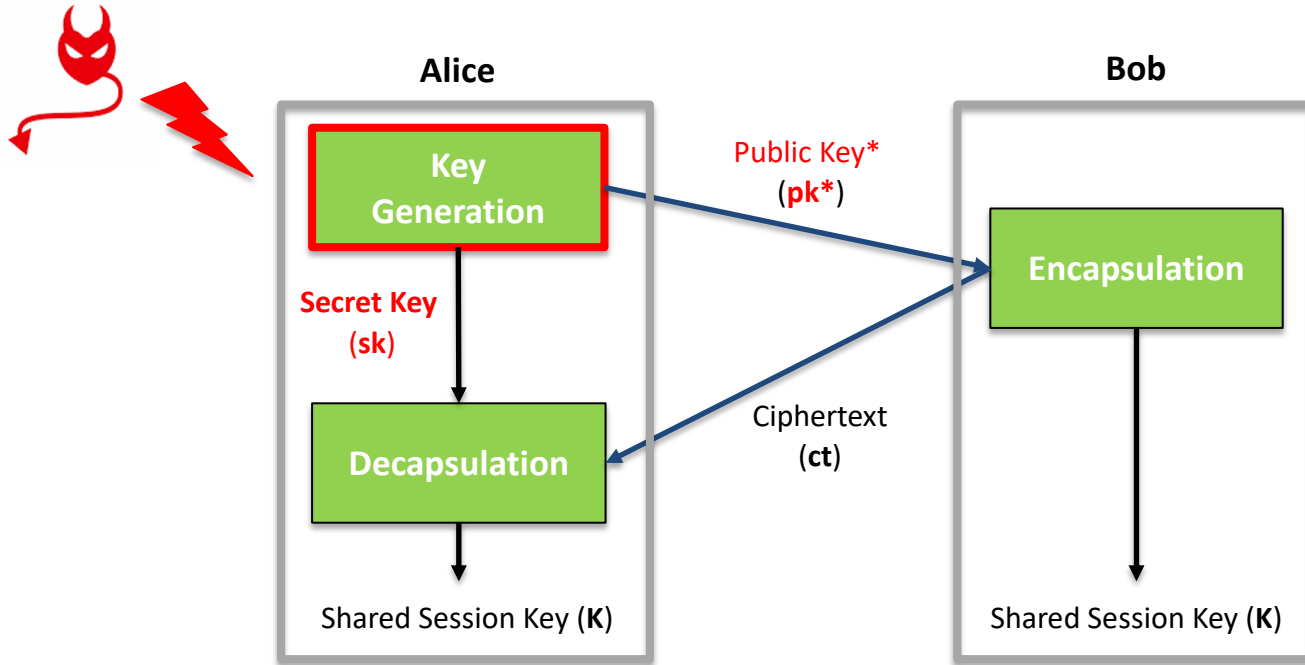
Key Encapsulation Mechanisms (KEMs)



Two Modes Possible:

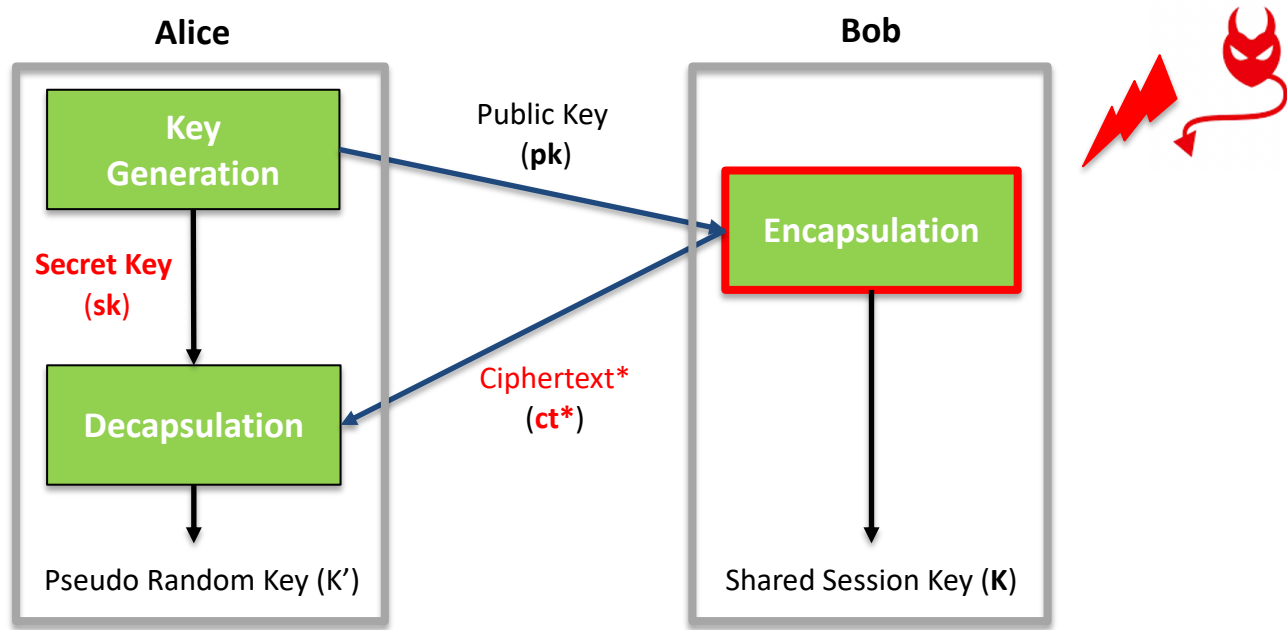
1. Ephemeral Key
2. Static Key

KEM in Ephemeral Mode: Attacking Alice



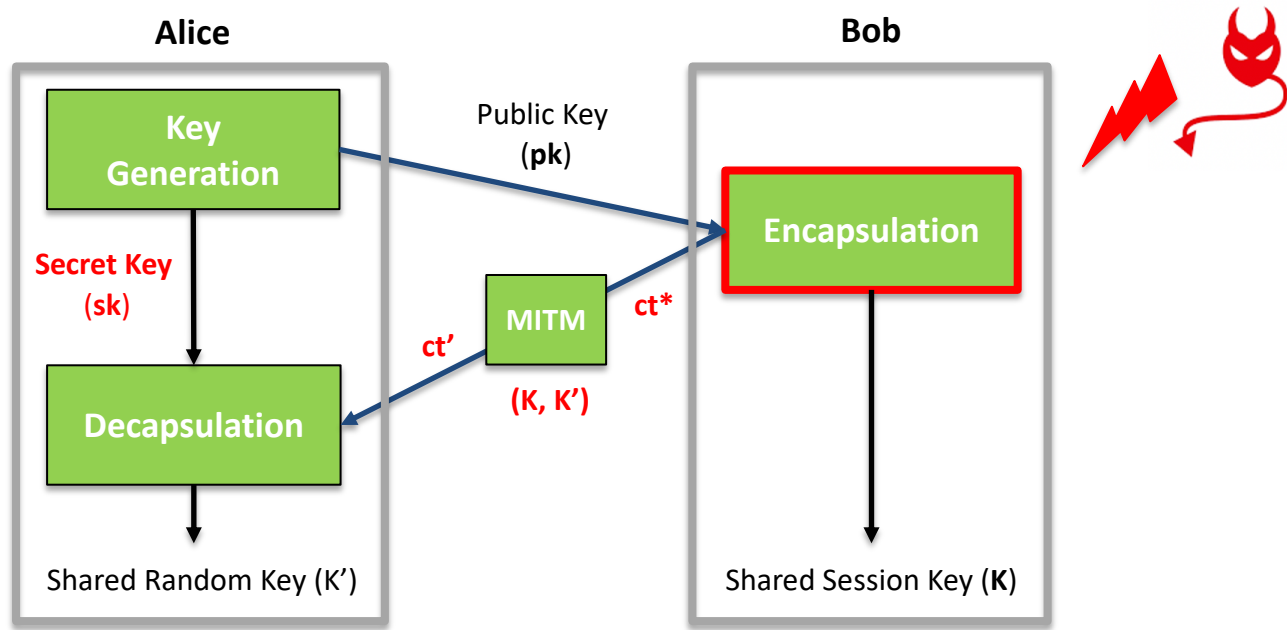
- ❑ Single execution to target Key Generation: Key Recovery Attack
 - ❑ Recover Secret key from Faulty but valid Public Key
- ❑ Decapsulation does not serve as an effective target – Attacker can only observe binary output (1-bit)

KEM in Ephemeral Mode: Attacking Bob



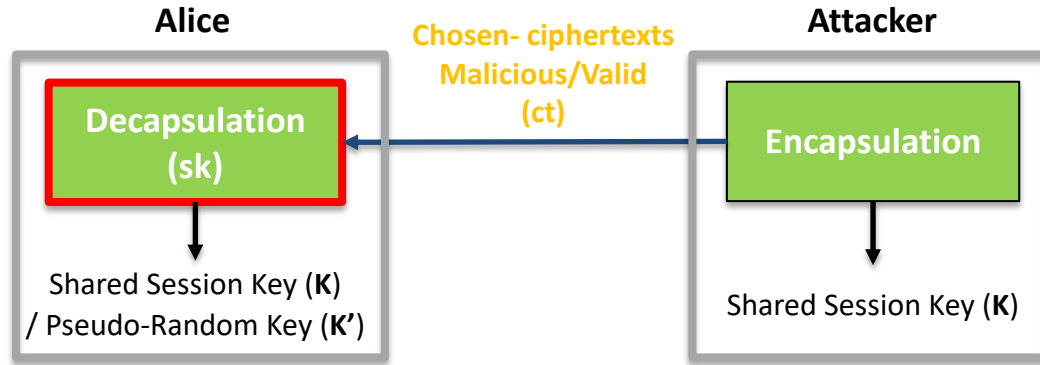
- ❑ Single execution to target Encapsulation Procedure: Message Recovery Attack
 - ❑ Recover Message from Faulty Ciphertext
 - ❑ Results in Decapsulation Failure (CCA Secure)

KEM in Ephemeral Mode: Attacking Bob

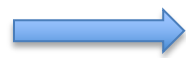


- ❑ Single execution to target Encapsulation Procedure: Message Recovery Attack
 - ❑ Recover Message from Faulty Ciphertext
 - ❑ Results in Decapsulation Failure (CCA Secure)
 - ❑ Attacker can still perform a MITM Attack!!!

KEM in Static Mode: Attacking Alice



| Chosen Ciphertext | Output |
|-------------------|---------|
| CT1 | Success |
| CT2 | Failure |
| CT3 | Success |
| ... | ... |



Key
Recovery

**Fault Assisted Chosen-Ciphertext
Attacks**

**Decryption Failure (DF) Oracle
through Faults**

Fault Attack Characteristics

- We describe known fault attacks with the following characteristics:
 - **Attacker's ability to communicate with DUT (DUT_IO_Access):**
 - Observe_DUT_IO: Can only passively observe target's IO
 - Communicate_DUT_IO: Can communicate with target
 - **Targeted or Not (Profiling, Knowledge of Implementation):**
 - Targeted_Fault
 - Random_Fault
 - **Type of Fault:** Control Flow, Data
 - **Number of Faults within Single Computation:** Single/Multiple
 - **Total Number of Faulty Computations**

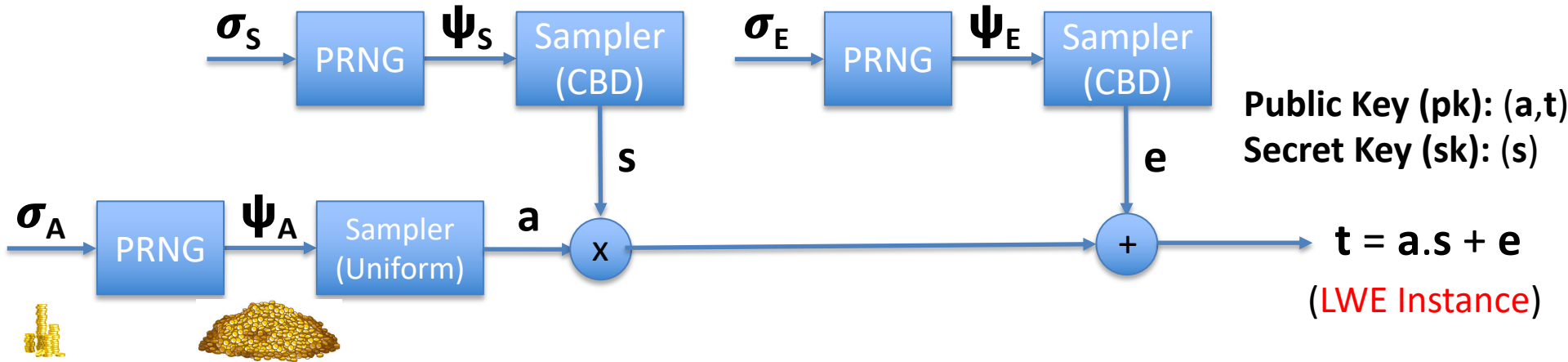
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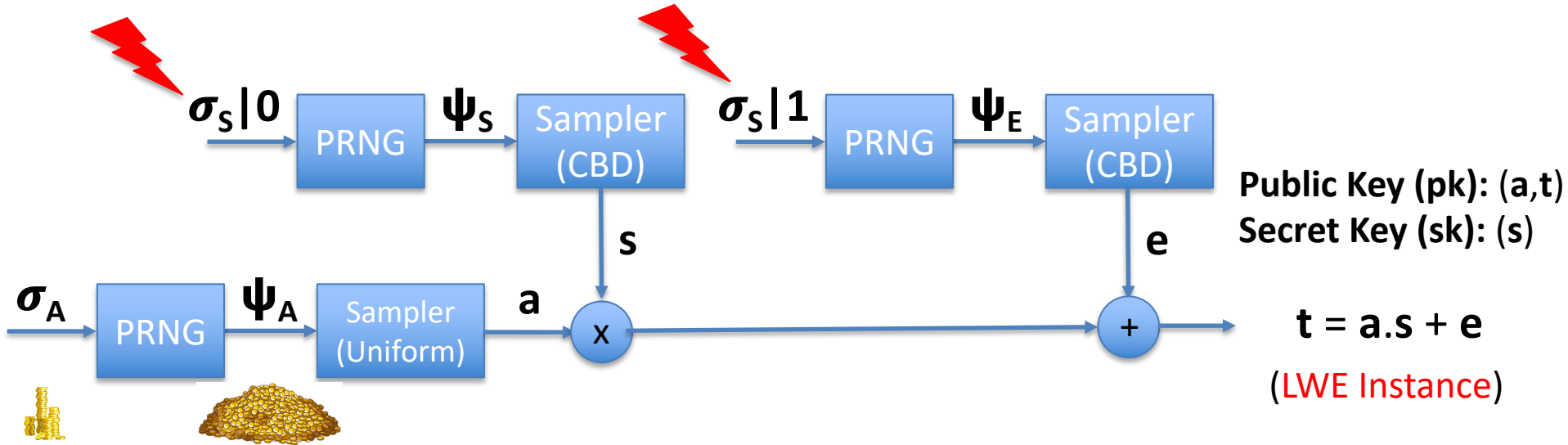
- ❑ Conclusion

FIA on Key Generation



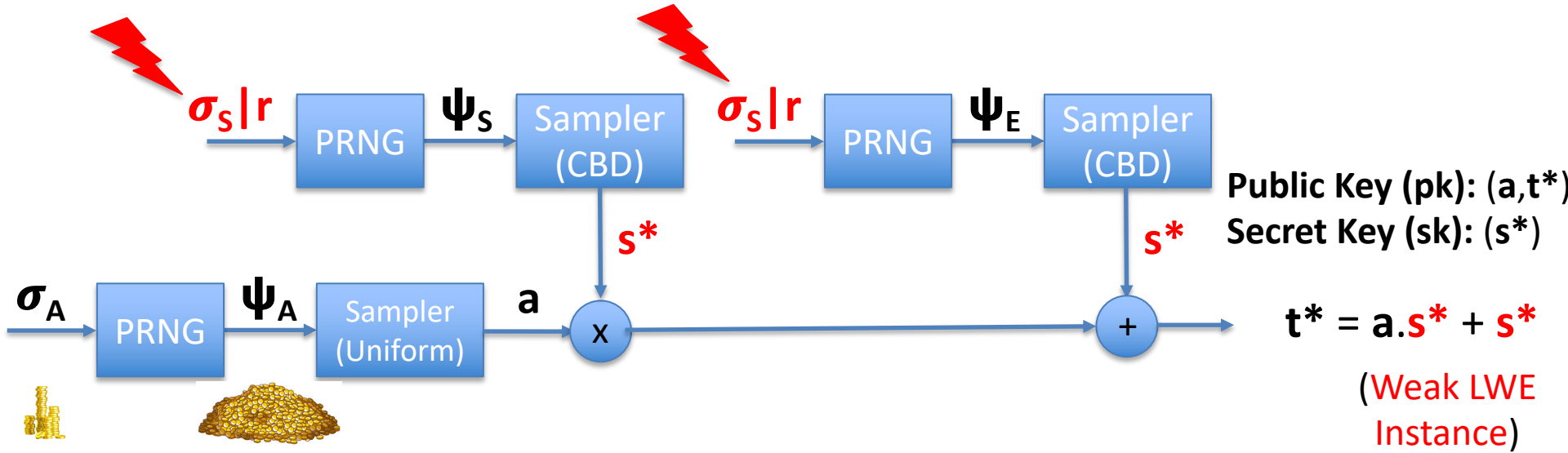
- **Aim of FIA on KeyGen:**
 - Weak LWE instances (easily solveable)
 - Secrets with Low Entropy

FIA on KeyGen: Weak LWE instances [RRB+19]



- **Fault Vulnerability:** Seed used to sample s and e only differ by a single byte
 - Inject fault to force using same seed for s and e

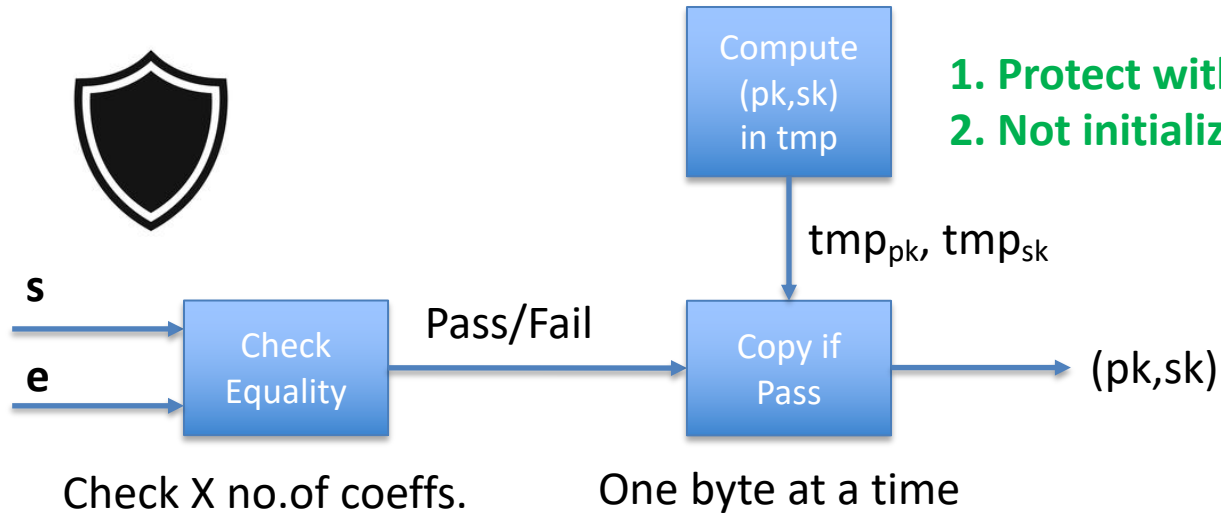
FIA on KeyGen: Weak LWE instances [RRB+19]



- **Weak LWE Instance:** Solved by Gaussian Elimination
- Applicability to Kyber KEM: Inject **k-2k** targeted faults
- Round-1 Kyber used rounded public keys, but rounding was removed from Round-2

FIA on KeyGen: Weak LWE instances [RRB+19]

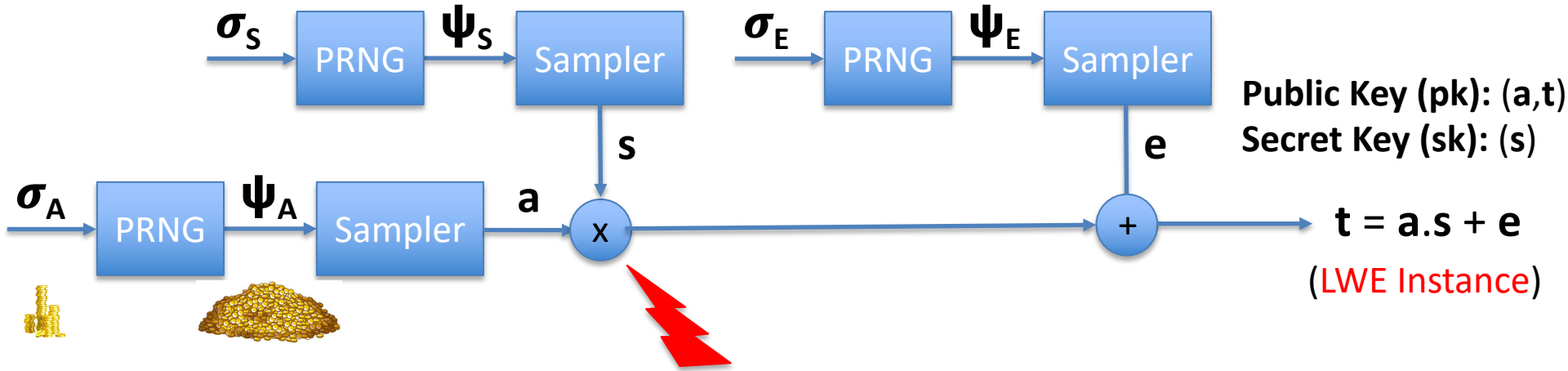
- Single fault enough for NewHope (Ring-LWE) and FrodoKEM (Standard-LWE)
- **Impact:** The algorithm of FrodoKEM (Finalist NIST PQC candidate) was modified in Round 2 to eliminate the fault vulnerability.
 - Completely different seeds were used to sample \mathbf{s} and \mathbf{e}



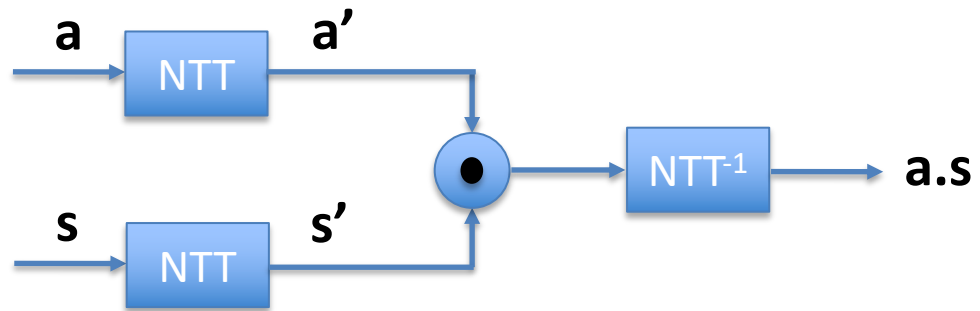
FIA on Kyber KeyGen/Encaps: Summary

| Attack Name | DUT_IO_Access | Type of Fault | Targeted_or_Not | No. of Faults within Single Computation | Total No. of Faulty Computations | Countermeasure |
|---------------------------|----------------|---------------|-----------------|---|----------------------------------|--|
| Nonce_Fault_Attack | Observe_DUT_IO | Data /Control | Targeted | k-2k | 1 | Check equality of s and e Copy Public Key if Pass |

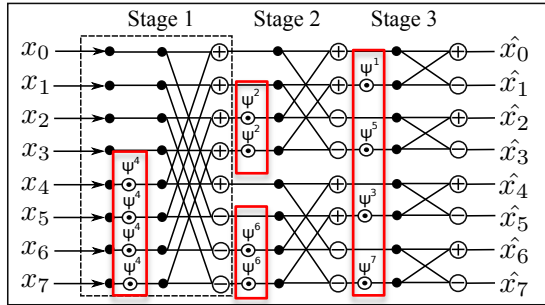
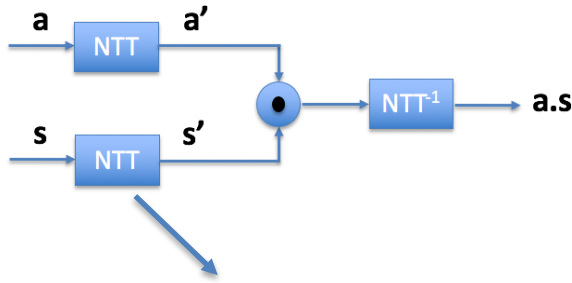
FIA on KeyGen: Reduce Entropy of Secret [RYB⁺23]



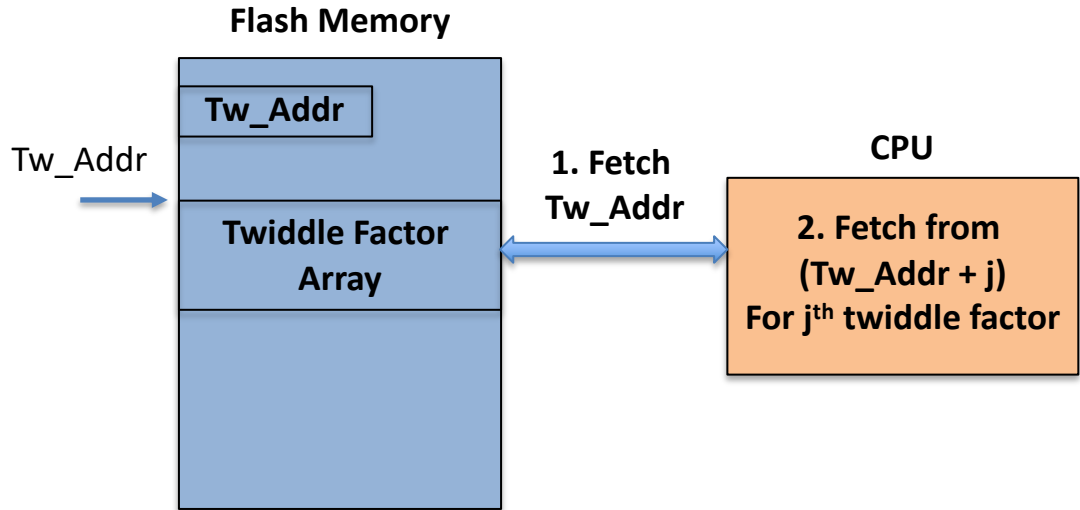
Polynomial multiplication is done using Number Theoretic Transform (NTT)



FIA on KeyGen: Reduce Entropy of Secret [RYB+23]



In MCU, Twiddle Constants are stored in Flash Memory as part of Firmware Binary

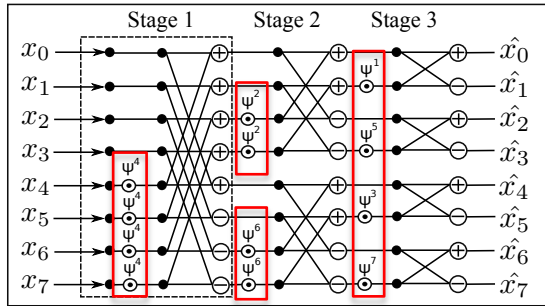
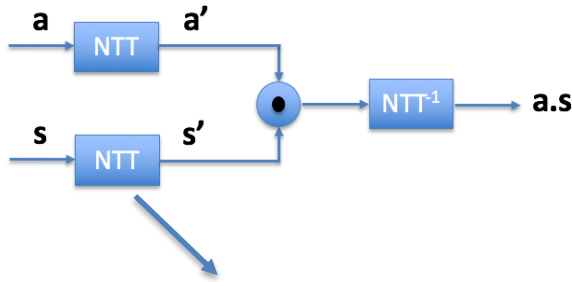


Main Observation: Tw_Addr is used as **base-address** to calculate address for all constants

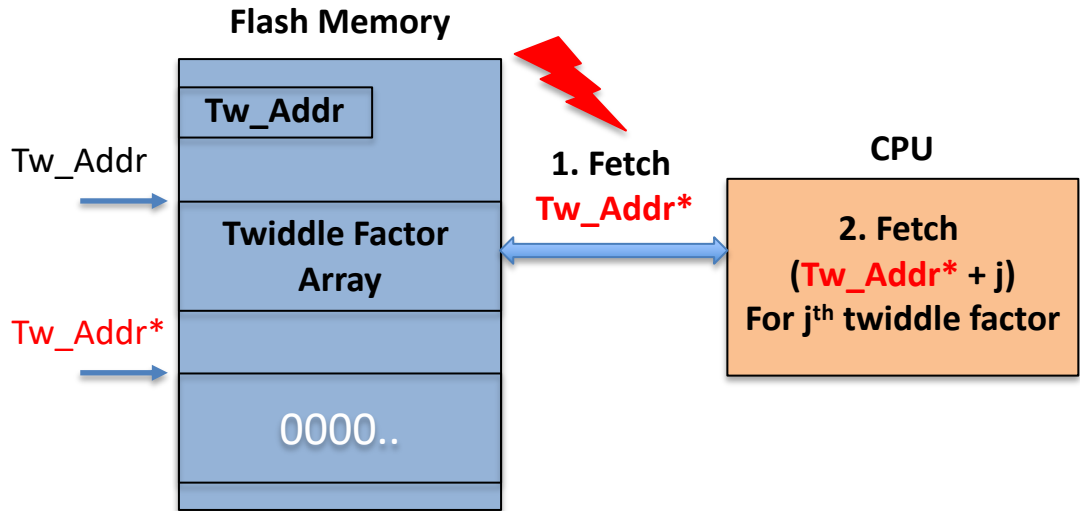
Fault Vulnerability: Can an attacker fault the base address?

Implementation Style used in all publicly available optimized implementations of Kyber and Dilithium for ARM Cortex-M4 Processor

FIA on KeyGen: Reduce Entropy of Secret [RYB+23]



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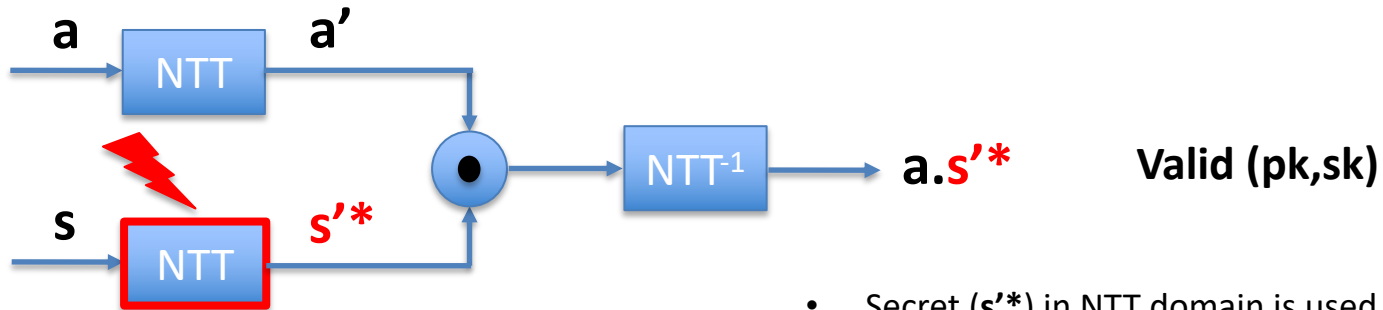


Observation: Can zeroize the entire twiddle factor array in a single fault

25% of random memory locations yield zeros on ARM Cortex-M4 processor

What happens when twiddle factors are zeroized???

FIA on KeyGen: Reduce Entropy of Secret [RYB+23]



s0 s1 s2 s3 s4 s5 s6 s7

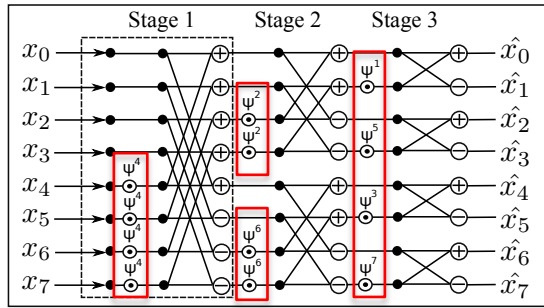
s0 s1 s0 s1 s0 s1 s0 s1

The effective secret s^*

s0 s1 0 0 0 0 0 0

- Secret (s^*) in NTT domain is used for Decaps
 - Kyber saves secret in NTT domain
 - To avoid extra NTT/INTT conversions
- Originally sampled secret s is forgotten!!!
 - Memoryless property of Kyber
 - Not applicable to Dilithium
- Attack also applies to masked implementations

FIA on KeyGen: Reduce Entropy of Secret [RYB+23]



- Sanity Check on Twiddle Constants:
 - Check Arithmetic Properties of Twiddle Constants:
 - n^{th} root of unity
 - Check Entropy of Twiddle Constants
- Not rely on single base address to access Twiddle Constant Array
- Check Entropy of NTT output

FIA on Kyber KeyGen/Encaps: Summary

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| Nonce_Fault_Attack | Observe_DUT_IO | Data /Control | Targeted | k | 1 | Check equality of s and e Copy Public Key if Pass |
| NTT_Fault_Attack | Observe_DUT_IO | Data | Targeted | 1 | 1 | Sanity Check on Twiddle Constants or NTT outputs |

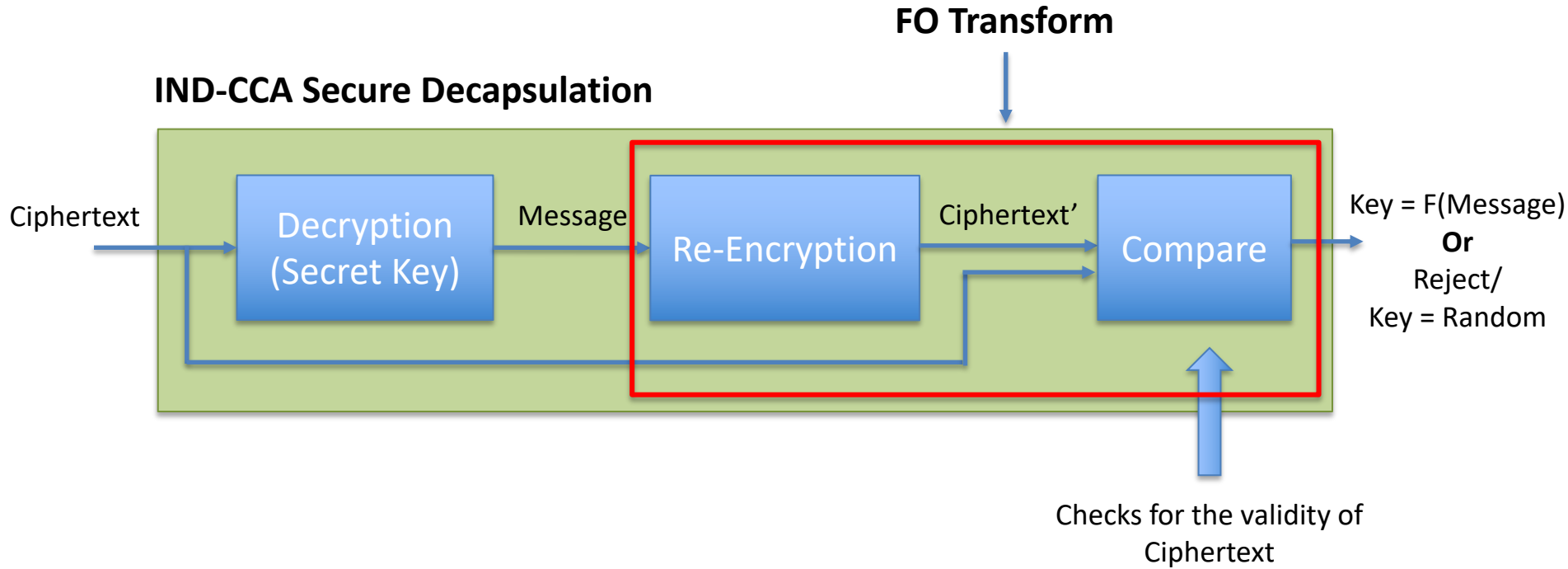
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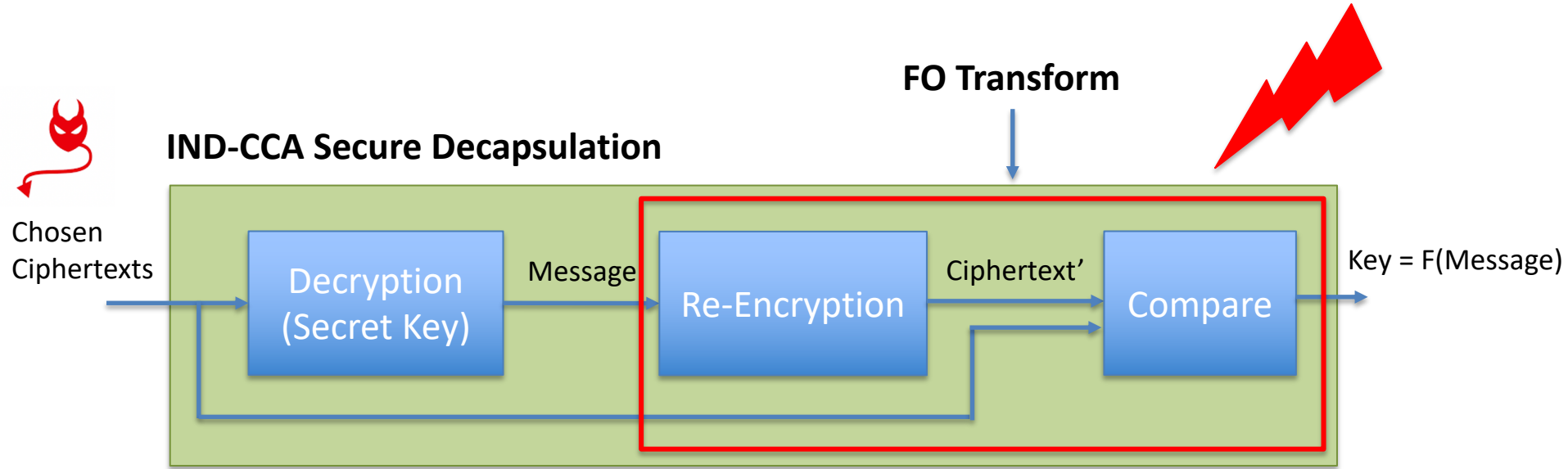
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FIA on Decapsulation



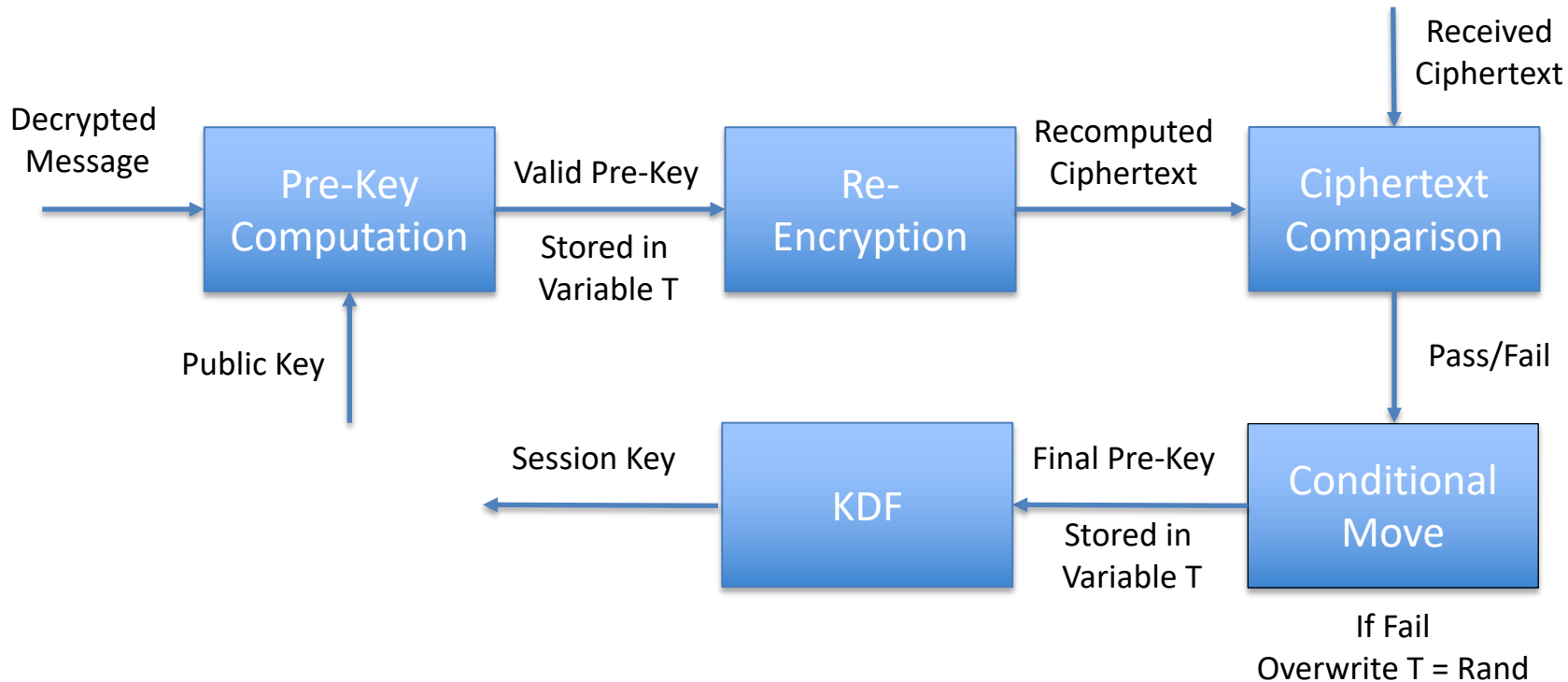
FIA on Decapsulation: Skip CT Comparison [XIU+21]



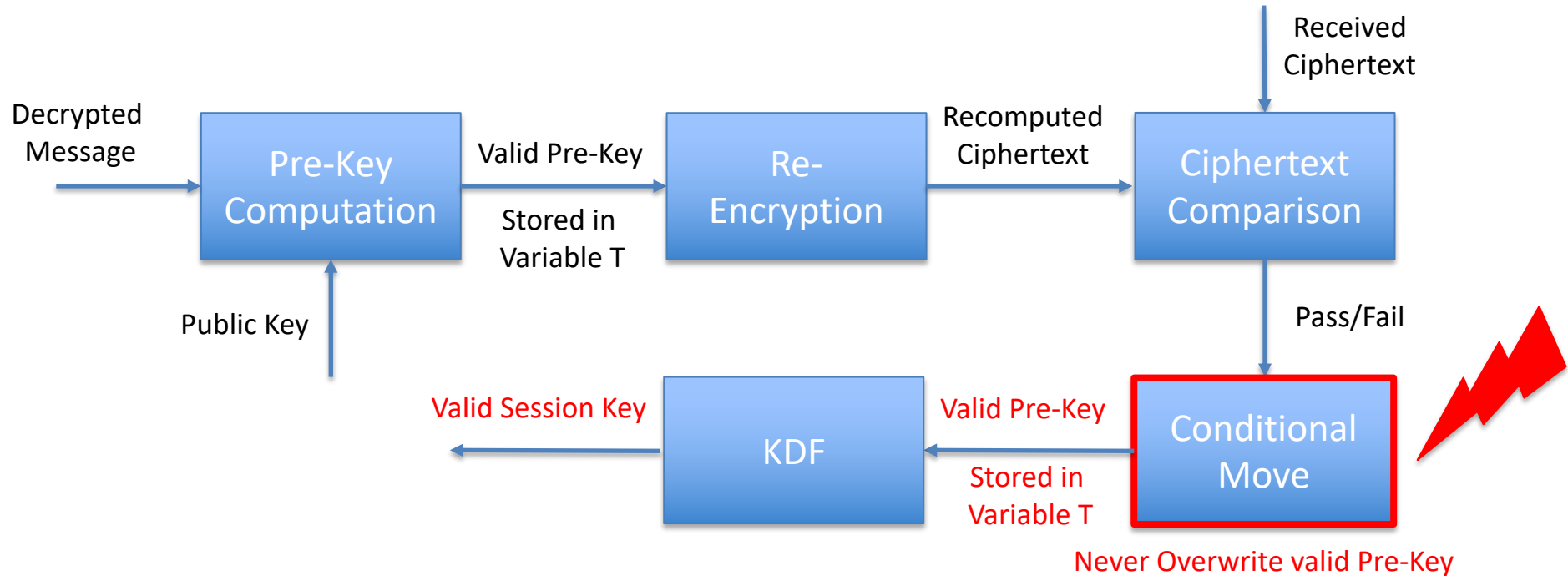
Bypass Comparison → **Downgrade from CCA to CPA Security**

This attack could have been easily avoided with a more careful implementation

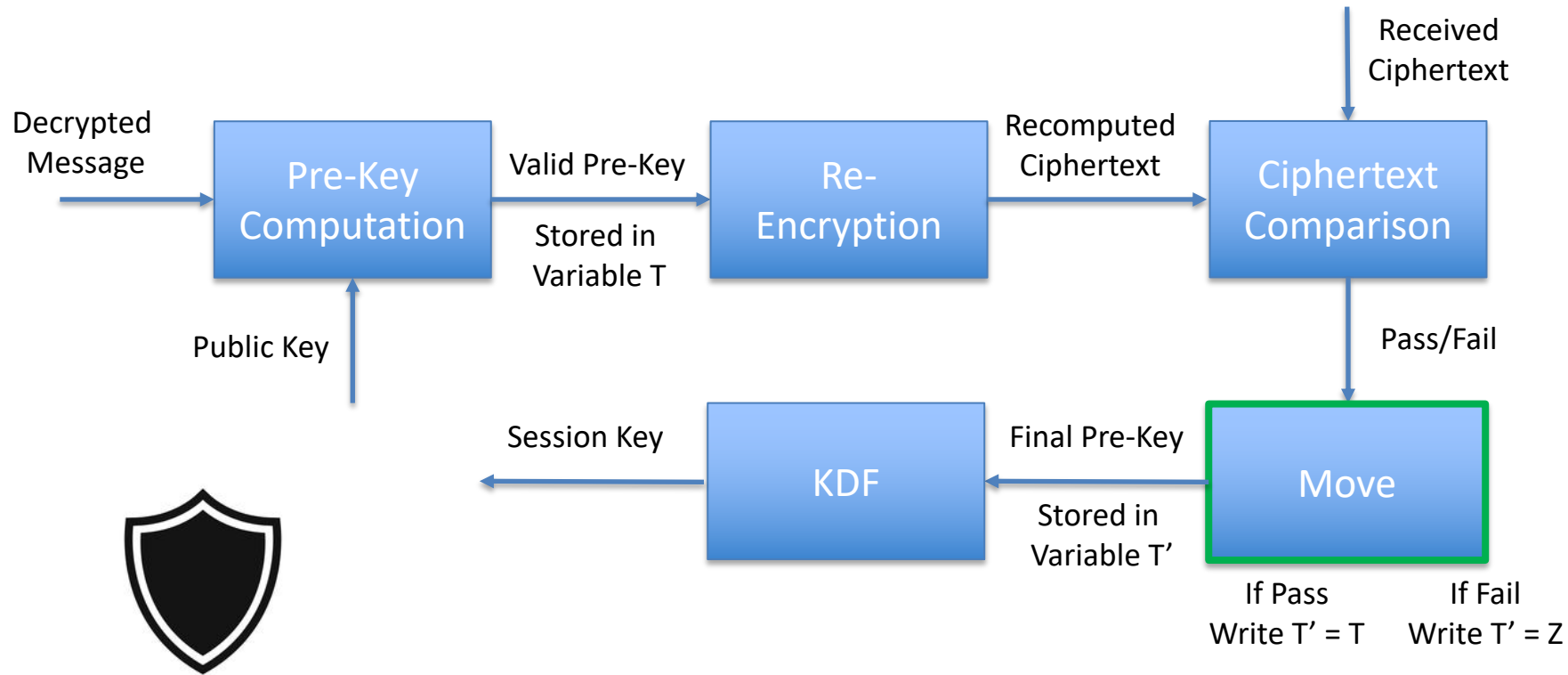
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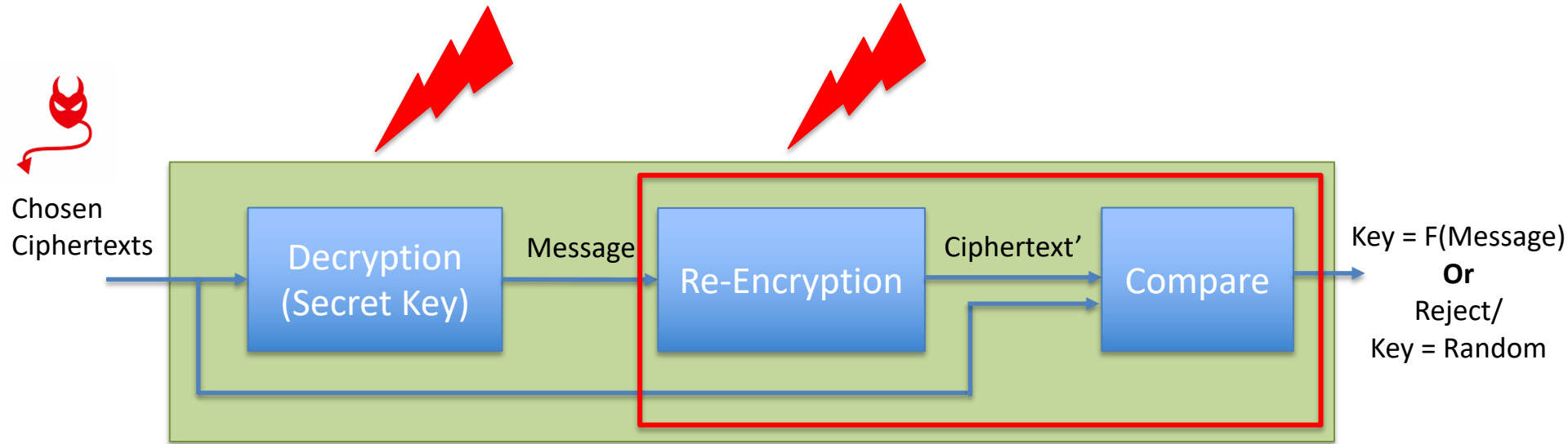


[XIU+21] Xagawa, Keita, Akira Ito, Rei Ueno, Junko Takahashi, and Naofumi Homma. "Fault-injection attacks against NIST's post-quantum cryptography round 3 KEM candidates." In *Advances in Cryptology—ASIACRYPT 2021: 27th International Conference on the Theory and Application of Cryptology and Information Security, Singapore, December 6–10, 2021, Proceedings, Part II* 27, pp. 33-61. Springer International Publishing, 2021.

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|------------------------|--------------------|---------------|-----------------|---|----------------------------------|----------------------------|
| Skip_CT_Compare | Communicate_DUT_IO | Control | Targeted | 1 | Few thousand. | Protected Conditional Move |

FIA on Decapsulation: Fault Assisted CCA [PP21,HPP21,D22]



Modus Operandi: Inject Faults to realize a Decryption Failure (DF) Oracle

[PP21] Pessl, Peter, and Lukas Prokop. "Fault attacks on CCA-secure lattice KEMs." *IACR Transactions on Cryptographic Hardware and Embedded Systems* (2021): 37-60.

[HPP21] Hermelink, Julius, Peter Pessl, and Thomas Pöppelmann. "Fault-enabled chosen-ciphertext attacks on Kyber." In *Progress in Cryptology—INDOCRYPT 2021: 22nd International Conference on Cryptology in India, Jaipur, India, December 12–15, 2021, Proceedings 22*, pp. 311-334. Springer International Publishing, 2021.

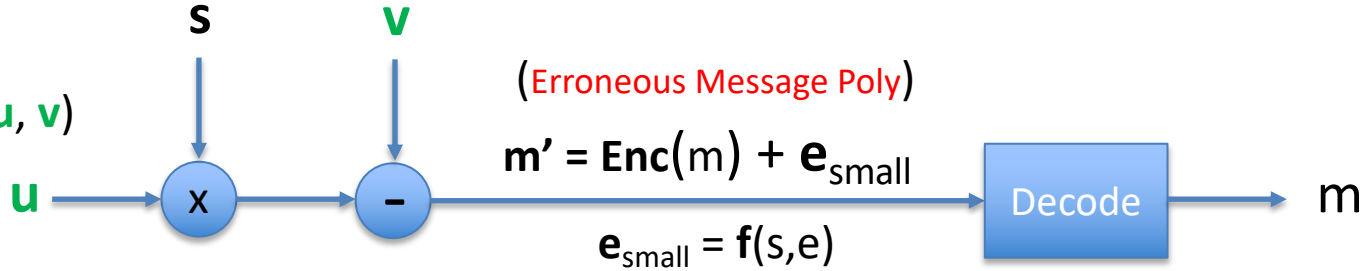
[D22] Delvaux, Jeroen. "Roulette: A Diverse Family of Feasible Fault Attacks on Masked Kyber." *IACR Transactions on Cryptographic Hardware and Embedded Systems* (2022)

FIA on Decapsulation: Fault Assisted CCA [PP21,HPP21,D22]

Decryption

Chosen

Ciphertext = (u, v)



(Erroneous Message Poly)

$$m' = \text{Enc}(m) + e_{\text{small}}$$

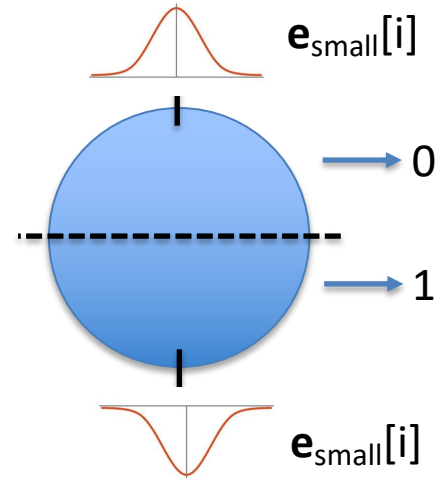
$$e_{\text{small}} = f(s, e)$$

Linear Relation

Knowledge of $e_{\text{small}}[i] < 0$ or $e_{\text{small}}[i] > 0$
for chosen-ciphertexts

Modus Operandi:

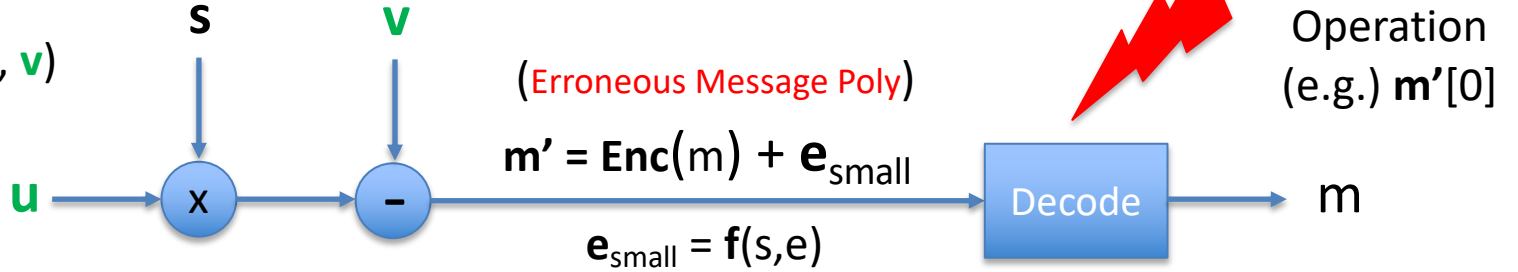
Use faults to learn $e_{\text{small}}[i] < 0$ or $e_{\text{small}}[i] > 0$
for chosen-ciphertexts



FIA on Decapsulation: Ineffective Fault Analysis [PP21]

Decryption

Ciphertext = (u, v)



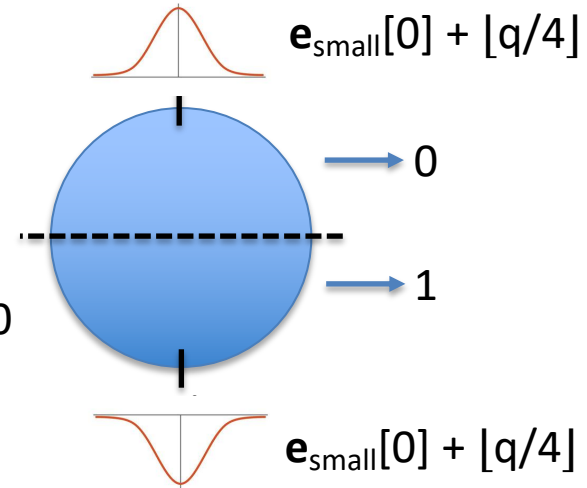
Linear Relation

Effective Fault

Bit Flip when $e_{\text{small}}[0] < 0$
Decapsulation Failure

Ineffective Fault

No Bit Flip when $e_{\text{small}}[0] > 0$
Decapsulation Success



Can also tolerate error (1%)

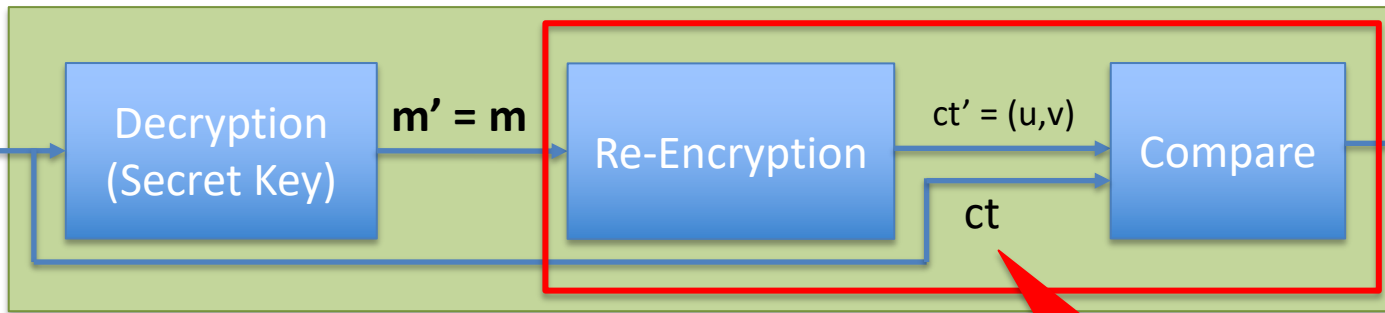
FIA on Kyber Decaps: Summary

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|------------------------|--------------------|---------------|---------------------|--|-------------------------------------|-------------------------------|
| Skip_CT_Compare | Communicate_DUT_IO | Control | Targeted | 1 | Few thousands (1k-3k) | Protected Conditional Move |
| Ineffective_FIA | Communicate_DUT_IO | Control | Targeted | 1 | Few thousands (5k-7k) | Shuffle Message Decoding |

FIA on Decapsulation: Fault Correction Attack [HPP21]

CCA Secure Decapsulation

Only differ by single bit



Success

$ct = (u, v + [q/4]. x^0)$

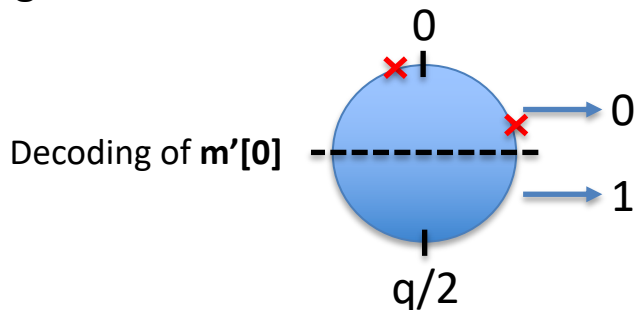
$m' = m$

$ct' = (u, v)$

ct

No Decryption Failure

Single bit change
in valid ct

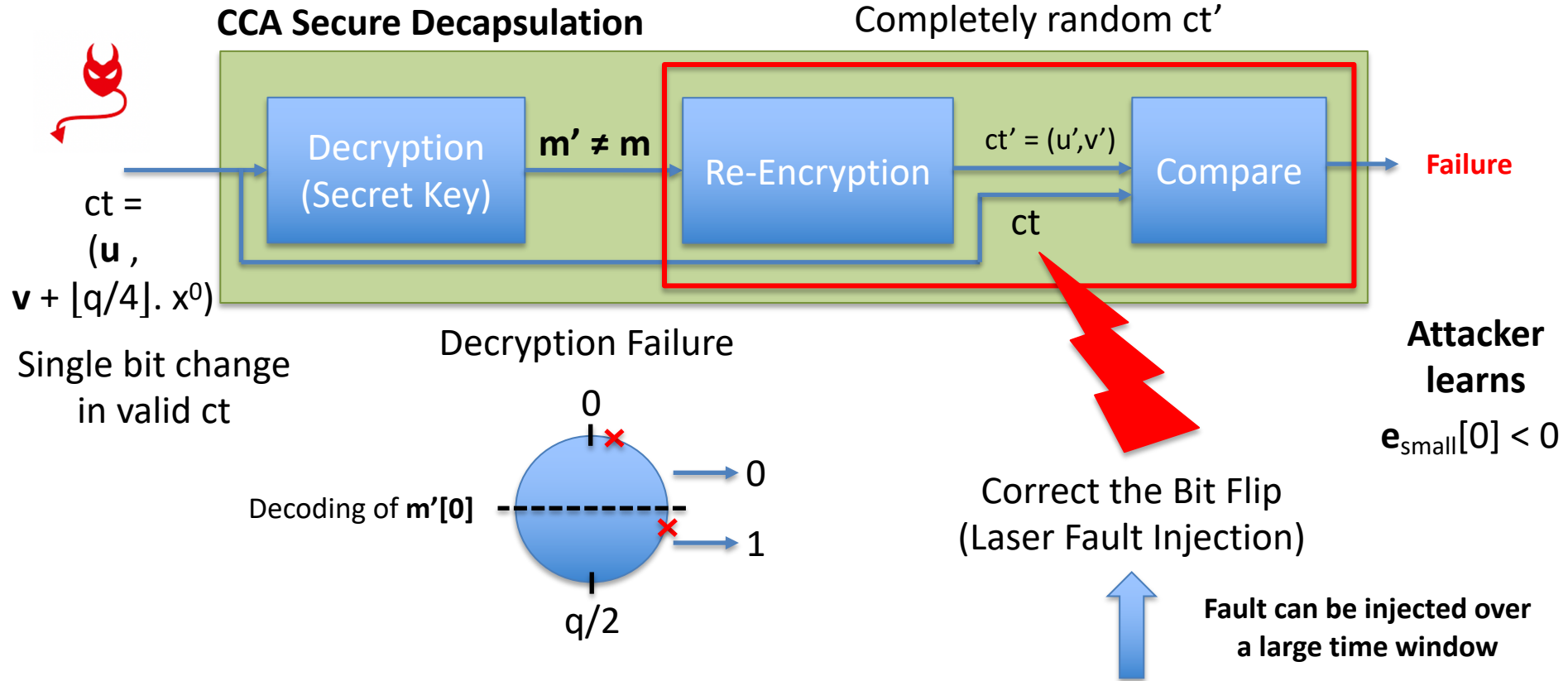


Correct the Bit Flip
(Laser Fault Injection)

Attacker
learns

$e_{\text{small}}[0] < 0$

FIA on Decapsulation: Fault Correction Attack [HPP21]



FIA on Decapsulation: Fault Correction Attack [HPP21]



- Compute Hash of Ciphertext and Compare with Hash
 - Inject fault before ciphertext sent to hash function
 - Reduces Attack Surface, but does not prevent attack
- Attack of [HPP21] improved in Roulette Attack [D22]:
 - **Larger Attack Surface:**
 - Operations in Re-Encryption: Sampler, INTT
 - **Relaxed Fault Models:**
 - Set-to-0, random faults, arbitrary bit flips, instruction skips
 - **Allows larger errors:**
 - $\cong 25\%$

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| Ineffective_FIA | Communicate_DUT_IO | Control | Targeted | 1 | Few thousands (5k-7k) | Shuffle Message Decoding |
| Fault_Correction_Attack_1 | Communicate_DUT_IO | Control | Targeted (Bit Flip) | 1 | Few thousands (5k-7k) | Redundancy |
| Fault_Correction_Attack_2 | Communicate_DUT_IO | Control | Targeted (More Relaxed Fault Models) | 1 | Few tens-hundred thousands (10k-100k) | Redundancy |

Open Research Directions: Development of Algorithmic Countermeasures against fault assisted DF oracle attacks

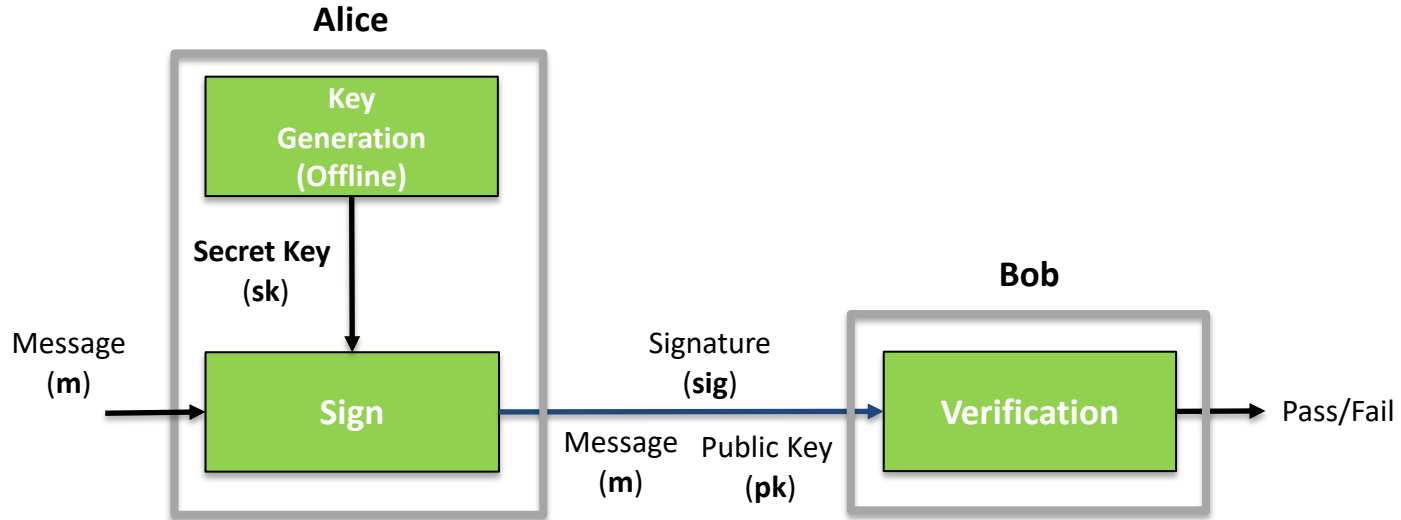
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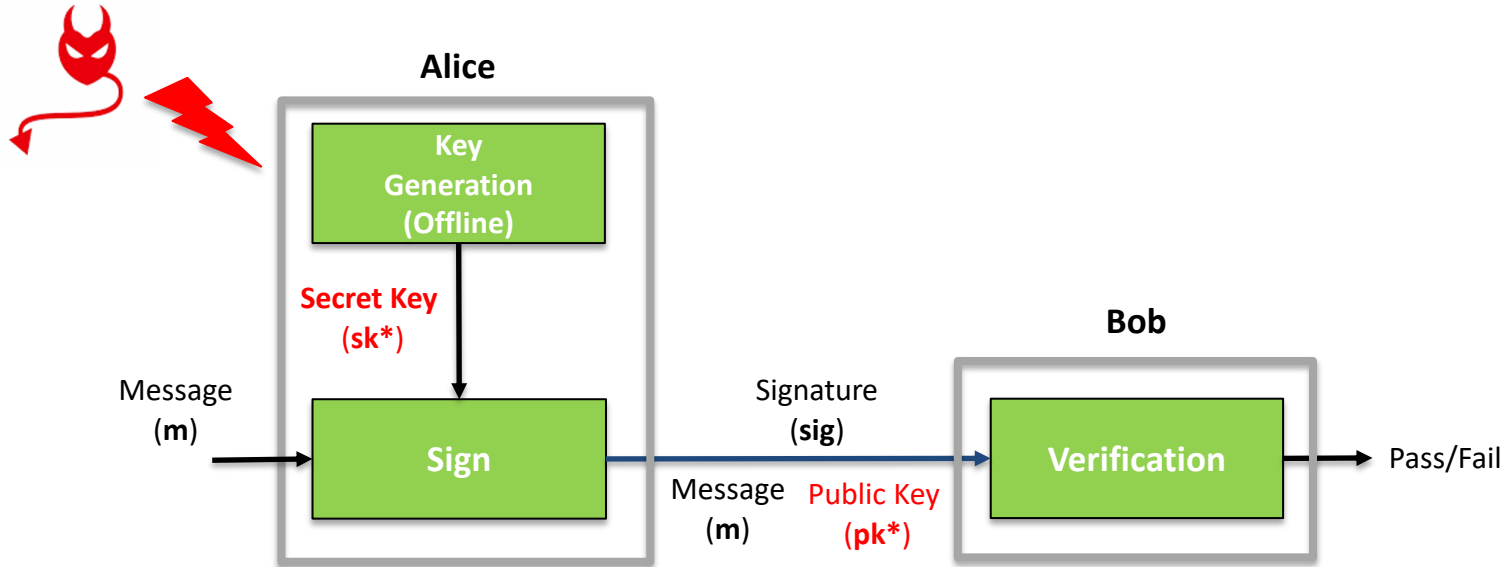
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- ❑ Conclusion

Signature Scheme: Background

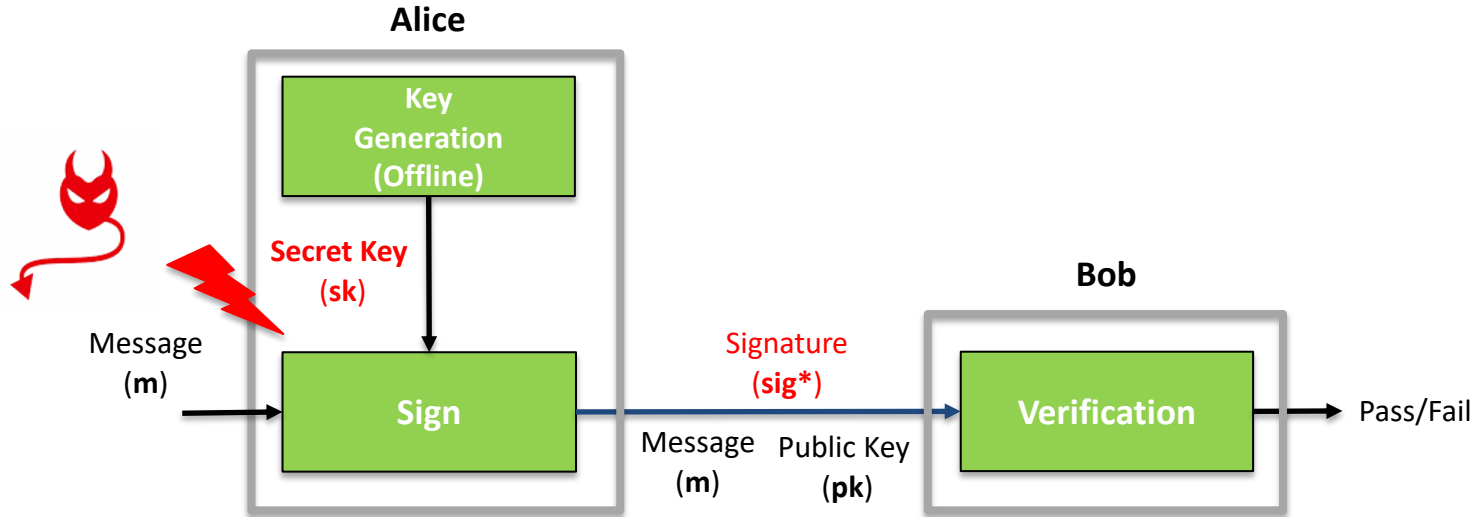


Signature Scheme: Attacking Alice



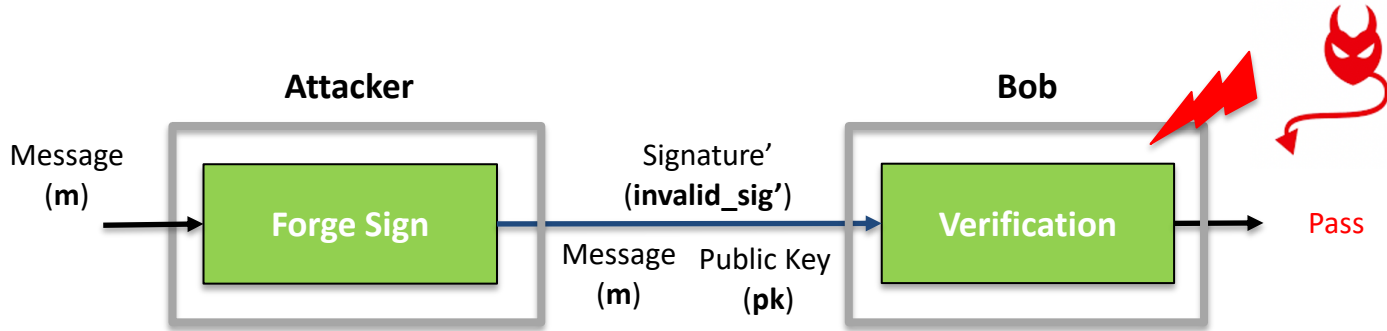
- ❑ Single execution to target Key Generation: Key Recovery Attack
 - ❑ Recover Secret key from Faulty but valid Public Key
 - ❑ But, key generation mostly done offline (PKI Infrastructure)

Signature Scheme: Attacking Alice



- ❑ Multiple Executions to target Signing Procedure:
 - ❑ Single/Multiple faulty signatures used to recover secret key
 - ❑ Most attractive target for attacker

Signature Scheme: Attacking Bob



- ❑ Single Execution to bypass Signature Verification:
 - ❑ Application: Secure Boot

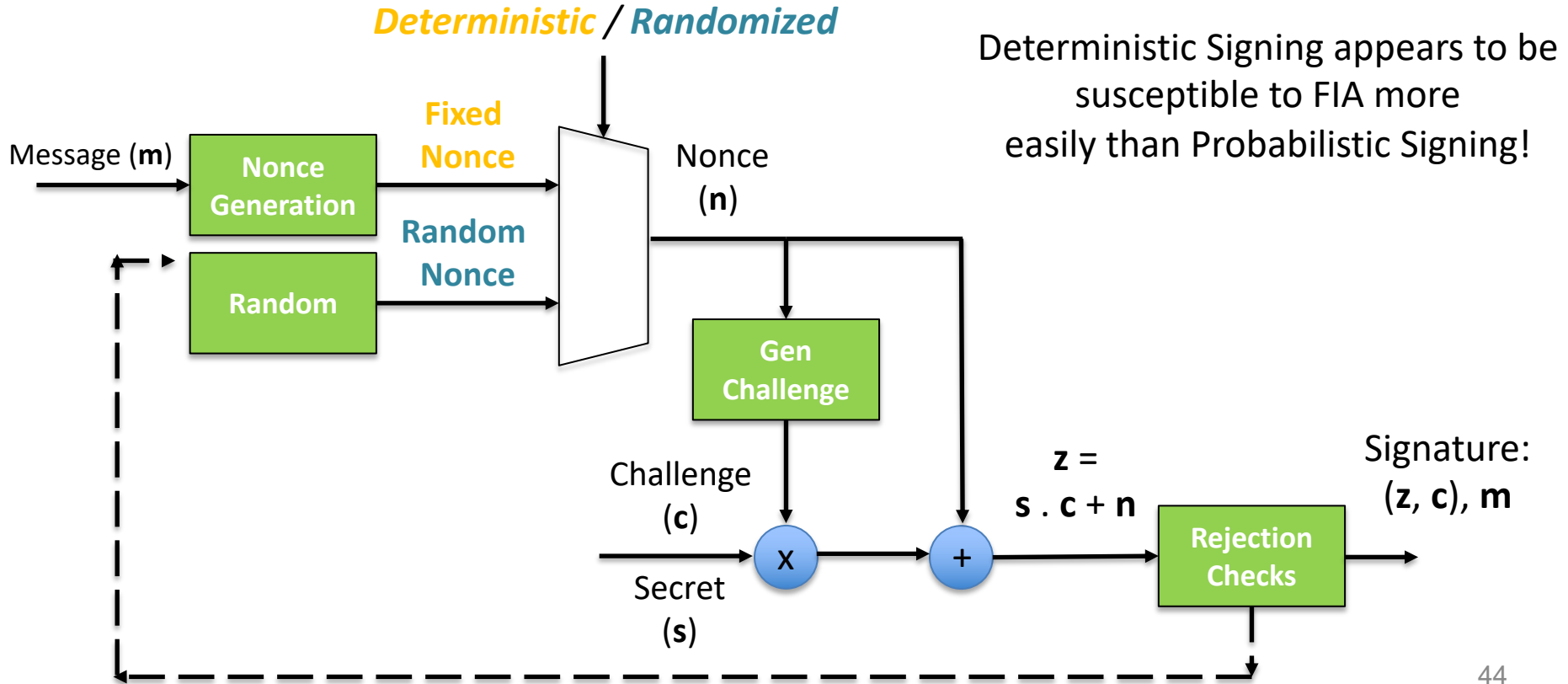
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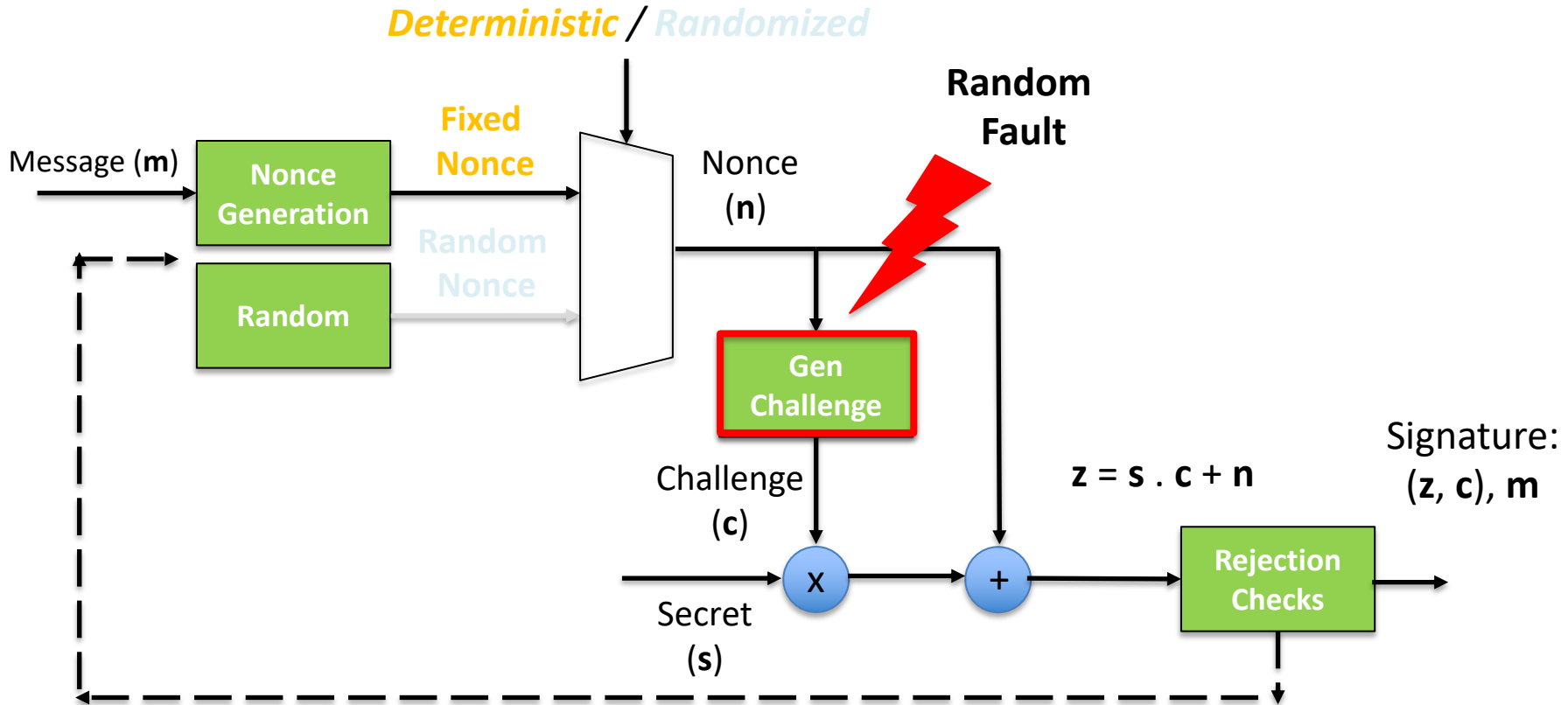
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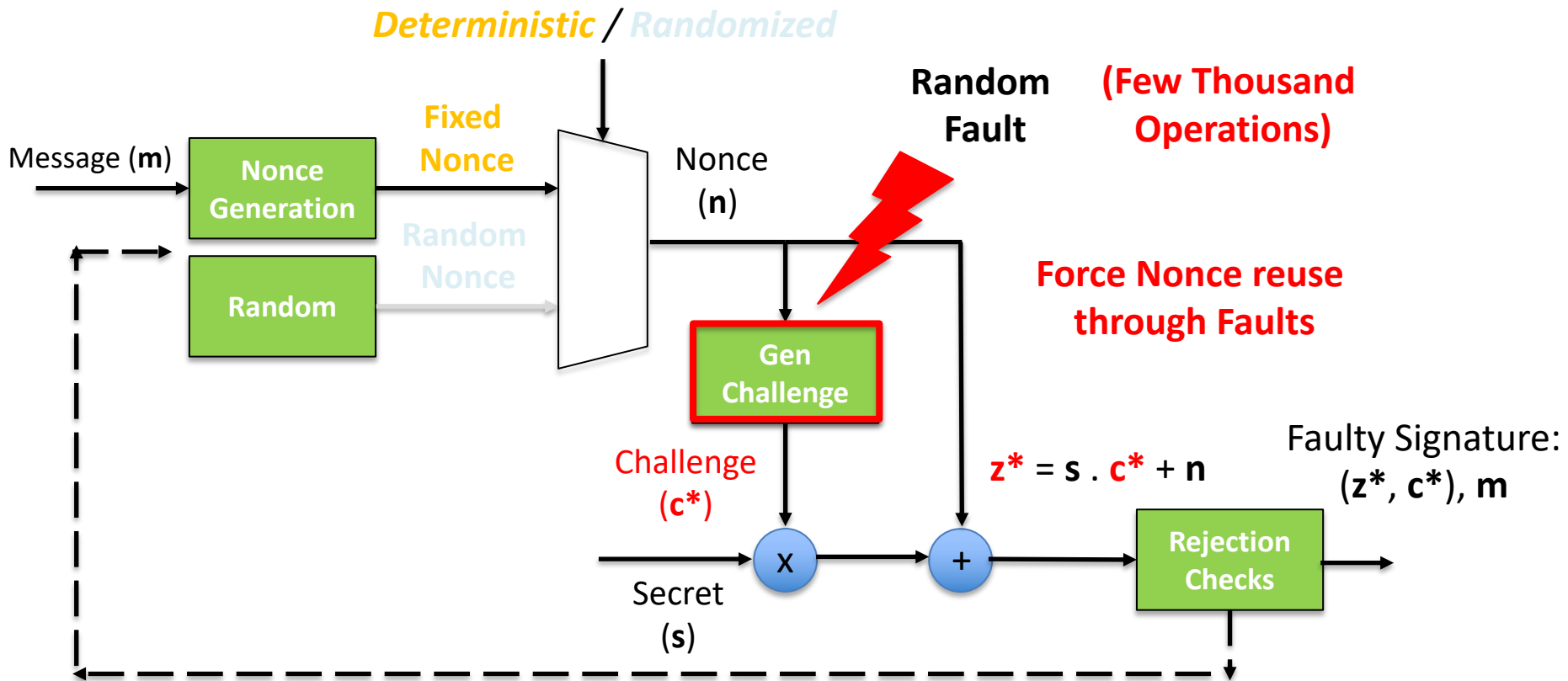
FIA on Signing Procedure: Background



FIA on Signing Procedure: DFA (Deterministic) [BP18]



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FIA on Signing Procedure: DFA (Deterministic) [BP18]

Valid Sig.

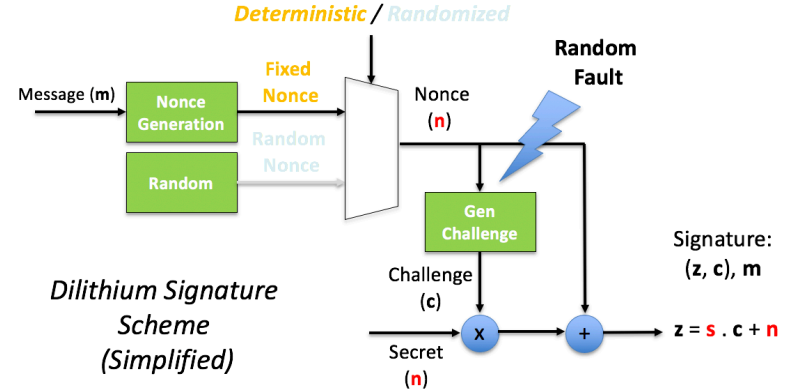
$$z = s \cdot c + n$$

$$(z - z^*) = s \cdot (c - c^*)$$

$$s = (z - z^*) \cdot (c - c^*)^{-1}$$

Faulty Sig.

$$z^* = s \cdot c^* + n$$

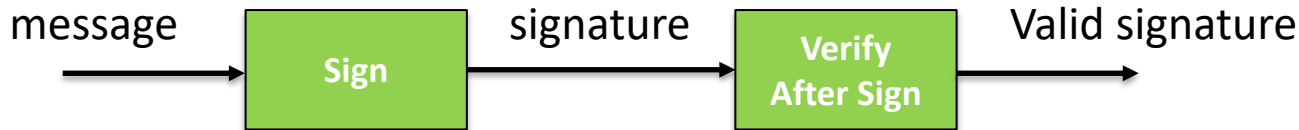


60% operations are vulnerable to DFA

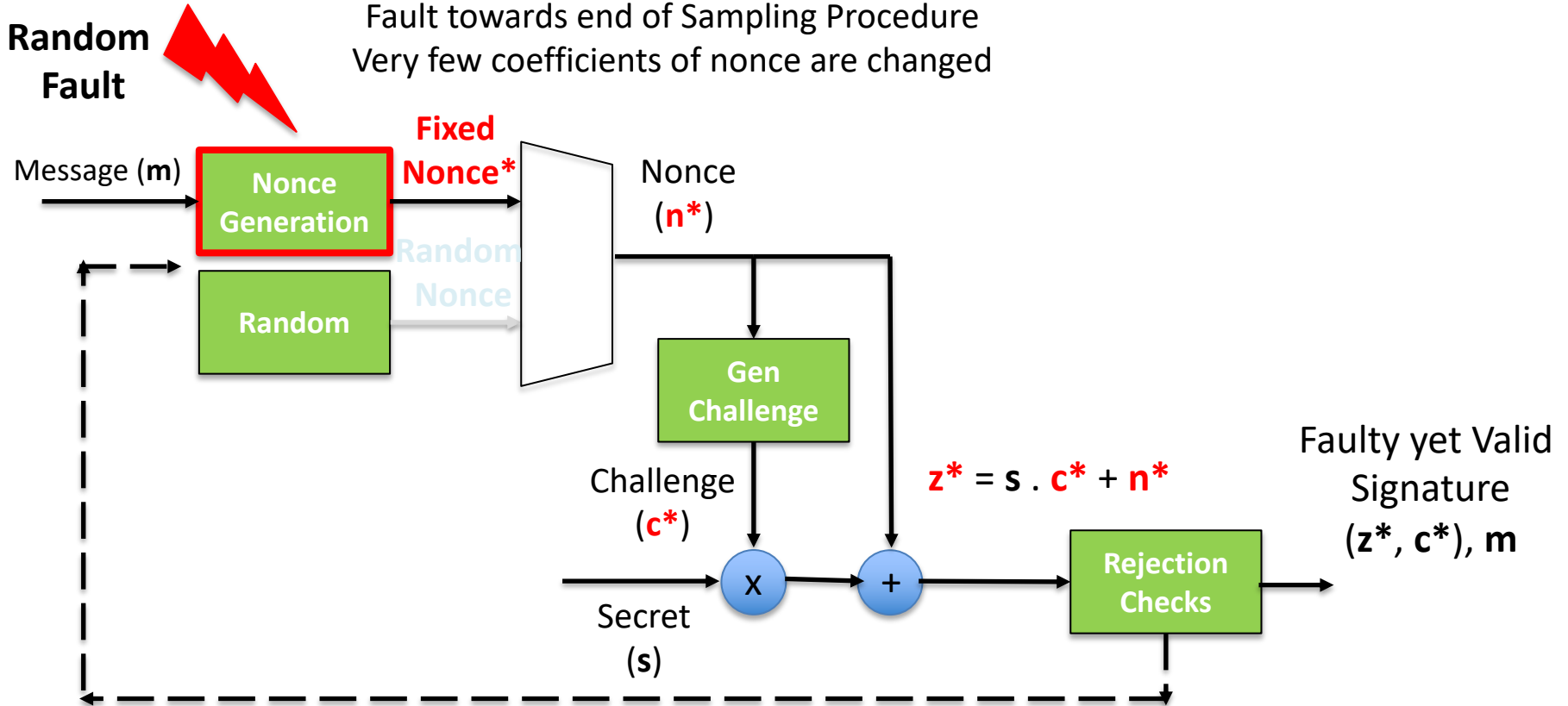


Mitigation

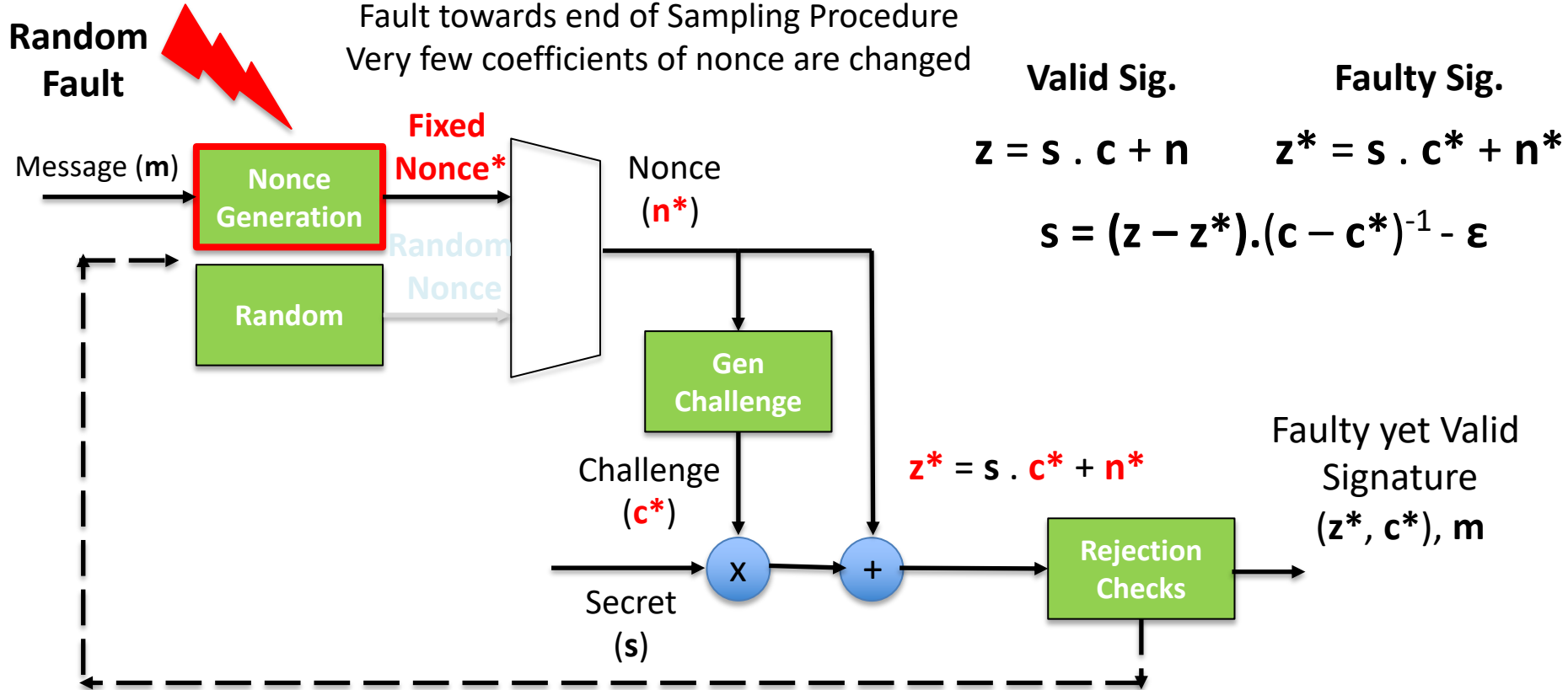
Verify the generated Signatures



FIA on Signing Procedure: DFA (Deterministic) [BP18]

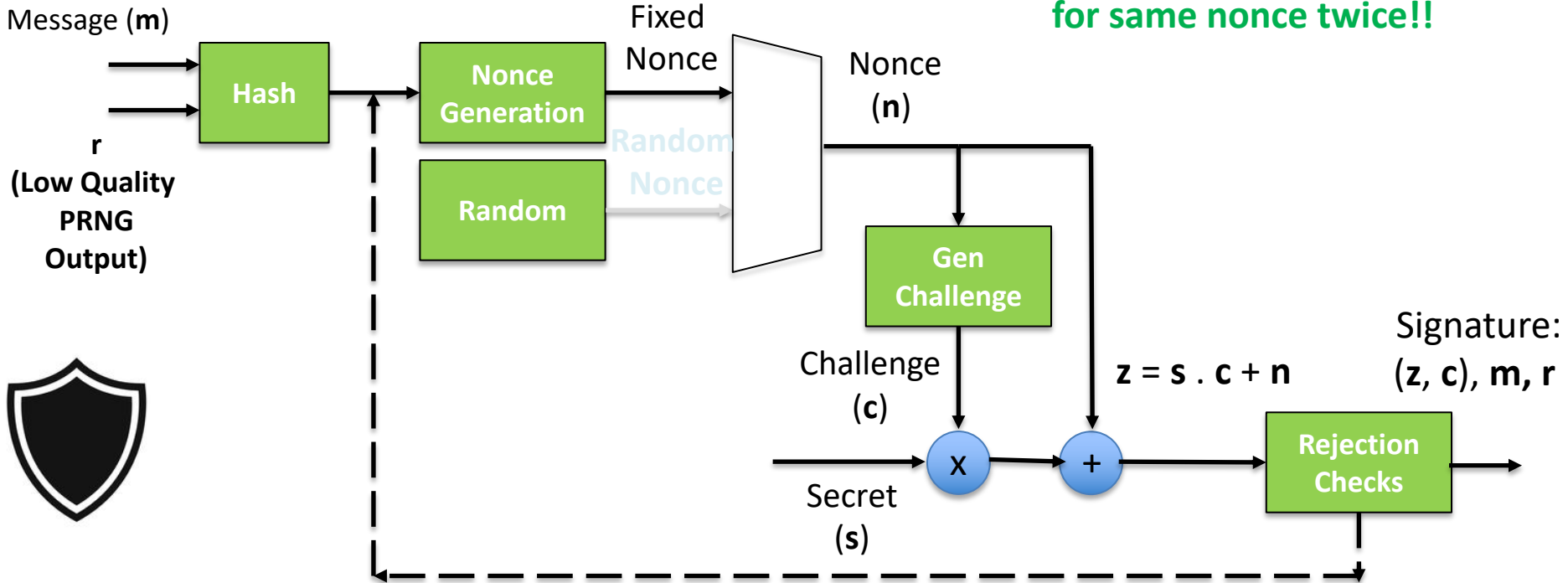


FIA on Signing Procedure: DFA (Deterministic) [BP18]



FIA on Signing Procedure: DFA (Deterministic)

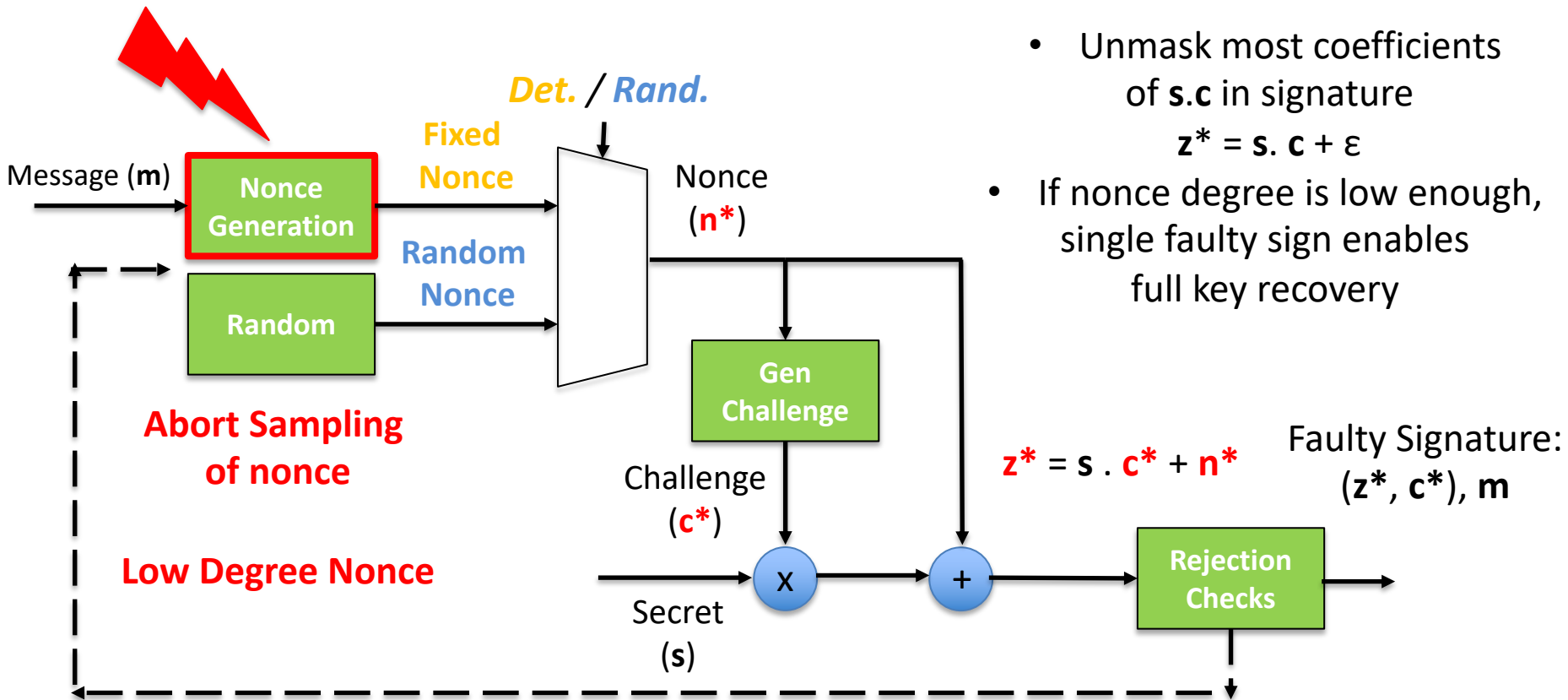
New Proposal (This Talk):



FIA on Dilithium Signing: Summary

| Attack Name | DUT_IO_Access | Type of Fault | Targeted_or_Not | No. of Faults within Single Computation | Total No. of Faulty Computations | Countermeasure |
|--------------------|--------------------|---------------|-----------------|---|----------------------------------|--|
| Generic_DFA | Communicate_DUT_IO | Data/Control | Non_Targeted | 1 | 1 | Verify_after_Sign Redundancy (Nonce Sample) |

FIA on Signing Procedure: Loop Abort Fault [EFG+18]



- Unmask most coefficients of s.c in signature

$$z^* = s \cdot c + \epsilon$$

- If nonce degree is low enough, single faulty sign enables full key recovery

FIA on Signing Procedure: Loop Abort Fault [EFG+18]

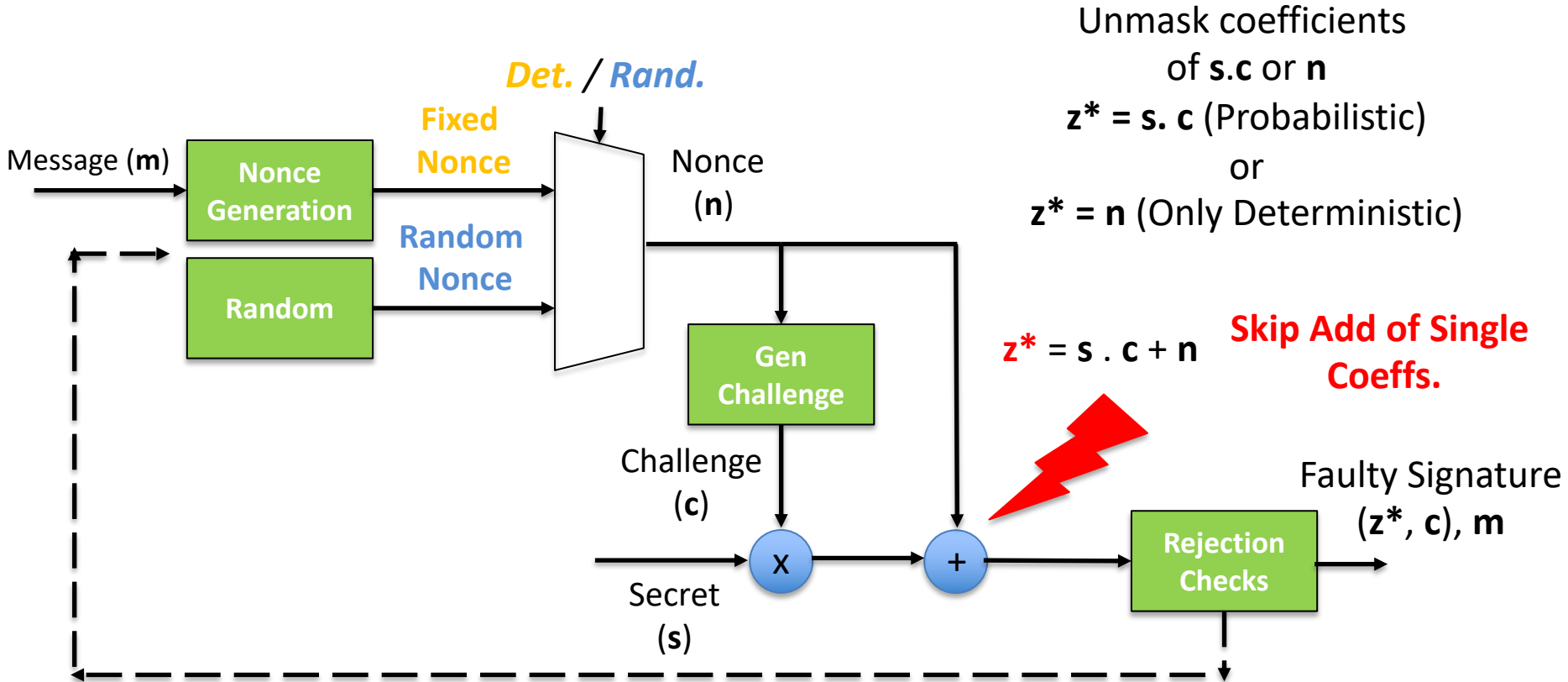


- Initializing nonce to random values and then sample
- Sanity Check of Nonce
 - Check for high number of zeros
 - Entropy Check

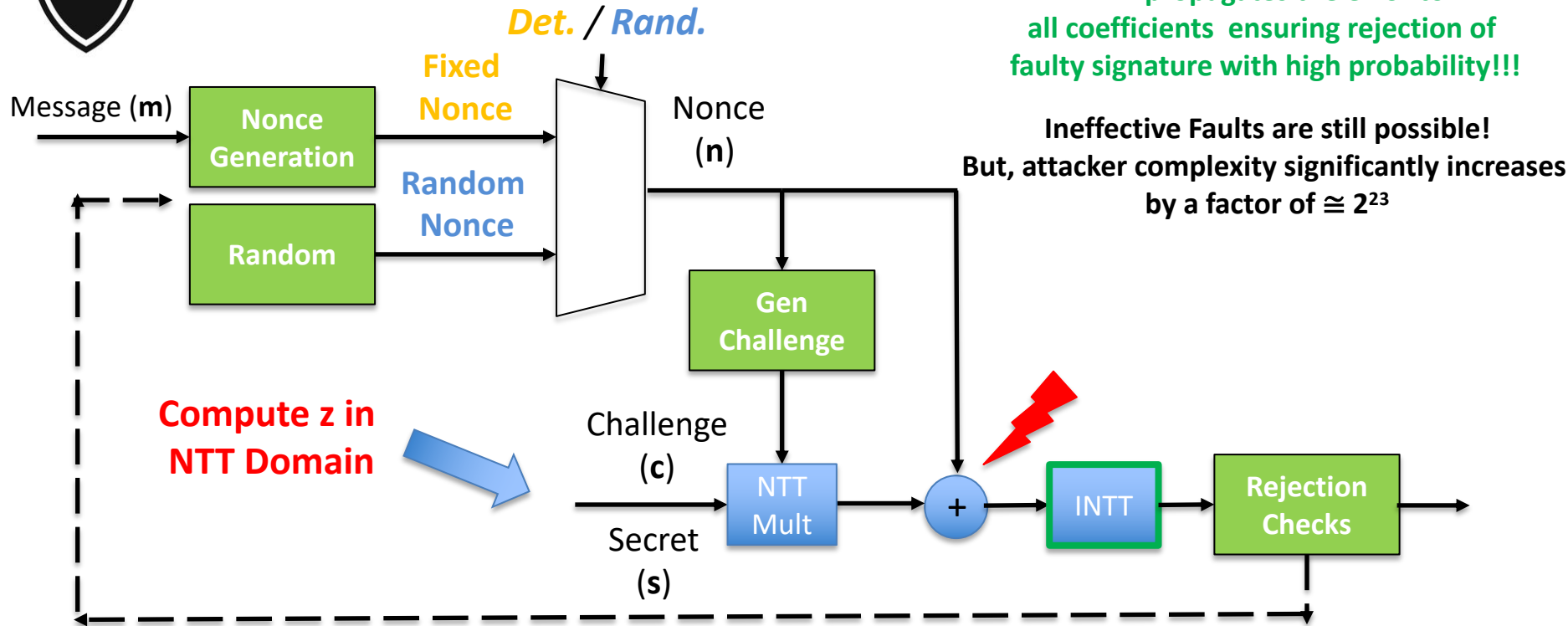
FIA on Dilithium Signing: Summary

| Attack Name | DUT_IO_Access | Type of Fault | Targeted_or_Not | No. of Faults within Single Computation | Total No. of Faulty Computations | Countermeasure |
|--------------------|--------------------|---------------|-----------------|---|----------------------------------|---|
| Generic_DFA | Communicate_DUT_IO | Data/Control | Non_Targeted | 1 | 1 | Verify_after_Sign Redundancy (Nonce Sample) |
| Loop_Abort | Observe_DUT_IO | Control | Targeted | 1 | 1 | Redundancy (Nonce Sample) Sanity Check on Nonce Random Initialization |

FIA on Signing Procedure: Skip Addition [RJH+19]



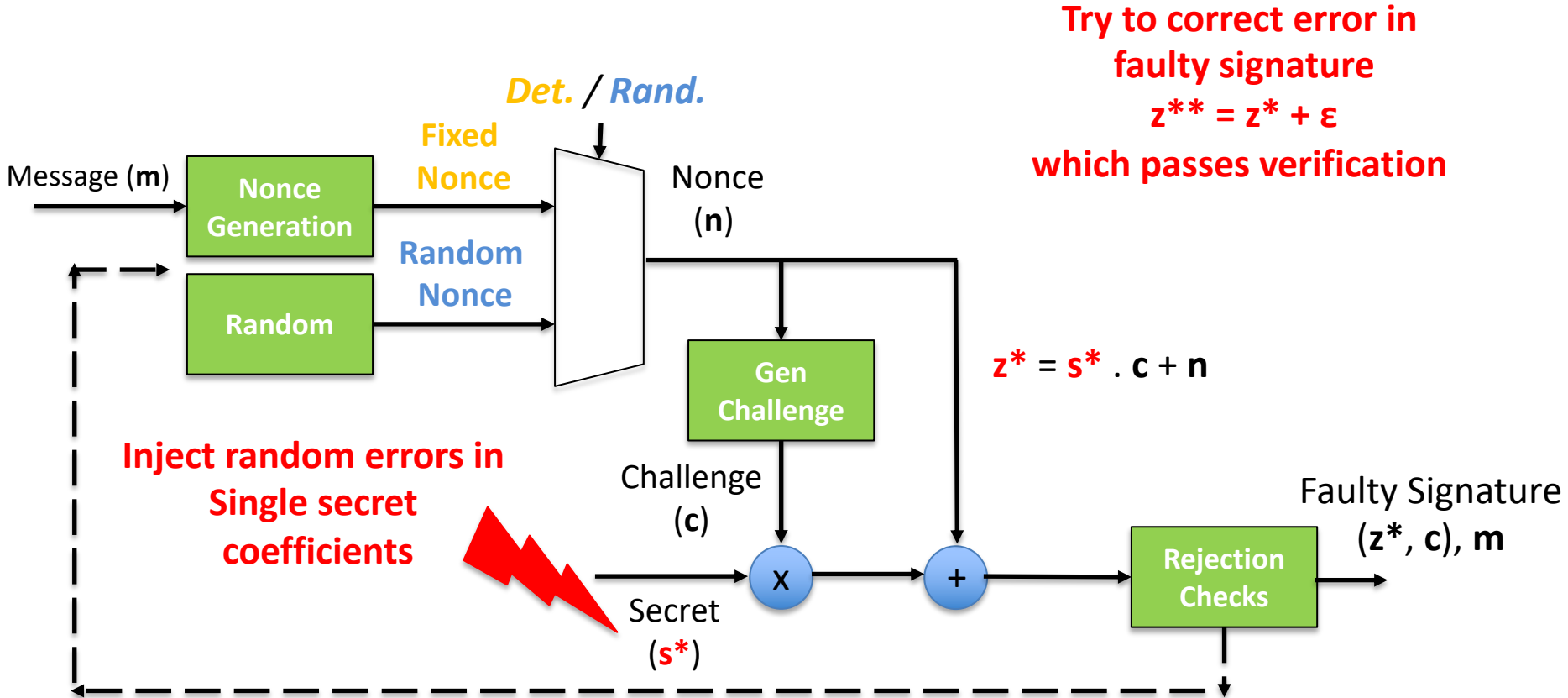
FIA on Signing Procedure: Skip Addition [RJH+19]



FIA on Dilithium Signing: Summary

| Attack Name | DUT_IO_Access | Type of Fault | Targeted_or_No t | No. of Faults within Single Computation | Total No. of Faulty Computations | Countermeasure |
|----------------------|-----------------------------------|---------------|------------------|---|----------------------------------|---|
| Generic_DFA | Communicate_DUT_IO | Data/Control | Non_Targeted | 1 | 1 | Verify_after_Sign Redundancy (Nonce Sample) |
| Loop_Abort | Observe_DUT_IO | Control | Targeted | 1 | 1 | Redundancy (Nonce Sample) Sanity Check on Nonce Random Initialization |
| Skip_Addition | Communicate_DUT_IO/Observe_DUT_IO | Data/Control | Targeted | 1 | Few thousand | Verify After Sign Redundancy (Final Addition) NTT Addition |

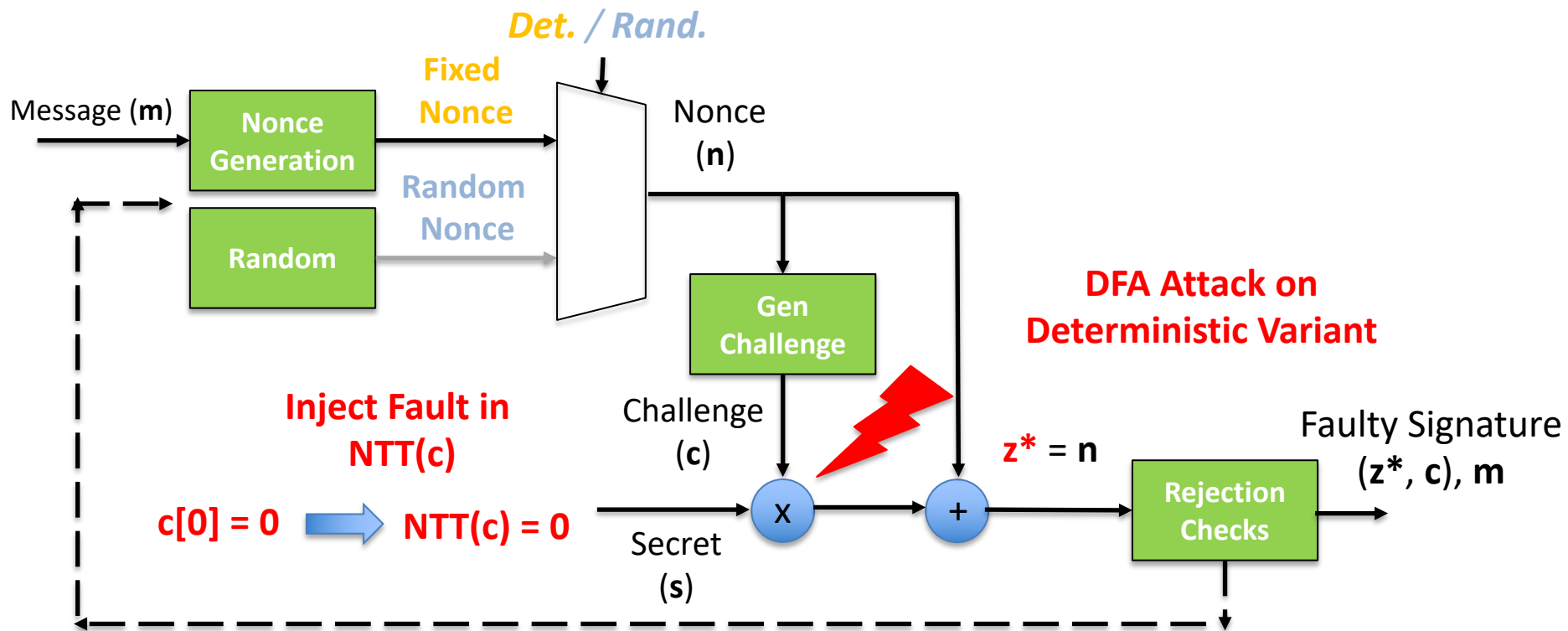
FIA on Signing Procedure: Erroneous Secret Key [IMS+22]



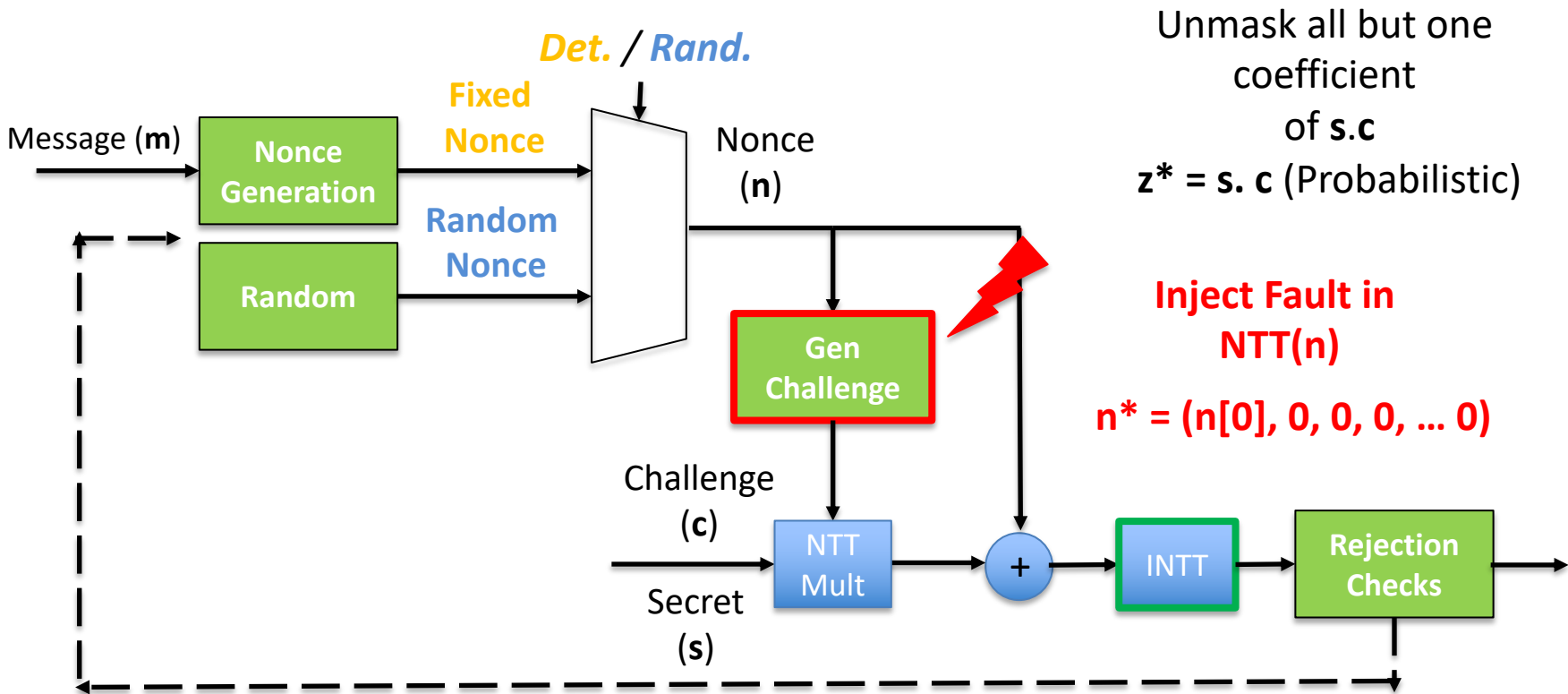
FIA on Dilithium Signing: Summary

| Attack Name | DUT_IO_Access | Type of Fault | Targeted_or_No t | No. of Faults within Single Computation | Total No. of Faulty Computations | Countermeasure |
|-----------------------------|-----------------------------------|---------------|------------------|---|----------------------------------|---|
| Generic_DFA | Communicate_DUT_IO | Data/Control | Non_Targeted | 1 | 1 | Verify_after_Sign Redundancy (Nonce Sample) |
| Loop_Abort | Observe_DUT_IO | Control | Targeted | 1 | 1 | Redundancy (Nonce Sample) Sanity Check on Nonce Random Initialization |
| Skip_Addition | Communicate_DUT_IO/Observe_DUT_IO | Data/Control | Targeted | 1 | Few thousand | Verify After Sign Redundancy (Final Addition) NTT Addition |
| Erroneous_Secret_Key | Observe_DUT_IO | Data | Targeted | 1 | Few thousand | Verify After Sign |

FIA on Signing Procedure: NTT Fault Attack [RYB+23]



FIA on Signing Procedure: NTT Fault Attack [RYB+23]



FIA on Dilithium Signing: Summary

| Attack Name | DUT_IO_Access | Type of Fault | Targeted_or_No_t | No. of Faults within Single Computation | Total No. of Faulty Computations | Countermeasure |
|-----------------------------|-----------------------------------|---------------|------------------|---|----------------------------------|---|
| Generic_DFA | Communicate_DUT_IO | Data/Control | Non_Targeted | 1 | 1 | Verify_after_Sign Redundancy (Nonce Sample) |
| Loop_Abort | Observe_DUT_IO | Control | Targeted | 1 | 1 | Redundancy (Nonce Sample) Sanity Check on Nonce Random Initialization |
| Skip_Addition | Communicate_DUT_IO/Observe_DUT_IO | Data/Control | Targeted | 1 | Few thousand | Verify After Sign Redundancy (Final Addition) NTT Addition |
| Erroneous_Secret_Key | Observe_DUT_IO | Data | Targeted | 1 | Few thousand | Verify After Sign |
| NTT_Fault_Attack | Communicate_DUT_IO/Observe_DUT_IO | Data | Targeted | 1 | 1 | Sanity Check on Twiddle Constants or NTT outputs |

Outline

- ❑ FIA on Kyber:
 - ❑ FIA on Key Generation
 - ❑ FIA on Decapsulation

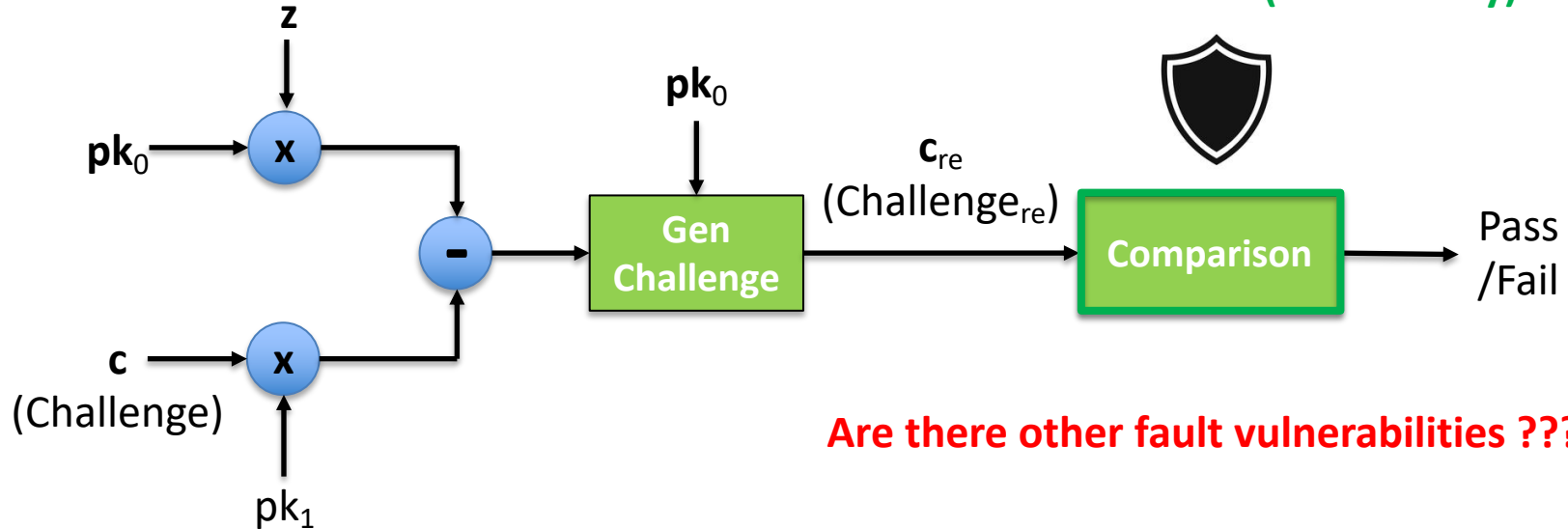
- ❑ **FIA on Dilithium**
 - ❑ FIA on Signing
 - ❑ **FIA on Verification**

- ❑ Conclusion

FIA on Verification Procedure

Signature: (z, c) , m

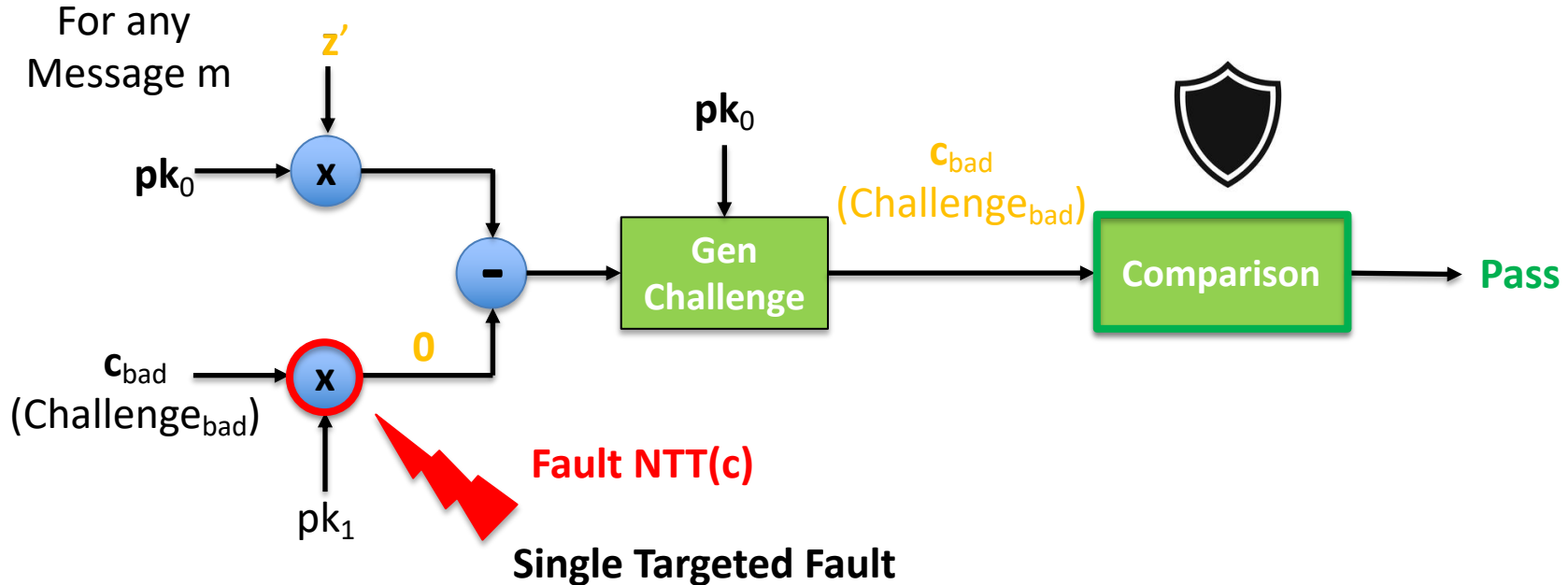
Public Key: (pk_0, pk_1)



FIA on Verification: NTT Fault Attack [RYB⁺23]

Signature: (z, c) , m

Public Key: (pk_0, pk_1)



Outline

- ❑ FIA on Kyber:
 - ❑ FIA on Key Generation
 - ❑ FIA on Decapsulation

- ❑ FIA on Dilithium
 - ❑ FIA on Signing
 - ❑ FIA on Verification

- ❑ **Conclusion and Open Questions**

Conclusion

- ❑ Variety of FIA are possible on unprotected Kyber and Dilithium
 - ❑ Non-Monolithic nature of scheme: Arithmetic Domain, Boolean Domain
- ❑ Underlying algorithmic features and implementation choices render them susceptible to SCA and FIA
- ❑ Questions:
 - ❑ Should we use dedicated countermeasures for every attack?
 - ❑ Can we exploit algorithmic features for efficient and low-cost countermeasures?
- ❑ Refer to [\[RDB⁺22\]](#) for a systematic study of SCA, FIA of Kyber and Dilithium (NIST PQC Standards)

“In a way, these things are like gold nuggets that God left in the forest. If I'm walking along in the forest and I stubbed my toe on it, who's to say I deserve credit for discovering it?”

-- Dr. Martin Hellman on the discovery of Public-Key Cryptography

Thank you!

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