

# FHE-Related Comments on NIST First Call for Multi-Party Threshold Schemes

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### **Public Comments**

- Our initial comments were submitted in April 2023
- This talk presents updated (conceptual) comments on FHE, focusing on classes of capabilities, use cases, and security models
- Andreea Alexadru's talk presents technical details and discusses potential gadgets for Threshold FHE

#### Comments on NIST First Call for Multi-Party Threshold Schemes

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Our comments are for Fully Homomorphic Encryption (FHE) schemes based on LWE and Ring/Module LWE over power-of-two cyclotomic rings, since that is what is most commonly implemented in open-source libraries. Our comments could apply to other FHE schemes with different hardness assumptions as well (e.g., NTRU).

## Background: FHE Schemes

- Practical FHE scheme instantiations target passive security
- Common FHE schemes can be separated into three categories:
  - Brakerski-Gentry-Vaikuntanathan [BGV14] (BGV) and Brakerski [Bra12]/Fan-Vercauteren [FV12] (BFV)
    - Support SIMD encrypted computations for arithmetic circuits modulo a prime power
  - Ducas-Micciancio [DM15] (DM, also called FHEW)/Chillotti-Gama-Georgieva-Izabachene [CGGI16] (CGGI, also called TFHE)
    - Support binary or small-precision arithmetic
    - Arbitrary functions are evaluated using lookup tables via functional/programmable bootstrapping
  - Cheon-Kim-Kim-Song [CKKS17] (CKKS, also called HEAAN)
    - Support SIMD fixed-point-like arithmetic circuits (for many real-number applications)
- All these schemes are based on LWE/RLWE/MLWE
  - There are also less-common variants based on NTRU
  - Thresholdization of FHE schemes in practice follows the same blueprint with key-homomorphism used for distributed key generation and masked partial decryptions for distributed decryption

### **Background: Current Standardization Efforts**

- <u>HomomorphicEncryption.org</u> open consortium of industry, government and academia to standardize homomorphic encryption
  - Founded in 2017
  - 6 meetings held from 2017 to present
  - Security recommendations are available since 2018
- ISO standardization (in progress)
  - Targeting BGV, BFV, CKKS, and CGGI
  - Technical Committee : ISO/IEC JTC 1/SC 27 Information security, cybersecurity and privacy protection

### Background: Main Active Open-Source Software Libraries

#### OpenFHE

- Implements in C++ single-key FHE for all five schemes
- Supports threshold FHE for BGV, BFV, and CKKS

#### Lattigo

• Implements in Go single-key and threshold FHE for BGV, BFV, and CKKS

#### TFHE-rs

Implements in rust the CGGI scheme (with many enhancements)

### **Classes of HE-Related Capabilities**

- We suggest separating lattice-based HE capabilities into three classes
- Class 1: Threshold key generation and decryption
- Class 2: HE schemes with linearly-homomorphic operations/additive homomorphic encryption (AHE)
- Class 3: FHE schemes that support nonlinear operations, which may use relinearization, bootstrapping, and/or other similar techniques
  - This class should probably be further broken down into subclasses
  - Circular security should be considered in the context of Class 3
- Targeting the first two classes should be easier, both under passive and active security

### Use Cases for HE

#### Class 1: Threshold key generation and decryption

- Same as for Category 1 of the NIST MPTS call
- Can build upon the NIST PQC standardization effort + thresholdization [CCMS21]

#### Class 2: Additive homomorphic encryption

- Often used as a gadget in hybrid cryptosystems, e.g., MP-SPDZ uses BGV for addition and plaintext-ciphertext computations to generate Beaver triplets
- Several other use cases are discussed in our public comments, e.g., secure voting, aggregation, some PIR schemes
- Class 3: FHE schemes
  - Certain PIR and PSI schemes, AES tranciphering, neural network & training
  - See our public comments for more details

#### **Passive Security**

- IND-CPA security is typically sufficient to achieve passive security (for data privacy) for exact FHE schemes, including BGV, BFV, DM, and CGGI
- IND-CPA security is not sufficient for **approximate** FHE schemes
  - Li and Micciancio showed that CKKS is not secure if access to a decryption oracle is provided, i.e., when the decryption result is shared with parties that do not have the secret key [LM21]
  - They proposed a new definition IND-CPA<sup>D</sup> that provides access to encryption, evaluation, and decryption oracles
- Threshold FHE schemes also require access to decryption oracles (for partial decryptions) and similar solutions, i.e., approximate encryption can be viewed as a special case of threshold FHE [KS23]

### **Active Security**

- What are the potential approaches for building actively-secure threshold encryption schemes based on lattices?
- Class 1: Fujisaki-Okamoto (FO) transformation (very challenging for scenarios with HE)?
  - Two of the three lattice-based finalists in the NIST PQC competition, Crystals-Kyber and Saber, built an IND-CPA secure encryption scheme and then applied the FO transform to create an IND-CCA hybrid scheme
  - There are some challenges in thresholdizing this approach; however, initial constructions for thresholdized Saber based on this approach are available [CCMS21]
- Class 2: Actively secure BGV-based threshold encryption/AHE
  - Overdrive [KPR18] employs BGV for addition and plaintext-ciphertext multiplication to generate Beaver triples, applying noise flooding for circuit privacy and zero-knowledge proofs for active security
  - Aranha et al. use actively secure threshold BGV (both threshold key generation and threshold decryption) and ZKPs to construct an efficient form of secure voting [ABGS22]

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# Thank You! Have questions?

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