



AONT: an essential gadget for Multi-Party Threshold Cryptography.

presented by Gilles Seghaier, Cofounder & CTPO of Astran
during Session 3c (13:20 EDT) Thursday, September 28th, 2023

MPTS 2023: NIST Workshop on Multi-Party Threshold Schemes 2023



THANK YOU

NLST

Astran Academic Partners



Astran Scientific Committee



Ludovic
Perret



Nigel
Smart

Professor Nigel Smart - Cryptographer and professor of computer science at the University of Leuven in Belgium, renowned for his work on elliptic curve cryptography and matching-based cryptography. He co-founded Unbound Security, a company specializing in the deployment of distributed cryptographic solutions based on multiparty computation (MPC).

Professor Ludovic Perret - Cryptography expert and lecturer at Sorbonne University, specializing in the standardization of post-quantum cryptography. Co-author of the GeMSS digital signature scheme and involved in several standardization bodies.

Zero Trust & Zero Knowledge Cloud Services

aws

S3

Simple
Storage
Service

THRESHOLD
CRYPTOGRAPHY



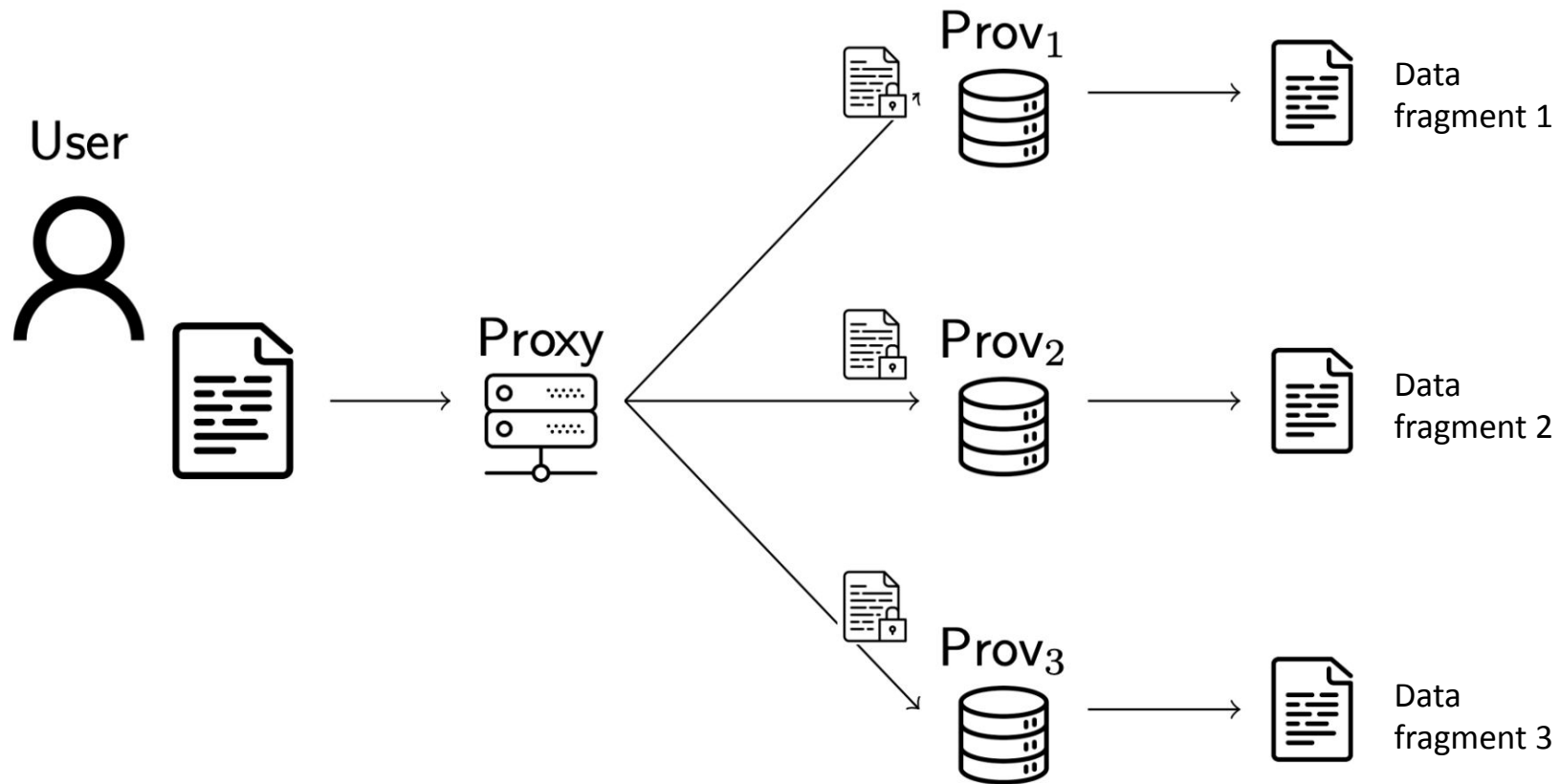
MULTI-PARTY
STORAGE

ASTRAN

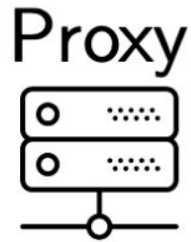
S5

Secret
Shared
S3

Multi-cloud storage with a proxy



MPS and threshold modeling possibilities



A single proxy

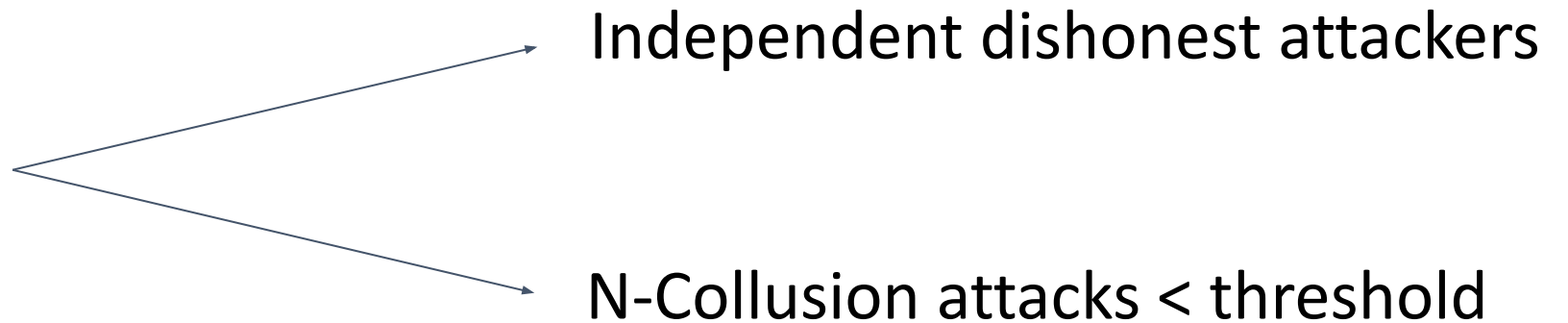


n storage providers



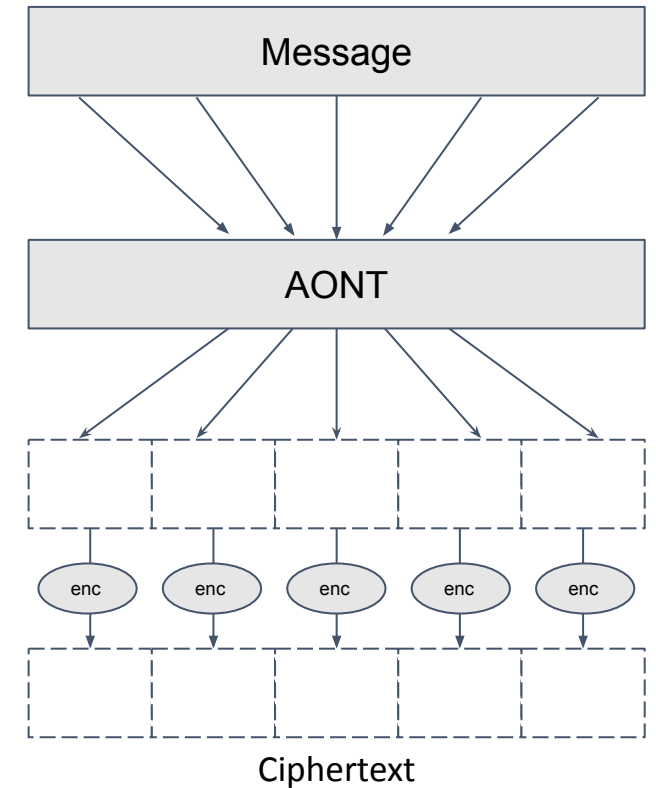
k out of n fragments threshold

n+1 adversaries



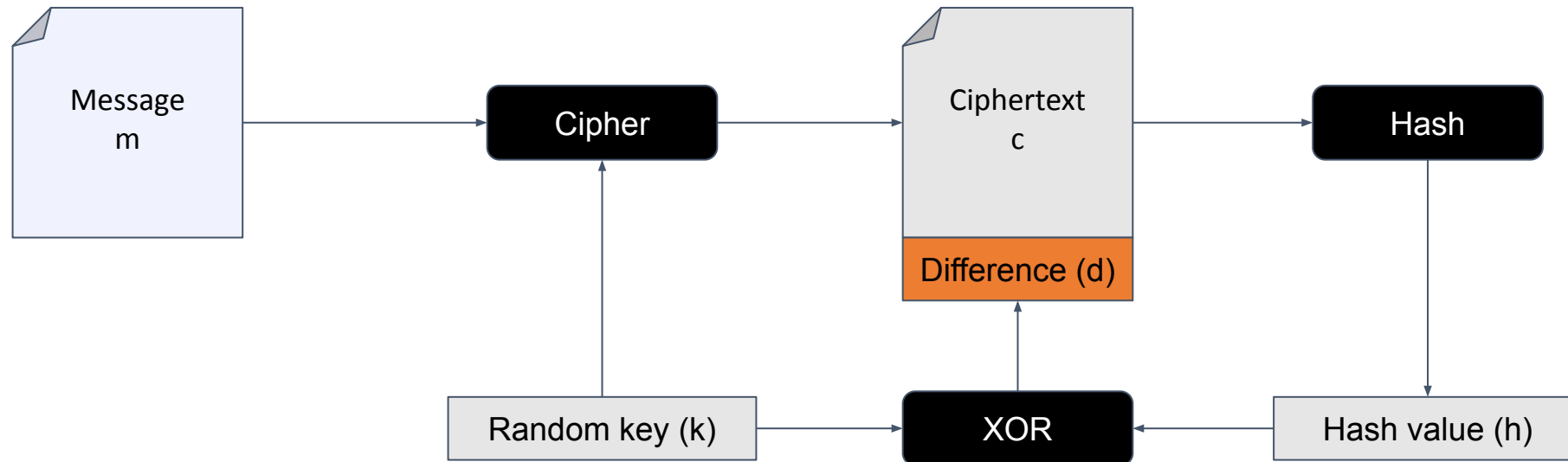
AONT

- All-Or-Nothing transform was introduced by Rivest[1] back in 1997.
- One must decrypt the entire ciphertext before one can determine even one message block.
- Original motivation: slowing down brute-force searches against all-or-nothing encryption blocks.



$$\mathbf{AONT(m) = Enc(m,k) \parallel XOR(k,h)}$$

where $h = \text{HASH}(\text{Enc}(m,k))$



Can be used in conjunction with error coding, secret sharing, IDA or others threshold schemes

<p style="text-align: center;">Secret Sharing Scheme (SSS)</p>	<p>We use Shamir's secret sharing [2]. It exploits the Lagrange interpolation theorem, specifically that k values suffice to uniquely determine a polynomial of degree $\leq k - 1$. Shamir's secret sharing has perfect secrecy.</p>
<p style="text-align: center;">Information Dispersal Algorithm (IDA)</p>	<p>Unlike secret sharing, an IDA does not provide perfect-secrecy. However, an IDA is very memory-efficient. We are using algorithms similar to Rabin's IDA using erasure codes.</p>
<p style="text-align: center;">Proxy Re-Encryption Scheme (PRE)</p>	<p>We use in our protocols the fully homomorphic encryption scheme BGV [3]. Since BGV is fully homomorphic, it commutes with the secret sharing described above. It can perform proxy re-encryption by using the key-switching method described in [3].</p>
<p style="text-align: center;">Multikey Encryption Scheme (MKE)</p>	<p>We use multi-key homomorphic encryption [4] scheme. That way, the providers can also each decrypt their own share, this time with the participation of the others.</p>

[2] A. Shamir. How to share a secret. Commun. ACM, 1979.

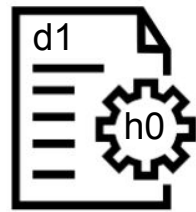
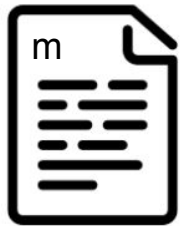
[3] Z. Brakerski, C. Gentry, and V. Vaikuntanathan. (leveled) fully homomorphic encryption without bootstrapping. In S. Goldwasser, editor, Innovations in Theoretical Computer Science 2012. ACM, 2012

[4] A. López-Alt, E. Tromer, and V. Vaikuntanathan. On-the-fly multiparty computation on the cloud via multikey fully homomorphic encryption. IACR Cryptol. ePrint Arch., 2013.

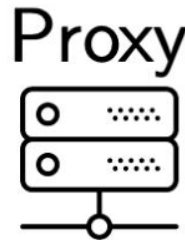
User



$$(pk_U, sk_U) \leftarrow \text{PRE.UKeyGen}(\lambda)$$

$$r_{ki} \leftarrow \text{PRE.ReKey}(sk_U, p_{ki})$$


$$h_0, d_1$$

$$r_{k1}, \dots, r_{kn}$$


$$r_i, h_i$$


$$d_0 || d_1 \leftarrow \text{AONT.Hide}(m)$$

$$h_0 \leftarrow \text{PRE.Enc}(d_0, pk_U)$$

$$s_1, \dots, s_n \leftarrow \text{SS.Split}(h_0, n, k)$$

$$h_i \leftarrow \text{PRE.ReEnc}(s_i, r_{ki})$$

$$r_1, \dots, r_n \leftarrow \text{IDA.Split}(d_1, n, k)$$

$$y_i \leftarrow \text{PRE.Dec}(h_i, ski)$$

$$\text{store } y_i, r_i$$


Transformation
+ Homomorphic Encryption



Fragmentation
+ Re-encryption



Dispersal

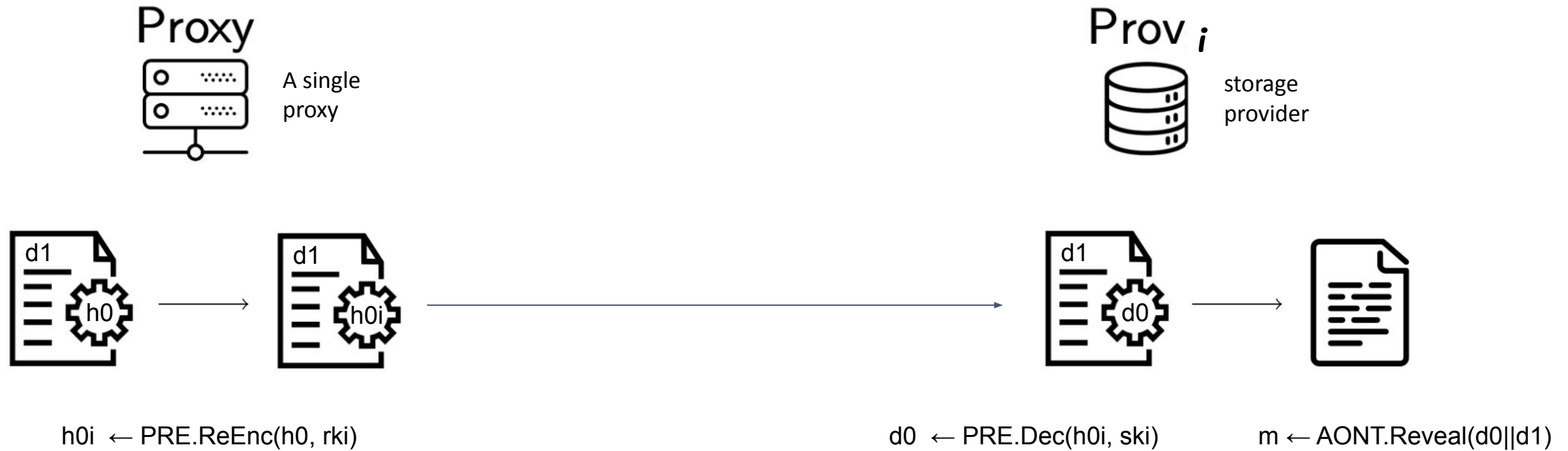


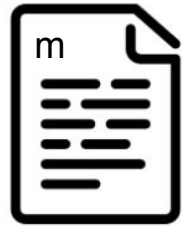
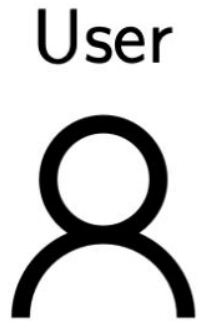
Decryption



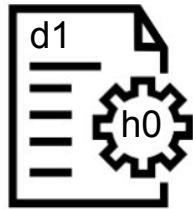
Storage

2-collusion : dishonest proxy and provider

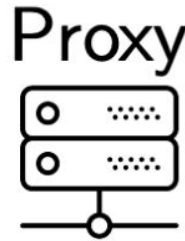




$d0||d1 \leftarrow \text{AONT.Hide}(m)$
 $h0 \leftarrow \text{MKE.Enc}(d0, \{pki\})$



$h0, d1$



$h1, \dots, hn \leftarrow \text{SS.Split}(h0, n, k)$
 $r1, \dots, rn \leftarrow \text{IDA.Split}(d1, n, k)$

hi



$(pki, ski) \leftarrow \text{MKE.KeyGen}(\lambda)$

$\{hij\} (i \neq j)$

$\{hij\} \leftarrow \text{MKE.PartDec}(hi, ski)$

$ri, \{hij\} (j \neq i)$



$yi \leftarrow \text{MKE.FinDec}(hij)$

store yi, ri

Transformation
+ Homomorphic Encryption

Fragmentation

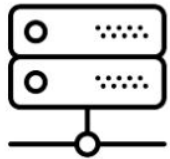
Dispersal

Partial Decryption
+ Decryption

Storage

k-collusion : dishonest proxy and k-1 providers

Proxy



A single
proxy

Prov i



n storage
providers

Frag



k out of n
fragments
threshold

All-or-nothing Transform (AONT)

n out of n is a threshold!

Brings strong secrecy and integrity guarantees

Memory-efficient for large volumes of data compared to SS

An essential building block in our current Multiparty Storage use case

A great and flexible gadget when combined with other schemes and algorithms

A generic scheme that can be implemented in many ways to provide additional threshold capabilities

Interested in our tech?

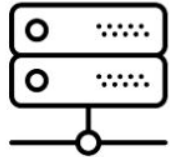
hello@astran.io

The background is a solid blue color with several faint, semi-transparent geometric shapes. There is a hexagon on the left side, a pentagon at the top center, and a large, rounded shape on the right side. The text is centered in the middle of the image.

**Ask Me (Almost)
Anything**

Appendices

Proxy



A single proxy

Prov i



n storage providers

Frag



k out of n fragments threshold

Provider-Secrecy	The data's confidentiality is preserved against cloud storage providers individually	The adversary plays the roles of a provider alone.
Proxy-Secrecy	The data's confidentiality is preserved against cloud proxy individually	The adversary plays the role of the proxy alone..
Provider-Collusion-Secrecy	The data's confidentiality is preserved against cloud storage providers collusion to a given threshold	The adversary plays the role of k colluding providers. k -provider-secrecy assumes the proxy is a trusted party.
Proxy-Provider-Collusion-Secrecy	The data's confidentiality is preserved against the proxy colluding with a given number of cloud storage provider	The adversary plays the role of the proxy and k colluding providers.