

Towards Standardization of Threshold Schemes at NIST

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National Institute of Standards and Technology
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Presentation at the
Theory of Implementation Security (TIS'19) Workshop
November 11, 2019 @ London, UK

Some slides are based on previous presentations ([NTCW'19](#); [ICMC'19](#); [ACS'19](#)).

The NIST Threshold Cryptography project, on which this presentation is based, has so far also counted with the participation of Apostol Vassilev, Michael Davidson, Nicky Mouha.

Outline

1. Crypto standards at NIST
2. Threshold intro
3. Threshold project
4. Threshold preliminary roadmap
5. Concluding remarks

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Goals of this presentation:

- ▶ Overview of the NIST standardization effort
- ▶ Present the new “preliminary roadmap” (NISTIR 8214A)
- ▶ Encourage feedback and collaboration

Outline 1

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Some NIST data

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(National Bureau of Standards 1901–1988 → NIST 1988–present)

- ▶ Non-regulatory federal agency (within the U.S. Department of Commerce)
- ▶ **Mission** (keywords): innovation, industrial competitiveness, measurement science, standards and technology, economic security, quality of life.



Aerial photo of Gaithersburg campus (source: Google Maps, August 2019)

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Wide spectrum of competences

- $\sim 6-7 \times 10^3$ workers
- Five laboratories and two centers
- Laboratories → Divisions → Groups → Projects
- Standards, research and applications



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Laboratories, divisions, groups

Information Technology Laboratory (ITL):

advancing measurement science, standards, and technology through research and development in information technology, mathematics, and statistics.



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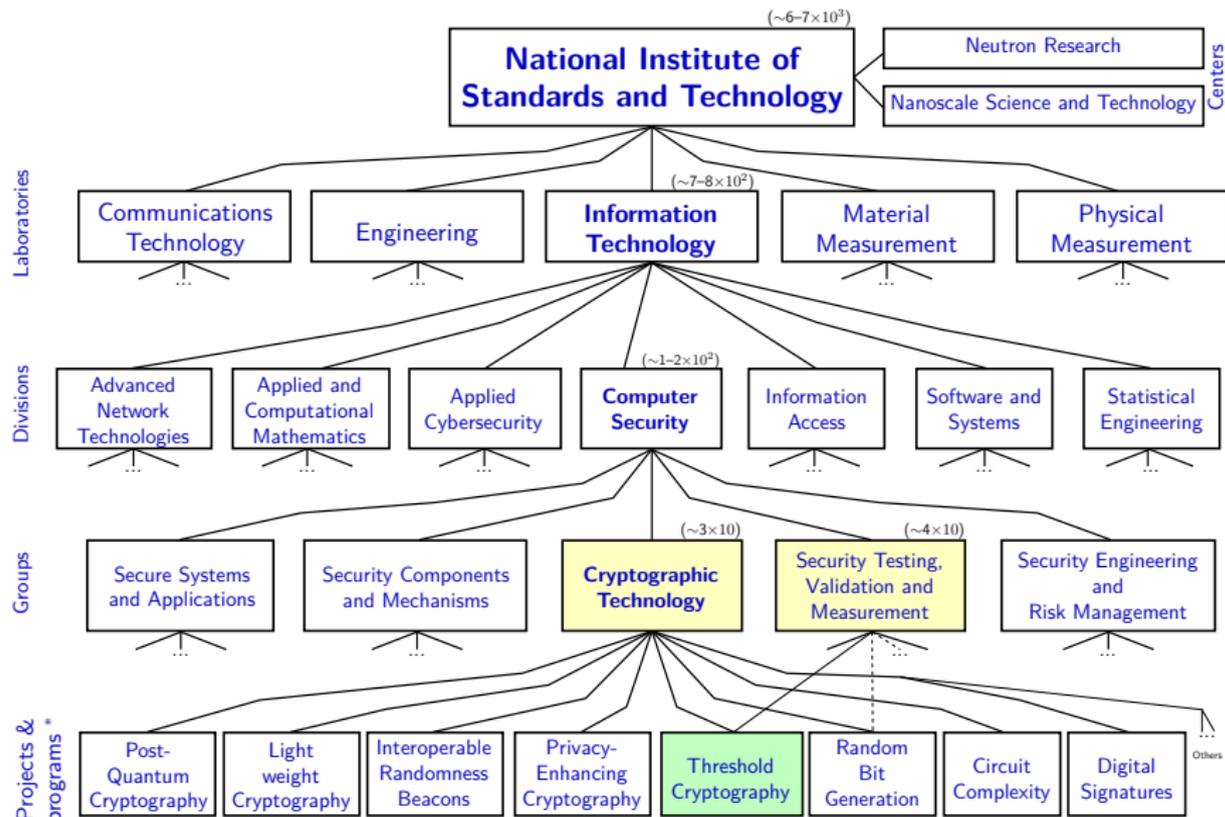
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▶ Documents: FIPS, SP 800, NISTIR.

▶ International cooperation: government, industry, academia, standardization bodies.

FIPS = Federal Information Processing Standards; SP 800 = Special Publications in Computer Security; NISTIR = NIST Internal or Interagency Report.

Some projects of crypto primitives/applications at NIST



* (Some projects/programs involve several groups, divisions or labs)

(in parenthesis: approximate range # workers, inc. associates and fed. employees)

Some standardized cryptographic primitives

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- ▶ Block ciphers
- ▶ Cipher modes of operation
- ▶ Hash functions
- ▶ Signatures
- ▶ Pair-wise key agreement
- ▶ DRBGs

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(Further details in “NIST Cryptographic Standards and Guidelines Development Program Briefing Book”)

Legend:

- AES = Advanced Encryption Standard
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Several methods:

- ▶ Internal or interagency developed techniques
- ▶ Adoption of external standards
- ▶ Open call, competition, “competition-like”

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- ▶ **Pairing-based Cryptography**: workshop (2008), study and call for feedback on use cases (2011), report (2012–2015) (forming NIST's position on standardization/recommendation: more research is needed).

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Development process:

- ▶ NISTIR 7977: NIST Cryptographic Standards and Guidelines Development Process (2016). Formalizes several **principles** to follow:
 - ▶ transparency
 - ▶ openness
 - ▶ balance
 - (and overarching considerations)
 - ▶ integrity
 - ▶ technical merit
 - ▶ usability
 - ▶ global acceptability
 - ▶ continuous improvement
 - ▶ innovation and intellectual property

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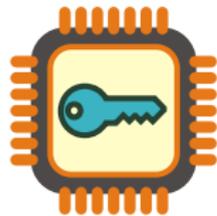
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Specially relevant: key-based cryptographic primitives



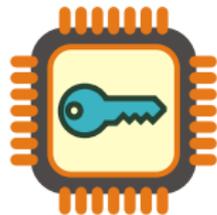
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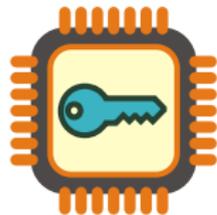
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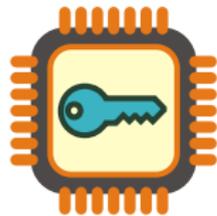
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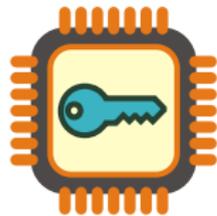
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Some things can go wrong!

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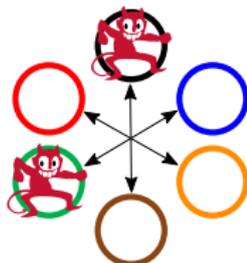
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The threshold approach

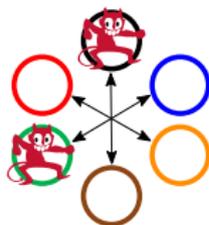


The red dancing devil is from
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The threshold approach

At a high-level:

use redundancy & diversity to mitigate the *compromise* of up to a threshold number (f -out-of- n) of components

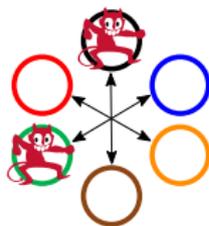


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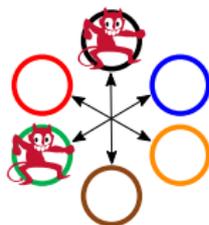
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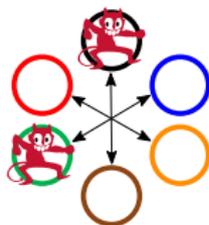
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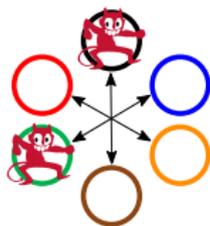
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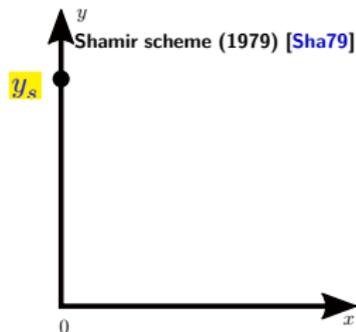
- ▶ **withstands** several *compromised* components;
- ▶ **needs** several un*compromised* components;
- ▶ **prevents** secret keys from being in one place;
- ▶ **enhances** resistance against side-channel attacks; ...

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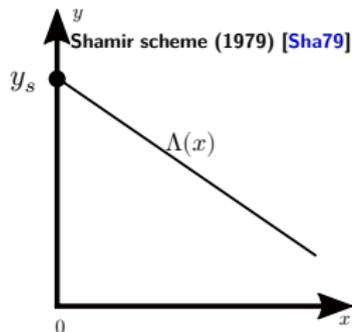


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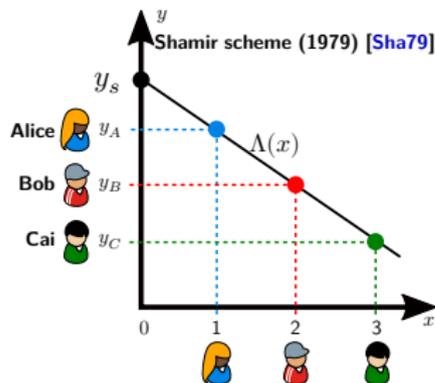


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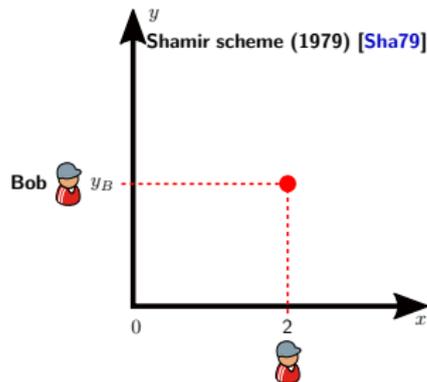


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- ▶ A random line Λ is drawn crossing the secret;
- ▶ Each share is a point $(\Lambda(i), i)$ in the line Λ ;

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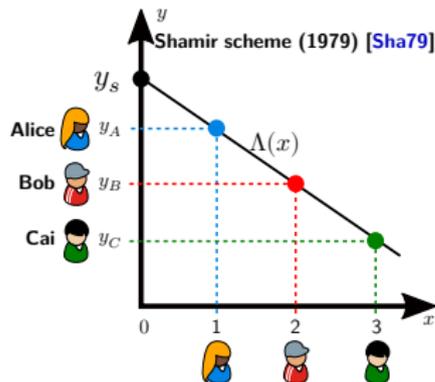
Example 2-out-of- n secret sharing

- ▶ The secret y_s is placed in the y -axis;
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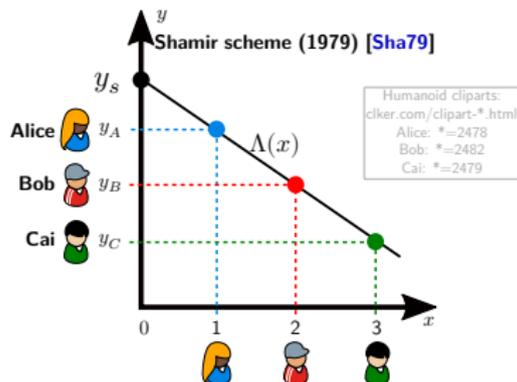
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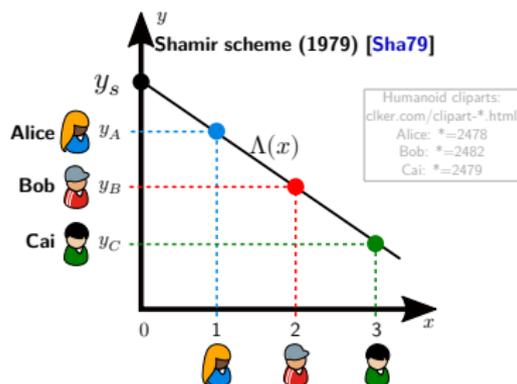
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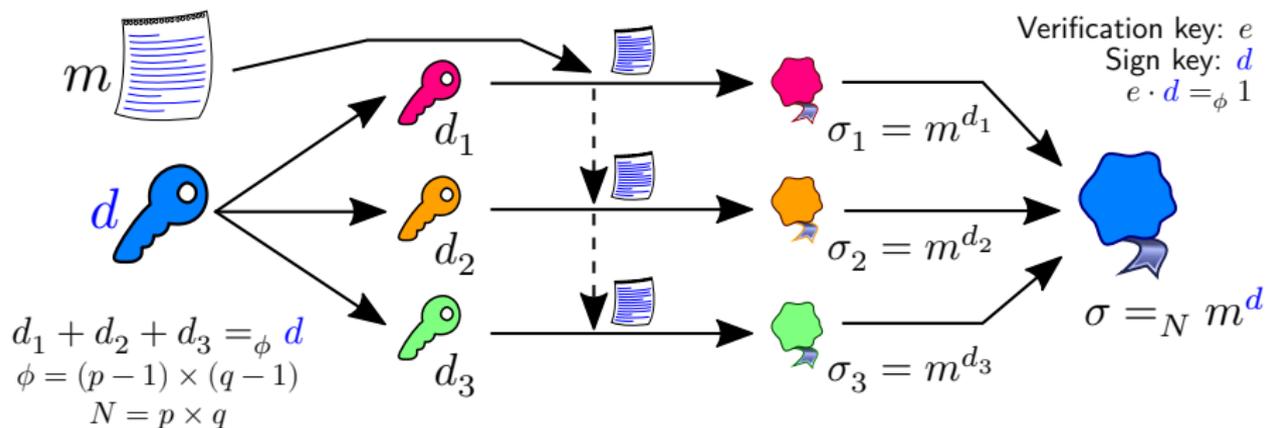
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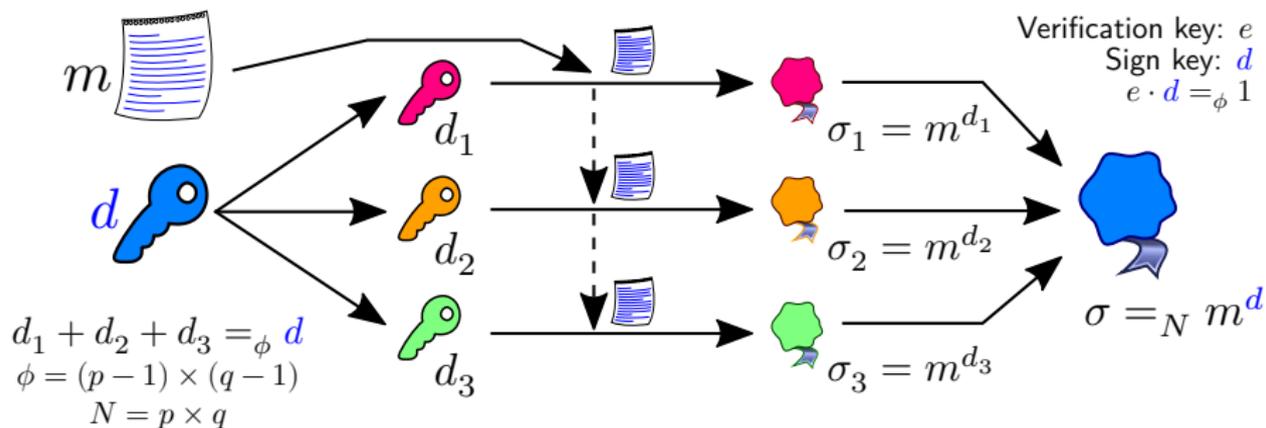
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Use [threshold schemes for cryptographic primitives](#) (next)

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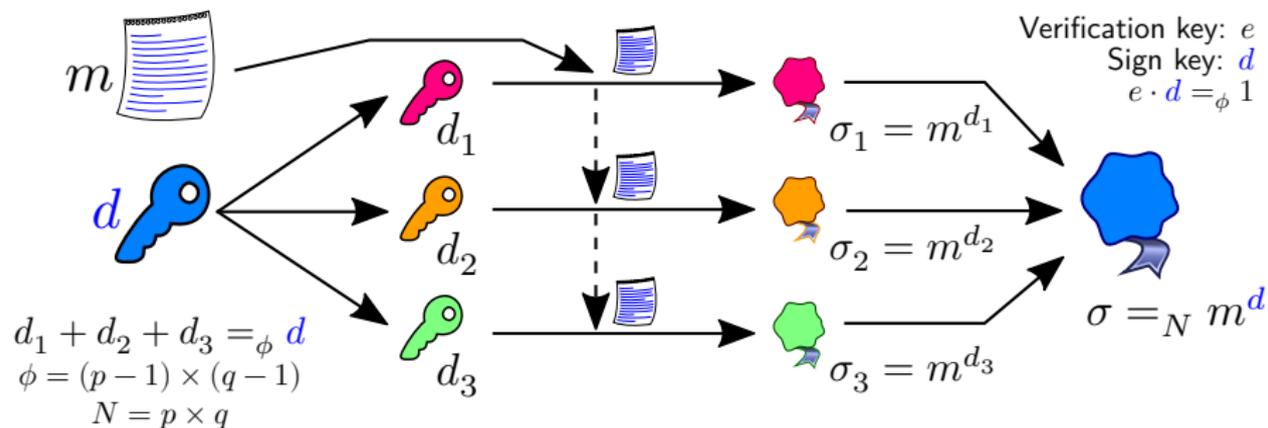
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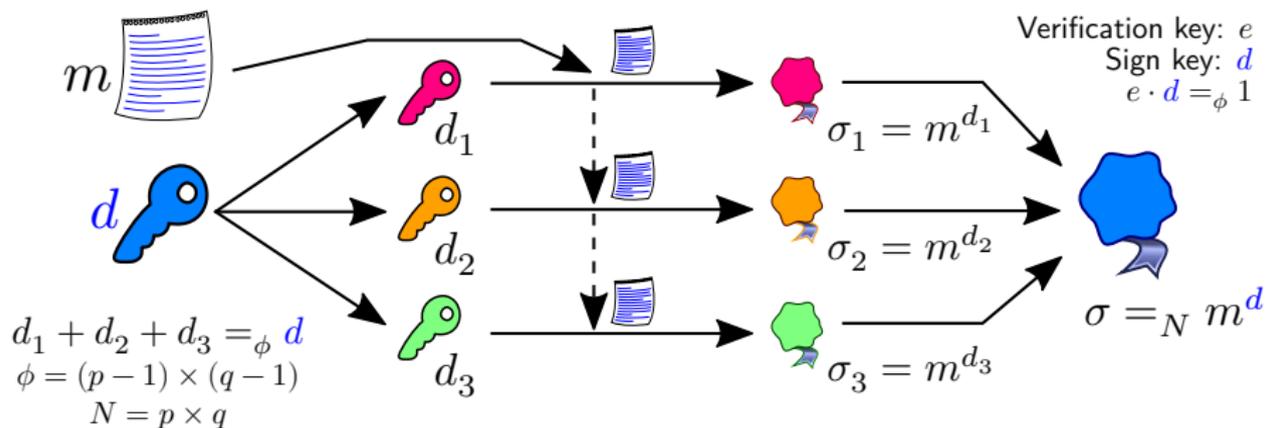
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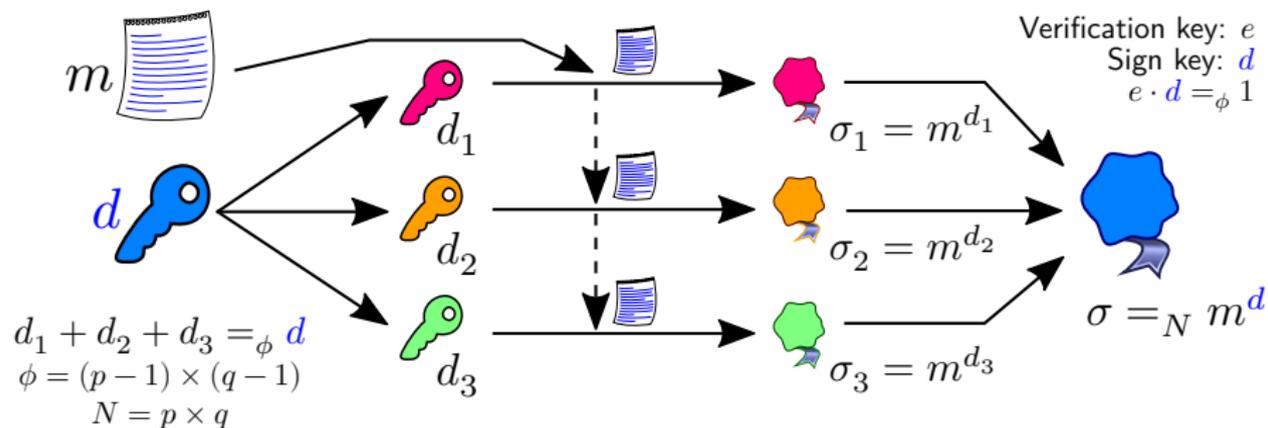
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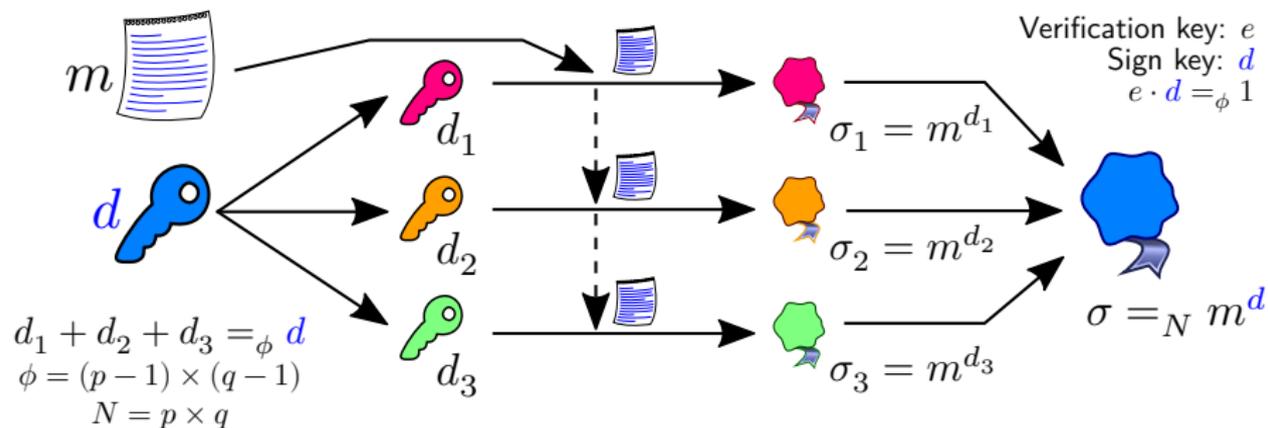
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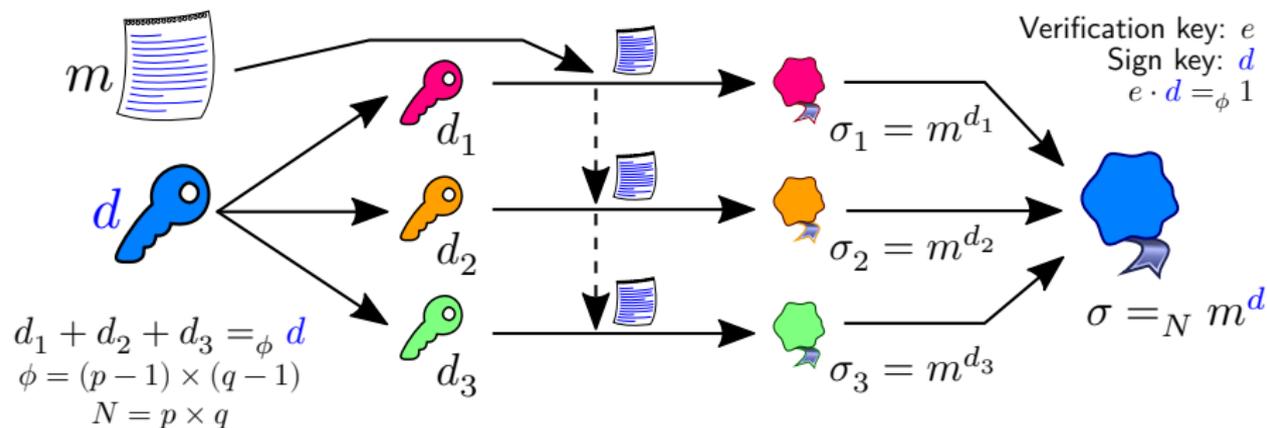


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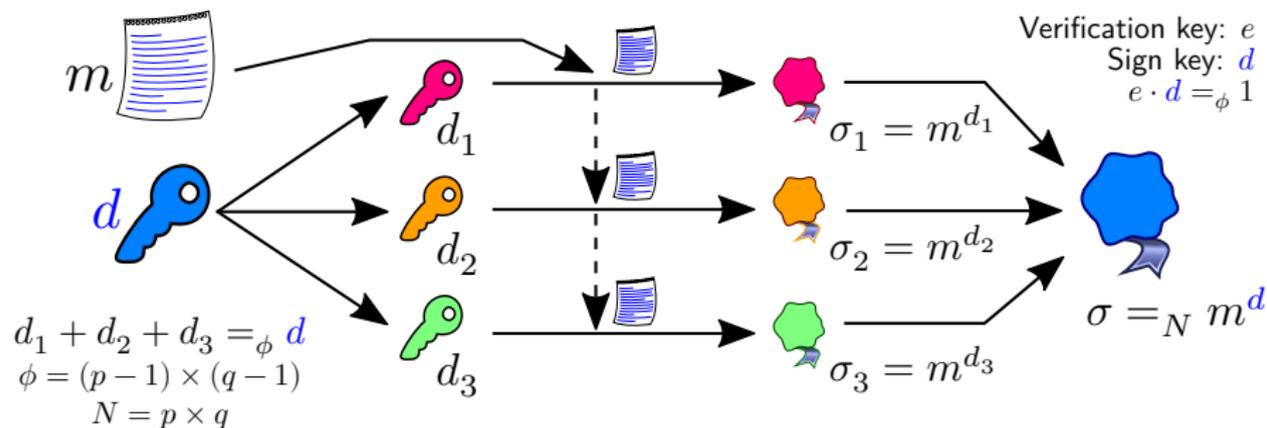


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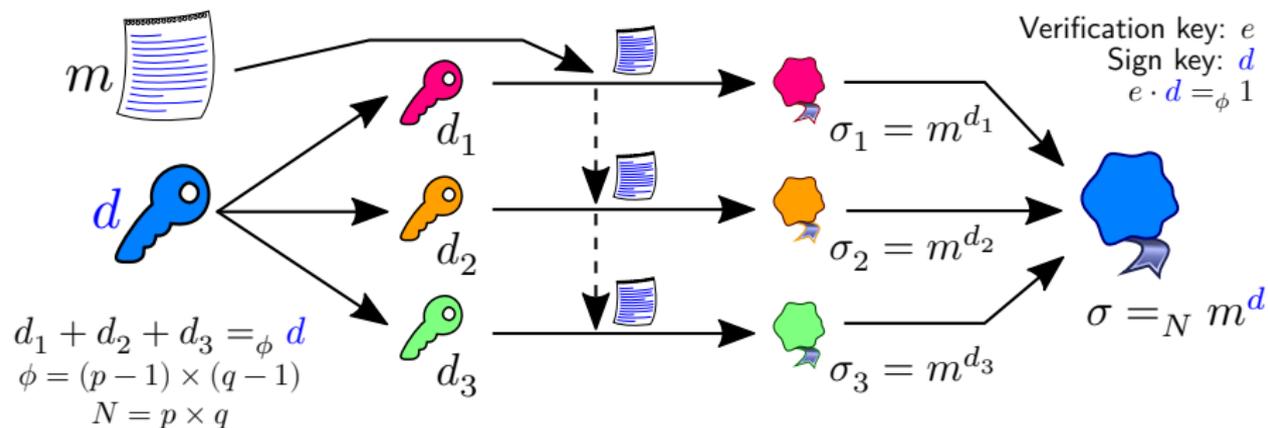


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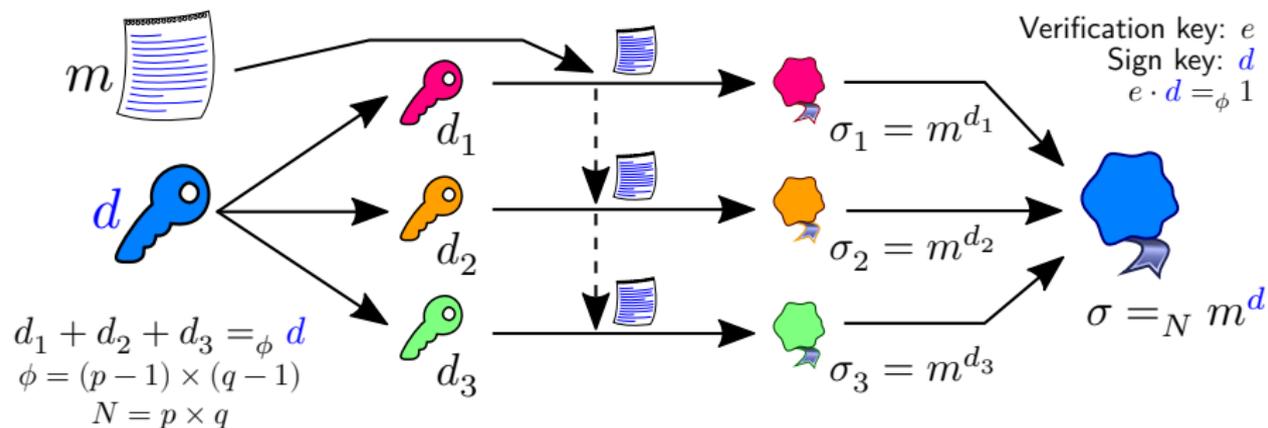
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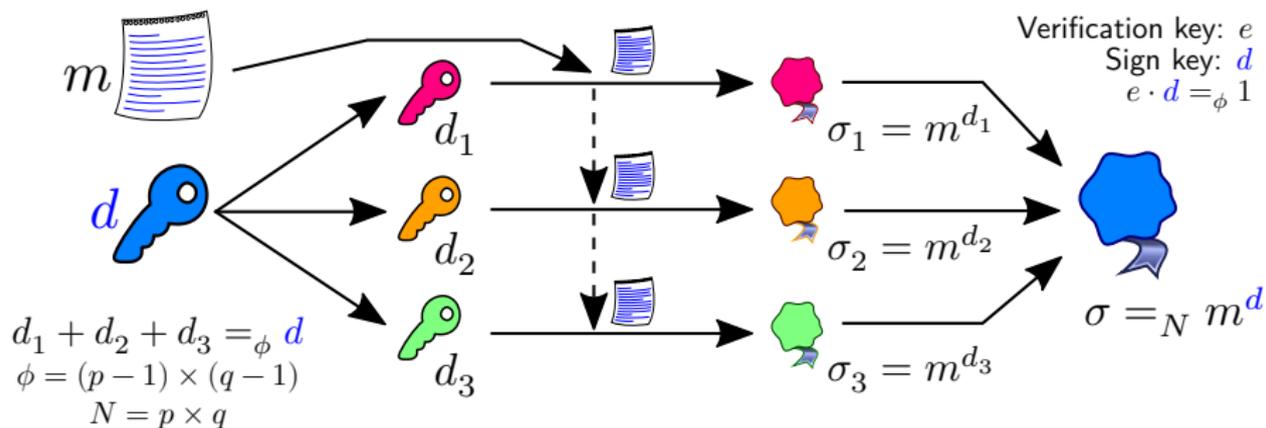
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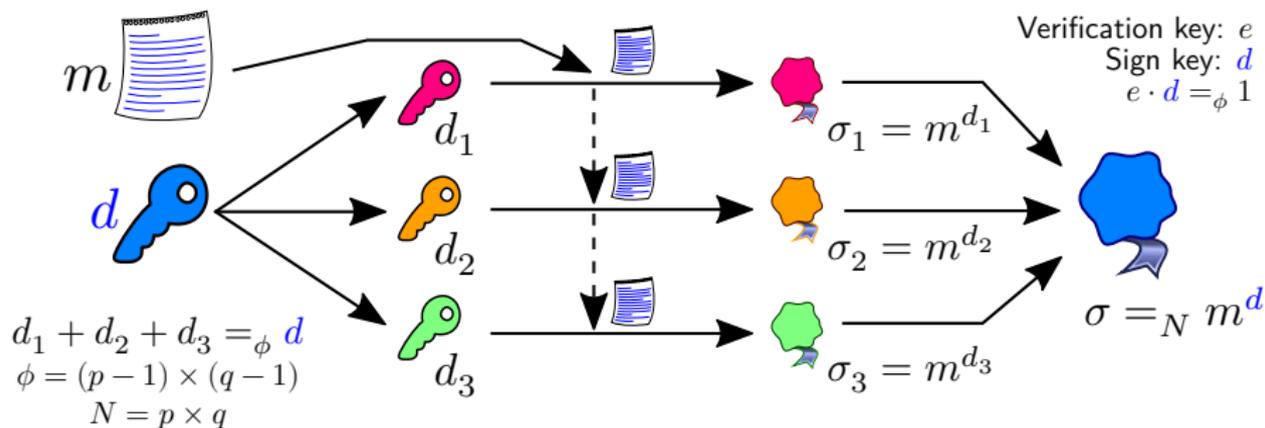
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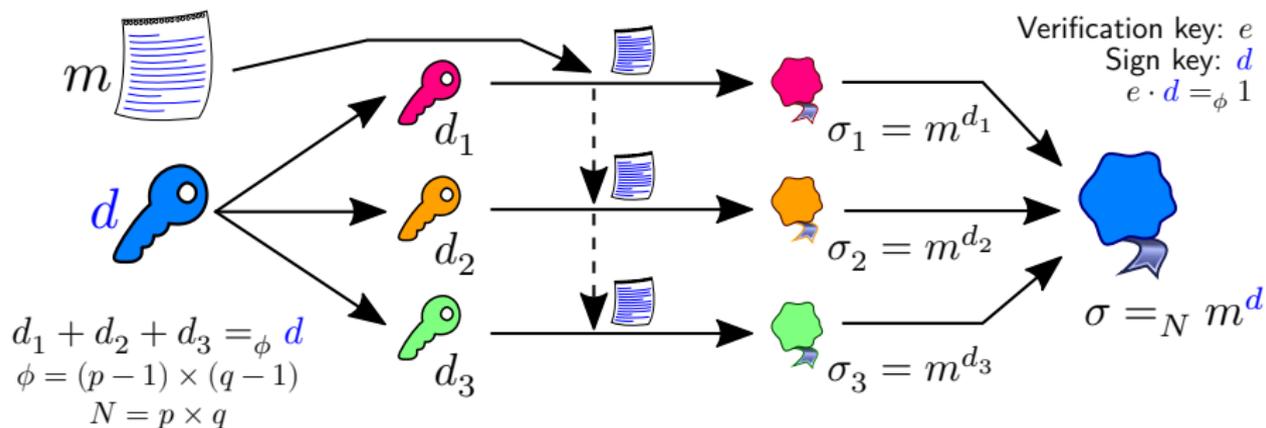
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Depends on attack model (e.g., attack surface, ...), system model (e.g., rejuvenations, ...), ...

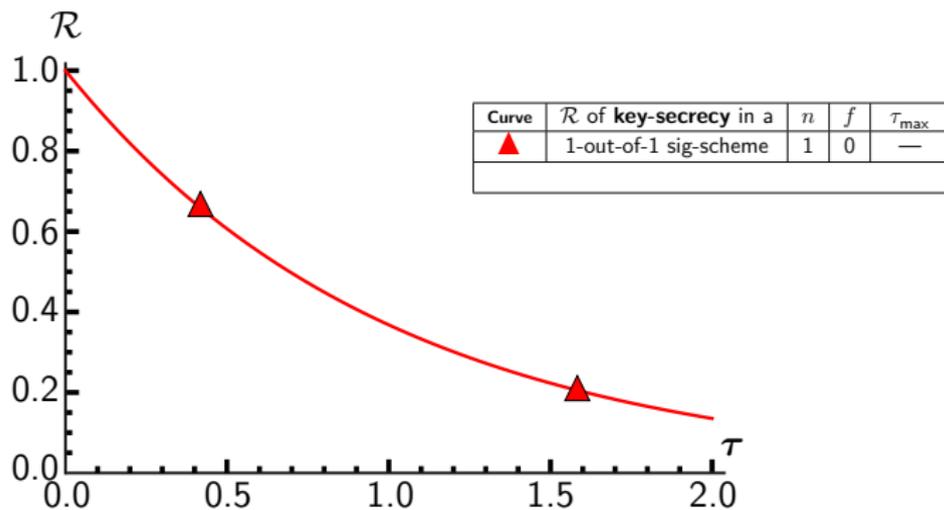
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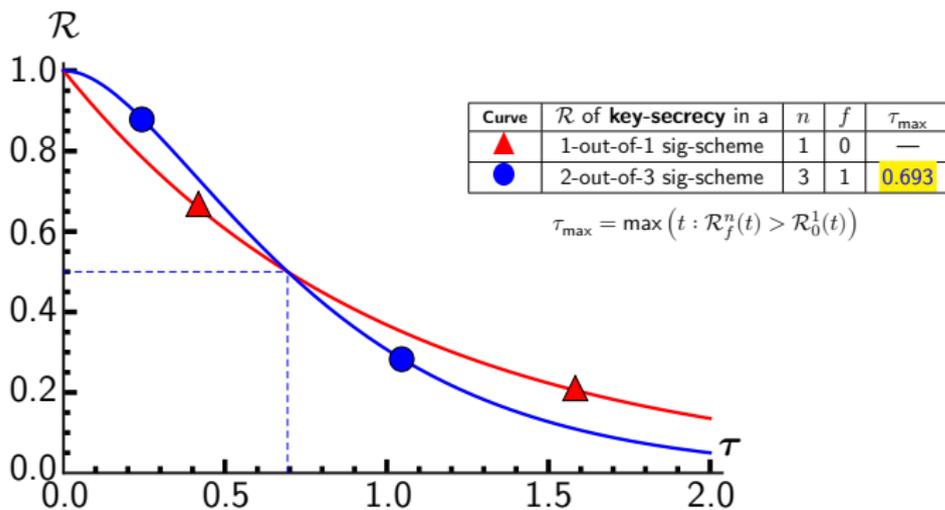


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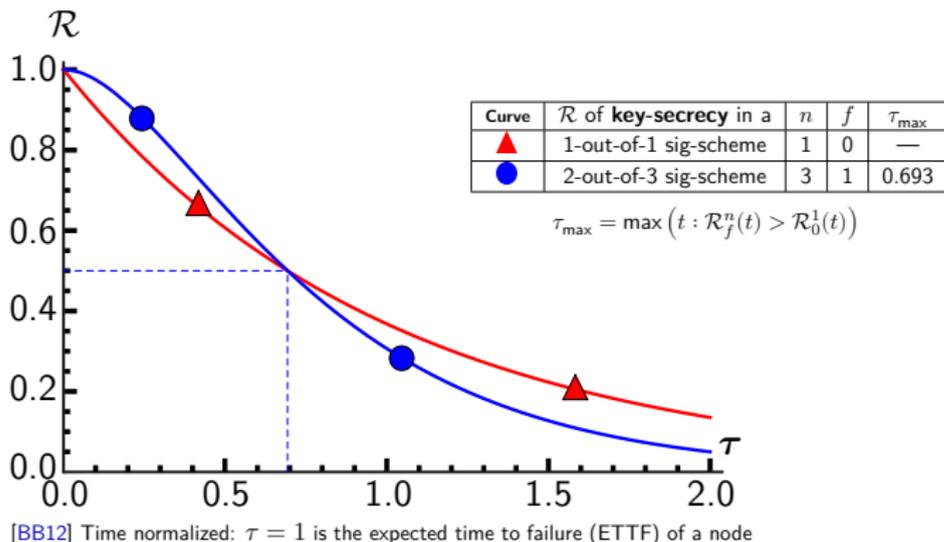
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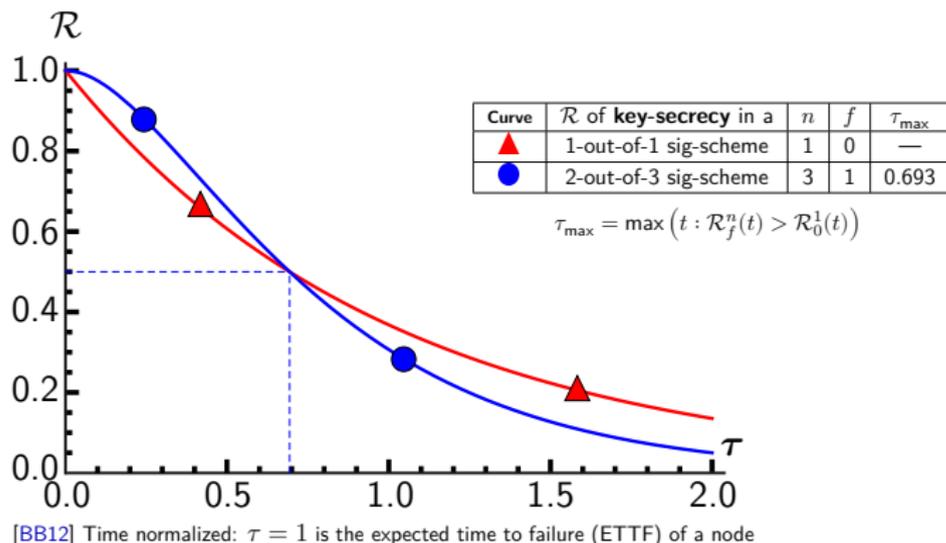


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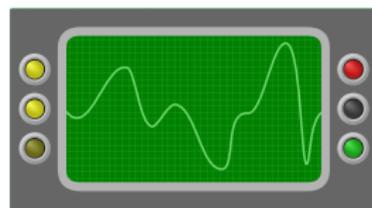
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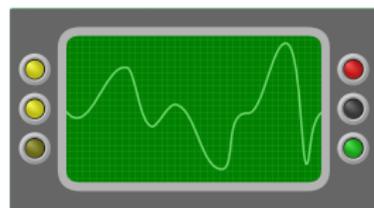
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Challenge questions:

- ▶ which models are realistic / match state-of-the-art attacks?
- ▶ what concrete parameters (e.g., n) thwart real attacks?

Outline 3

1. Crypto standards at NIST
2. Threshold intro
3. Threshold project
4. Threshold preliminary roadmap
5. Concluding remarks

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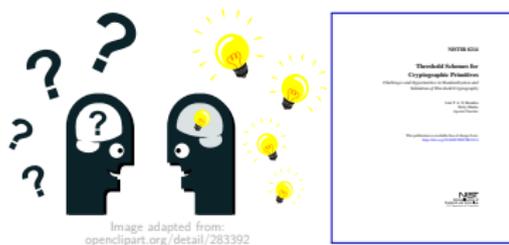
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Past timeline:

- ▶ 2018-July: Draft online 3 months for public comments
- ▶ 2018-October: Received comments from 13 external sources
- ▶ 2019-March: Final version online, along with “diff” and received comments

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- Kinds of threshold 

- Communication interfaces 

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- Setup and maintenance  

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But there are other factors ...

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▶ **Conceivable attack types.**



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- ▶ Active vs. passive
- ▶ Invasive (physical) vs. non-invasive
- ▶ Static vs. adaptive
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- ▶ Parallel vs. sequential (wrt attacking nodes)

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A threshold scheme **improving** security against an attack in an application **may be powerless or degrade** security for another attack in another application

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Validation is needed in the federal context:

- ▶ need to use **validated** implementations [tC96] of **standardized** algorithms
- ▶ FIPS 140-2/3 defines, for cryptographic modules, 4 security levels:
subsets of applicable security assertions [NIS01, NIS19]

(FIPS = Federal Information Processing Standards)

#NTCW2019

NIST Threshold Cryptography Workshop 2019

<https://csrc.nist.gov/Events/2019/NTCW19>

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www.nist.gov/image/surfgaithersburg.jpg

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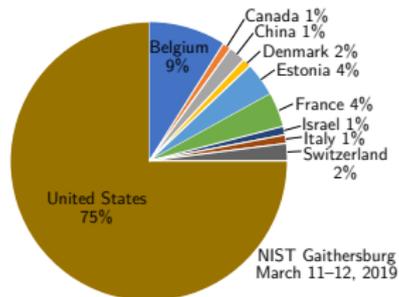
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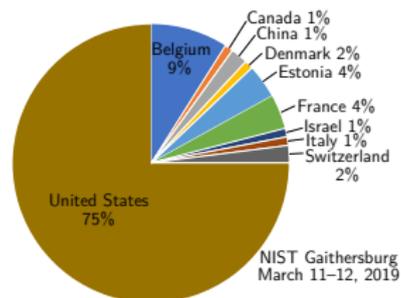
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A platform for open interaction:

- ▶ hear about experiences with threshold crypto;
- ▶ get to know stakeholders;
- ▶ get input to reflect on roadmap and criteria.

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Discussion of diverse topics:

- ▶ threshold schemes in general (motivation and implementation feasibility);
- ▶ NIST standardization of cryptographic primitives
- ▶ a post-quantum threshold public-key encryption scheme;
- ▶ threshold signatures (adaptive security; elliptic curve digital signature algorithm);
- ▶ validation of cryptographic implementations;
- ▶ threshold circuit design (tradeoffs, pitfalls, combined attacks, verification tools);
- ▶ secret-sharing with leakage resilience;
- ▶ distributed symmetric-key encryption;
- ▶ applications and experience with threshold cryptography.

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Some notes:

- ▶ differences in granularity (building blocks vs. full functionalities);
- ▶ separation of single-device vs. multi-party;
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- ▶ usefulness of explaining rationale (e.g., as complimented for the NISTIR);
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Results

A step in *driving an open and transparent process towards standardization of threshold schemes for cryptographic primitives.* (See [NISTIR 7977](#))

Some notes:

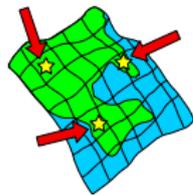
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These elements are helpful for the next step ... designing a roadmap

Outline 4

1. Crypto standards at NIST
2. Threshold intro
3. Threshold project
- 4. Threshold preliminary roadmap**
5. Concluding remarks

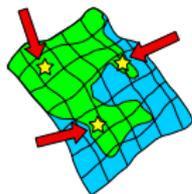
A “preliminary” roadmap



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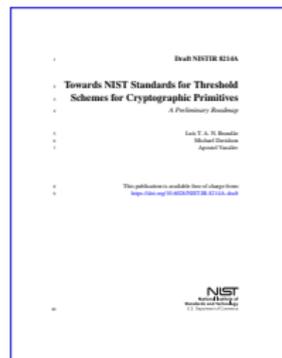
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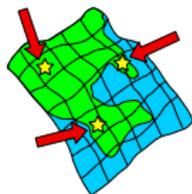
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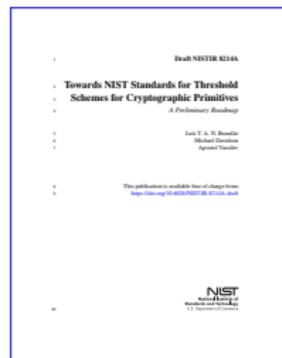
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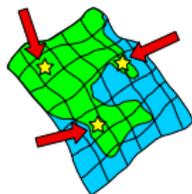
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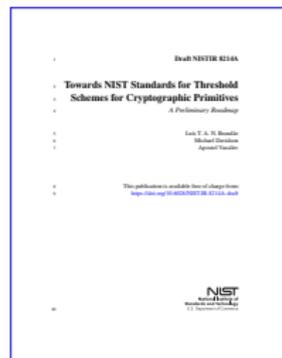
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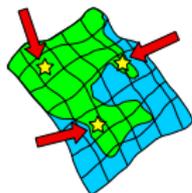
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Lays the basis towards a roadmap:

- ▶ Map/organize potential items for standardization
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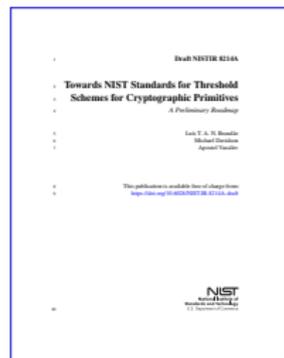
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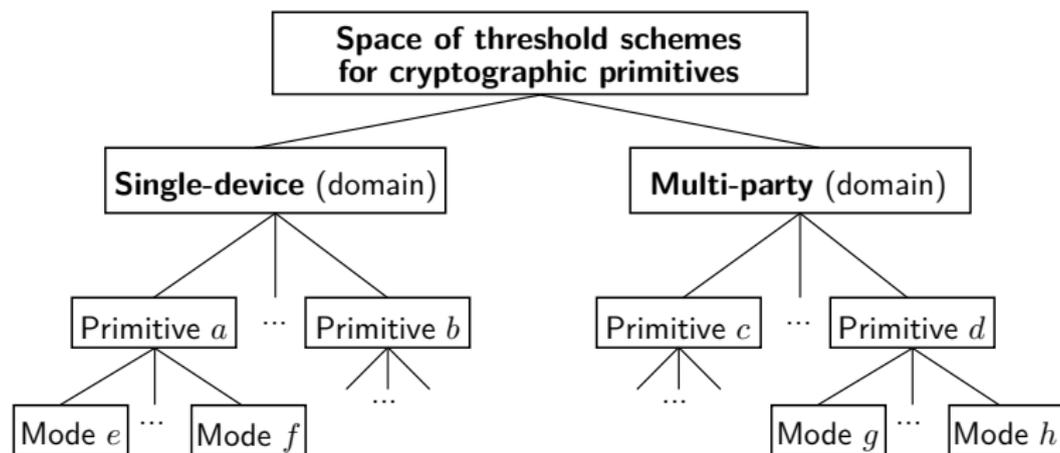
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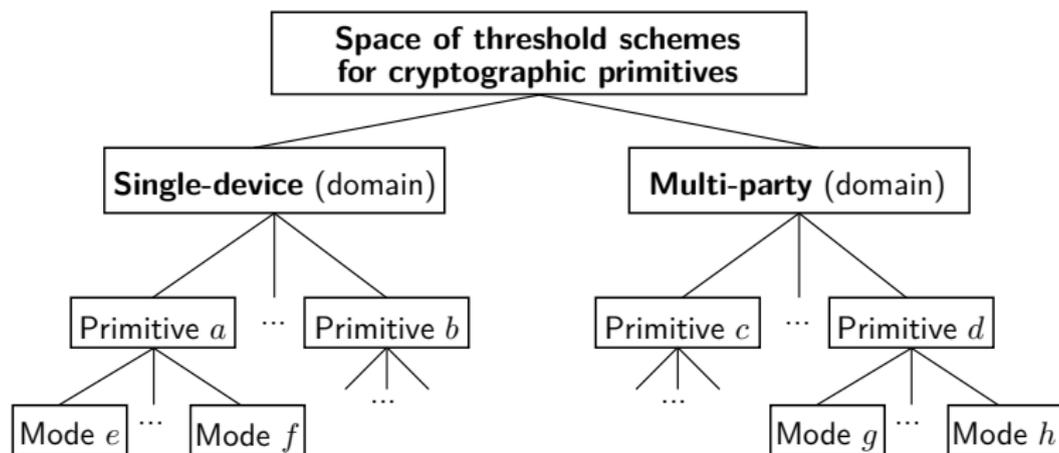


Open to public comments: 2019/Nov/11 – 2020/Feb/10.

Mapping the space of potential “schemes”



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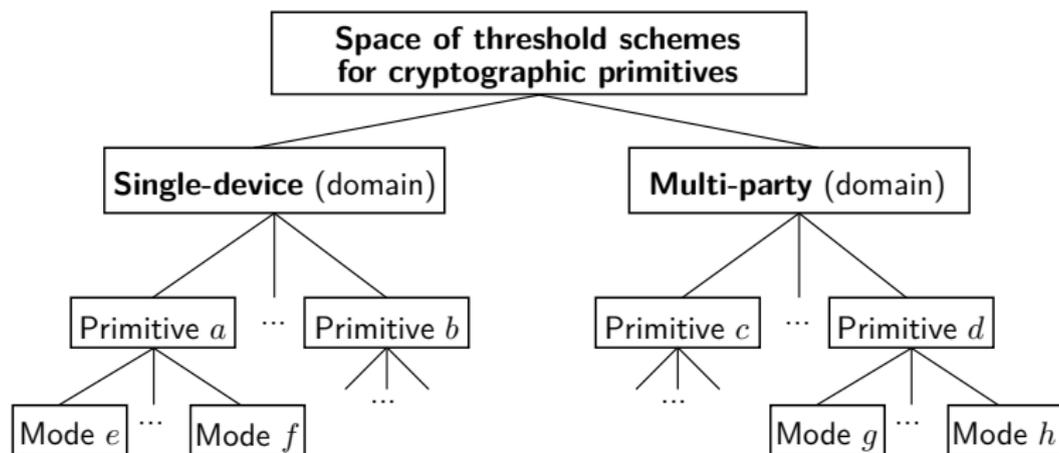
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Multi-party:

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Each **domain** also represents a **track** in the standardization effort.

Some conceivable primitives (focus on NIST-approved)

Less complex:

- ▶ Multi-party: RSA decrypt & sign; EdDSA/Schnorr* sign; ECC key-gen.
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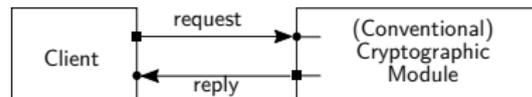
Notes:

- ▶ **Complexity:** depends on more factors, e.g., *, *mode* (next slide).
- ▶ **Other cases:** distributed RNG; some can have similarities across tracks.

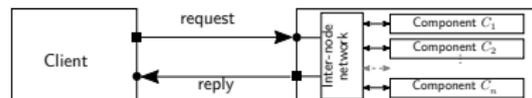
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Input/Output interface: client communication with the module / threshold entity?



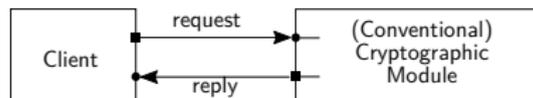
Conventional (non-threshold)



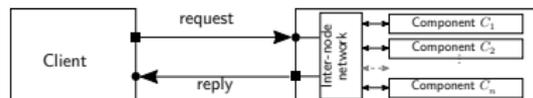
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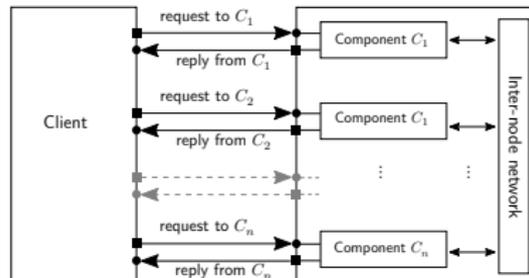
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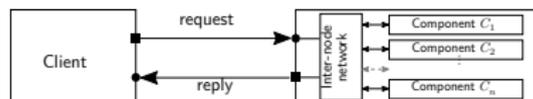
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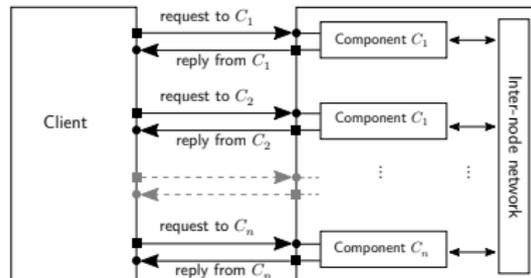
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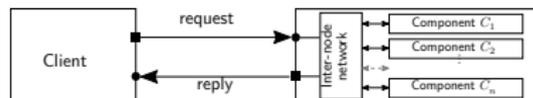
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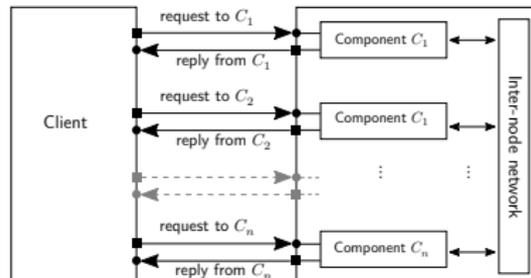
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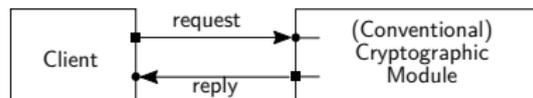
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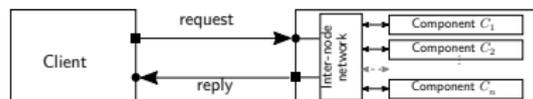
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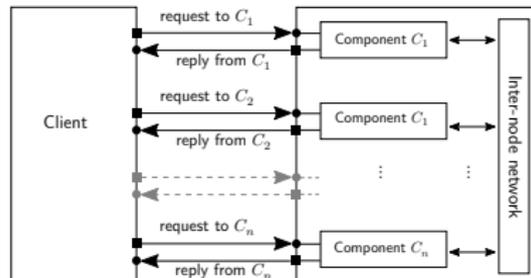
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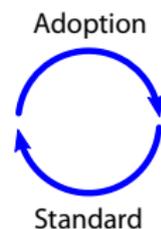
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Examples:

- ▶ Shared-I: signature protecting the secrecy of the input
- ▶ Shared-O: decryption protecting the secrecy of the output
- ▶ Auditable: succinct multi-signature verifiable against several public-keys

Standardization vs. adoption

“not every conceivable possibility is suitable for standardization”

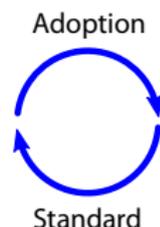


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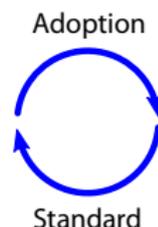
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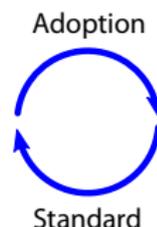
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best practices; minimum defaults; interoperability; innovation.



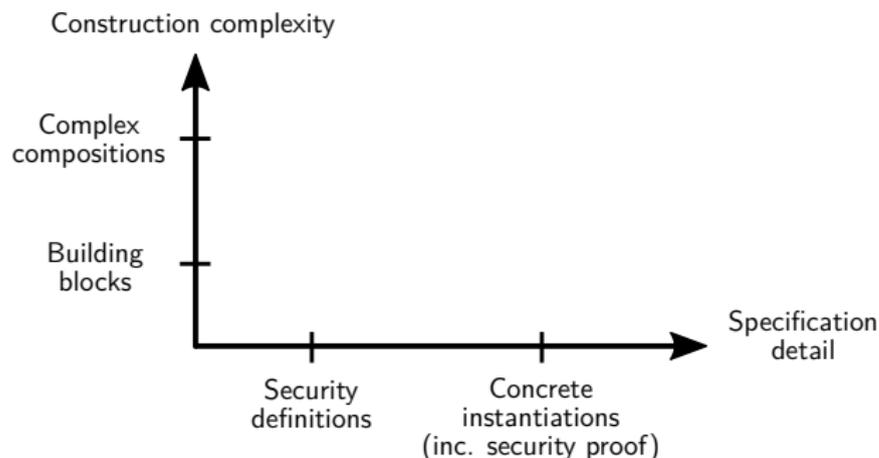
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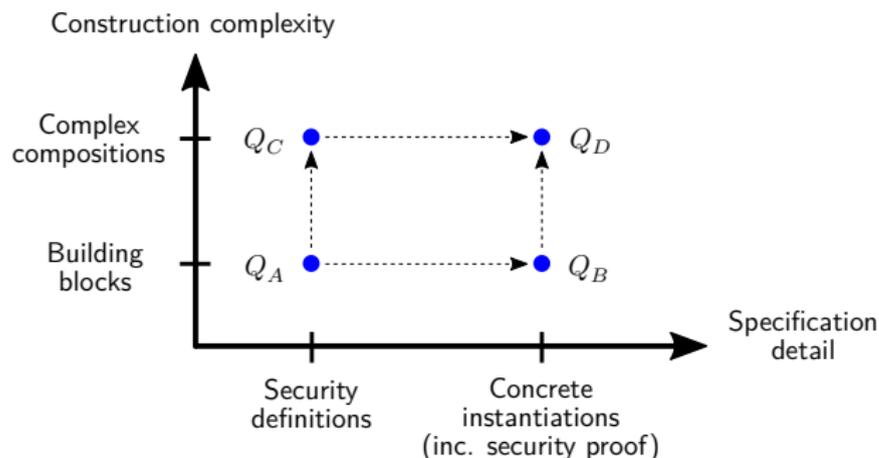
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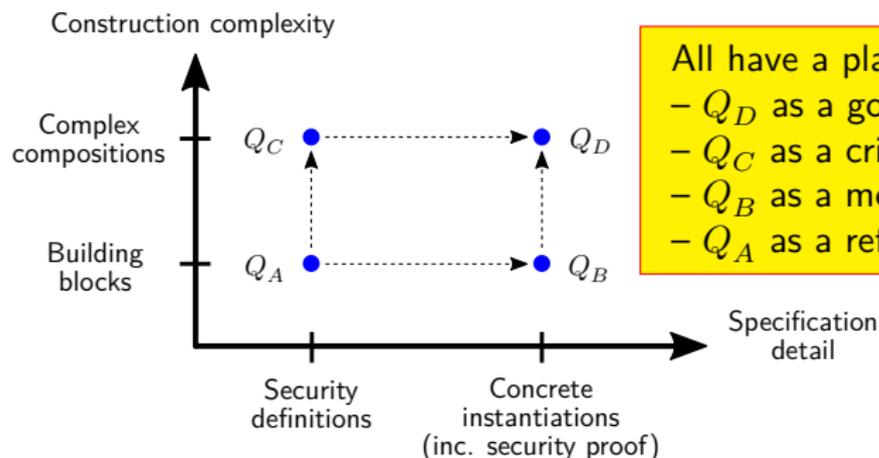
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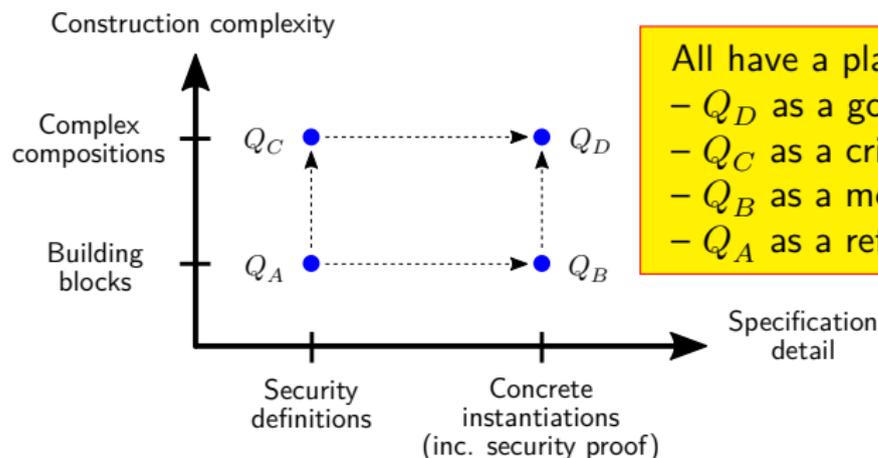
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Example possible *gadgets*: secret sharing; distributed/correlated randomness; consensus; oblivious transfer; garbled circuits; ...

Designing concrete threshold schemes

Additional features to consider:

- ▶ Configurability of threshold parameters
- ▶ Rejuvenation of components (shares, parties, ...)
- ▶ Security (functionality/properties): composable?, adaptive?, graceful degradation?, ...
- ▶ Suitability for testing and validation
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Important:

- ▶ Useful to get feedback from stakeholders about concrete examples
- ▶ These may help define criteria for calls / evaluation / selection

Development process

Generic possible sequence of phases:

1. Roadmap → **2. Calls with criteria** → **3. Evaluation** → **4. Issue standards**

(Each phase to include public feedback. Some Threshold Cryptography workshops along the way?)

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Different standardization *items* can have **different**:

- ▶ **calls for contributions:** feedback on reference protocols; new protocols; reference implementations showing feasibility; research results, ...
- ▶ **timelines** (e.g., depending on complexity; existing rationale for choices)
- ▶ **final formats:** addendum vs. standalone standard, reference to other standards, implementation/validation guidelines, reference definitions,

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- ▶ Asking for disclosure of patents: *call* for disclosure, conditions for submitting
- ▶ Promote “FRAND” license: **f**air, **r**easonable, **a**nd **n**on-discriminatory*
 - * the NIST-ITL patent policy puts it as “reasonable and demonstrably free from unfair discrimination”
- ▶ Cannot force third party to disclose or enable FRAND terms ... but can choose to specify guidance based on expectation of FRAND terms.

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Excerpt from NIST-ITL patent policy: *“assurance [...] that [...] party does not hold [...] any essential patent claim(s); or that a license [...] will be made available [...] under reasonable terms and conditions that are demonstrably free of any unfair discrimination;”*
[possibly without compensation]

Excerpt from NISTIR 7977: *“NIST has noted a strong preference among its users for solutions that are unencumbered by royalty-bearing patented technologies. NIST has observed that widespread adoption of cryptographic solutions that it has developed has been facilitated by royalty-free licensing terms.”* [...]

“NIST will explicitly recognize and respect the value of IP and the need to protect IP if it is incorporated into standards or guidelines.”

Outline 5

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- ▶ NIST-CSD is driving an effort to standardize threshold schemes for NIST-approved cryptographic primitives
- ▶ Collaboration with stake-holders is essential
- ▶ We are in the stage of building a roadmap ... your feedback can (and should) help determine the outcome
- ▶ A two track approach (multi-party and single-device)
- ▶ Various standardization items in each track, with various complexities

The test of time

70 years from now, will *threshold schemes* (still) be used to enable distributed trust in the implementation and operation of cryptographic primitives?

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Photo in 1948 *

Photo in 2018: https://www.nist.gov/sites/default/files/documents/2018/06/15/nist_gaithersburg_master_plan_may_7_2018.pdf

The NIST Stone Test Wall: “Constructed [in 1948] to study the performance of stone subjected to weathering. It contains 2352 individual samples of stone, of which 2032 are domestic stone from 47 states, and 320 are stones from 16 foreign countries.”

* <https://www.nist.gov/el/materials-and-structural-systems-division-73100/nist-stone-wall>

- ▶ Project webpage: <https://csrc.nist.gov/Projects/Threshold-Cryptography>
- ▶ Project email adress: threshold-crypto@nist.gov
- ▶ NISTIR 8214: <https://csrc.nist.gov/publications/detail/nistir/8214/final>
- ▶ NISTIR 8214A (draft): <https://csrc.nist.gov/publications/detail/nistir/8214a/draft>
- ▶ TC-forum: <https://list.nist.gov/tc-forum>

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- [BDV19] L. T. A. N. Brandão, M. Davidson, and A. Vassilev. *Towards NIST Standards for Threshold Schemes for Cryptographic Primitives: A Preliminary Roadmap*. Draft NISTIR 8214A, November 2019. DOI:[10.6028/NIST.IR.8214A-draft](https://doi.org/10.6028/NIST.IR.8214A-draft).
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