## Secure Multiparty Computation and Applications

Presenter: Steve Lu, CEO Stealth Software Technologies, Inc.

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#### Stealth Team – About Us

#### Our Mission

To help organizations maximize the utility of their data while mitigating security and privacy risks through applied cryptography and cybersecurity.



HQ in Los Angeles, co-founded by UCLA Distinguished Professor Rafail Ostrovsky. We additionally work with consultants who are world-renowned cryptographers.



Founded in 2006. We have competitively won over a dozen major contracts from leading agencies in cryptography.



Multiple breakthroughs in cryptography, matured for deployment, including privacy-preserving computation, analysis, and storage.



#### Centralized Systems

Registry



**Certificate Authority** 



Data Warehouse



# **Trust in Centralized Entity**

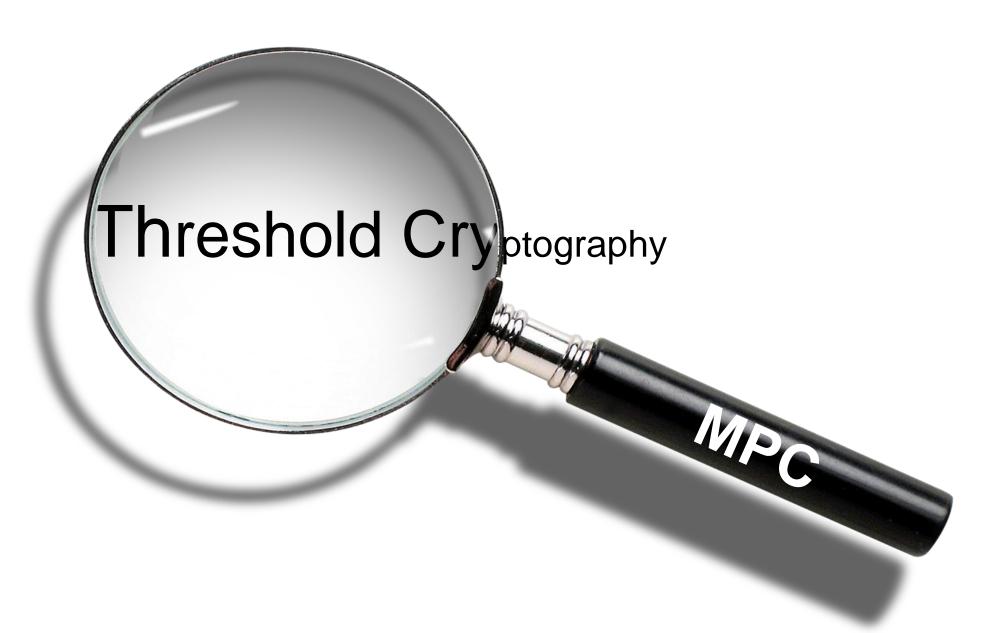


#### Secure Multiparty Computation (MPC) to the Rescue?

- MPC and Threshold Cryptography are clearly related, but often not mentioned in the same breath
- MPC usually deals with players who already have inputs and want to compute a joint (arbitrary) function
- Threshold cryptography wants to distribute an existing functionality – a specific goal for a specific problem



#### In This Talk









#### Simple Three-Round Multiparty Schnorr Sig with Full Simulatability

Yehuda Lindell

Coinbase

yehuda.lindell@gmail.com

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Abstr

Building on the Gennaro & Goldfeder and Lindell & No protocols, for any number of signatories and any thre art:

<u>UC</u> Non-Interactive, Proa

Rosario Gennaro<sup>§</sup>

- Only the last round of our protocols requires kr take place in a preprocessing stage, lending to a
- Our protocols withstand adaptive corruption of refresh mechanism and offer full proactive securi-
- Our protocols realize an ideal threshold signature random oracle model, assuming Strong RSA, DD a somewhat enhanced variant of existential unfor

Abstract. In a multiparty signing protocol, also known as a threshold signature scheme, the private signing key is shared amongst a set of parties and only a quorum of those parties can generate a signature. Research on multiparty signing has been growing in popularity recently due to its application to cryptocurrencies. Most work has focused on reducing the number of rounds to two, and as a result: (a) are not fully simulatable in the sense of MPC real/ideal security definitions, and/or (b) are no secure under concurrent composition, and/or (c) utilize non-standard assumptions of different types in their proofs of security. In this paper, we describe a simple three-round multiparty protocol for Schnorr signature and prove its security. The protocol is fully simulatable, secure under concurrent composition, and proven secure in the standard model of random-oracle model (depending on the instantiations of the commitmen and zero-knowledge primitives). The protocol realizes an ideal Schnor. signing functionality with perfect security in the ideal commitment and zero-knowledge hybrid model (and thus the only assumptions needed are for realizing these functionalities). We also show how to achieve proactive security and identifiable abort.

- Both protocols achieve accountability by identifying corrupted signatories in case of failure to generate a valid signature.

#### Threshold ECDSA from ECDSA Assumptions: The Multiparty Case

Jack Doerner j@ckdoerner.net Northeastern University

Evsa Lee eysa@ccs.neu.edu Northeastern University

Yashvanth Kondi ykondi@ccs.neu.edu Northeastern University abhi shelat

abhi@neu.edu Northeastern University

May 22, 2020

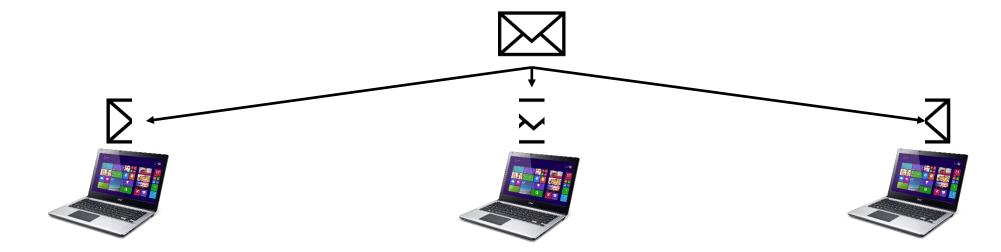
#### Abstract

Cryptocurrency applications have spurred a resurgence of interest in the computation of ECDSA signatures using threshold protocols—that is, protocols in which the signing key is secret-shared among n parties, of which any subset of size t must interact in order to compute a signature. Among the resulting works to date, that of Doerner et al. [DKLs18] requires the most natural assumptions while also achieving the best practical signing speed. It is, however, limited to the setting in which the threshold is two. We propose an extension of their scheme to arbitrary thresholds, and prove it secure against a malicious adversary corrupting up to one party less than the threshold under only the Computational Diffie-Hellman Assumption in the Global Random Oracle model, an assumption strictly weaker than those under which ECDSA is proven.



Ran Canetti<sup>†‡</sup>

- Modeling attributes of Threshold Cryptography can help identify and fine-tune the tradeoffs
- Obvious Attribute: Threshold (t-out-of-n)





## **Modeling Attributes**

• Just looking at threshold ECDSA, a survey paper [AHS20]:

Majority

Honest Dishonest Corruption

Static Adaptive **Assumptions** 

DDH Enhanced ECDSA etc.

**Building Blocks** 

VSS Commitments ZK Rounds

Setup Signing On/Offline Features

Identifiable Abort UC etc.



## **Modeling Attributes**

- These are the same kinds of attributes listed in MPC survey literature!
- As we standardize threshold cryptography, take a page from the categorization of MPC schemes
- Cryptographers generally do a good job writing down the model, can standardization help w/ communicating them?







### Implementation

Use an existing library

OR

Implement paper from scratch, possibly releasing a library

 Decision largely influenced by ease of integration and utility to application



- Non-distributed versions of these crypto primitives are already difficult to implement right
  - E.g. plain signature schemes have many flavors and optimizations, not to mention many buggy implementations
- Standardizing threshold cryptography should come with a lot more suggestions on getting it right
  - Interactive schemes are much more complex than primitives
- Viewing it as MPC can help modularize the problem



- One story: working on asynchronous MPC has helped in other areas such as Private Set Intersection (PSI)
  - Boost.Asio/libevent/Netty
- In one particular instance, work with Virginia Longitudinal Data System on a (honest-but-curious) PSI system, we applied our async knowledge:
  - 1Mx1M PSI in 5 seconds, 100Mx100M in 20 minutes
  - Despite high latency between the AUS, UK, US regions
- Hypothesis: large overlap in skillset between threshold cryptography and MPC engineering



#### "Generality"

#### **Threshold Schemes**

Custom computation

Custom messages

Faster, fewer rounds

#### **MPC Frameworks**

Off-the-shelf algo (maybe tweaked for MPC-friendliness) Run MPC

Can we find more common ground?



#### **MPC** Frameworks

- This talk is sandwiched by two threshold-from-MPC talks
- Demonstrates thresholding of many existing primitives
  - Easily extends to new ones as well
- But...
  - Might be tricky to integrate or deploy
  - Still many difficult modeling questions and choices
- Work on standardizing components, as well as good guiding principles a la NIST Privacy Framework



# **Additional Remarks**



## Standardizing MPC May Help Related Areas

- Zero-Knowledge Proofs
- Private Set Intersection
  - Apple (CSAM), Signal (Contact Discovery)

- Secure Analytics
  - Prio, Google Ads, Amazon Cleanrooms



#### Other Considerations

Access structures beyond t-out-of-n threshold

Formal Verification

• etc.



#### Info

Dr. Steve Lu, CEO <a href="mailto:steve@stealthsoftwareinc.com">steve@stealthsoftwareinc.com</a>

Dr. Rafail Ostrovsky, Co-Founder <a href="mail@stealthsoftwareinc.com">rafail@stealthsoftwareinc.com</a>

https://stealthsoftwareinc.com

#### Key Takeaways

MPC offers plenty of modeling concepts that apply to threshold cryptography

Engineering can be made more modular with MPC frameworks, can we make them easier to integrate and deploy?

Standardization effort for threshold cryptography can take advantage of the two-way street it has with MPC

