

A Side-Channel Assisted Attack on NTRU

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NTRU

KeyGen(*seed*)

1. $((\mathbf{f}, \mathbf{f}_p, \mathbf{h}_q), \mathbf{h}) \leftarrow \text{KeyGen}'(\textit{seed})$
2. $s \leftarrow_{\$} \{0, 1\}^{256}$
3. return $((\mathbf{f}, \mathbf{f}_p, \mathbf{h}_q, s), \mathbf{h})$

Encapsulate(\mathbf{h})

1. $\textit{coins} \leftarrow_{\$} \{0, 1\}^{256}$
2. $(\mathbf{r}, \mathbf{m}) \leftarrow \text{Sample_rm}(\textit{coins})$
3. $\mathbf{c} \leftarrow \text{Encrypt}(\mathbf{h}, (\mathbf{r}, \mathbf{m}))$
4. $k \leftarrow H_1(\mathbf{r}, \mathbf{m})$
5. return (\mathbf{c}, k)

Decapsulate($((\mathbf{f}, \mathbf{f}_p, \mathbf{h}_q, s), \mathbf{c})$)

1. $(\mathbf{r}, \mathbf{m}, \textit{fail}) \leftarrow \text{Decrypt}((\mathbf{f}, \mathbf{f}_p, \mathbf{h}_q), \mathbf{c})$
2. $k_1 \leftarrow H_1(\mathbf{r}, \mathbf{m})$
3. $k_2 \leftarrow H_2(s, \mathbf{c})$
4. if $\textit{fail} = 0$ return k_1
5. else return k_2



Decrypt Implementation

```
int owcpa_dec(unsigned char *rm, const unsigned char *ciphertext,
const unsigned char *secretkey) {
    int i;
    int fail;
    poly x1, x2, x3, x4;

    poly *c = &x1, *f = &x2, *cf = &x3;
    poly *mf = &x2, *finv3 = &x3, *m = &x4;
    poly *liftm = &x2, *invh = &x3, *r = &x4;
    poly *b = &x1;

    poly_Rq_sum_zero_frombytes(c, ciphertext);
    poly_S3_frombytes(f, secretkey);
    poly_Z3_to_Zq(f̄);

    poly_Rq_mul(cf, c, f̄);
    poly_Rq_to_S3(mf, cf);

    poly_S3_frombytes(finv3, secretkey+NTRU_PACK_TRINARY_BYTES);
    poly_S3_mul(m, mf, finv3);
```



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Mapping \mathbb{Z}_3 To \mathbb{Z}_q

- Maps $\{0, 1, 2\}$ to $\{0, 1, q-1\}$
- Highlighted intermediate result:
 - “...0000” if coefficient was 1 or 0
 - “...1111” if coefficient was 2

```
/* Map {0, 1, 2} -> {0,1,q-1} in place */
void poly_Z3_to_Zq(poly *r) {
    int i;
    for (i = 0; i < NTRU_N; i++) {
        r->coeffs[i] = r->coeffs[i] | ((-r->coeffs[i] >> 1) & (NTRU_Q-1));
    }
}
```



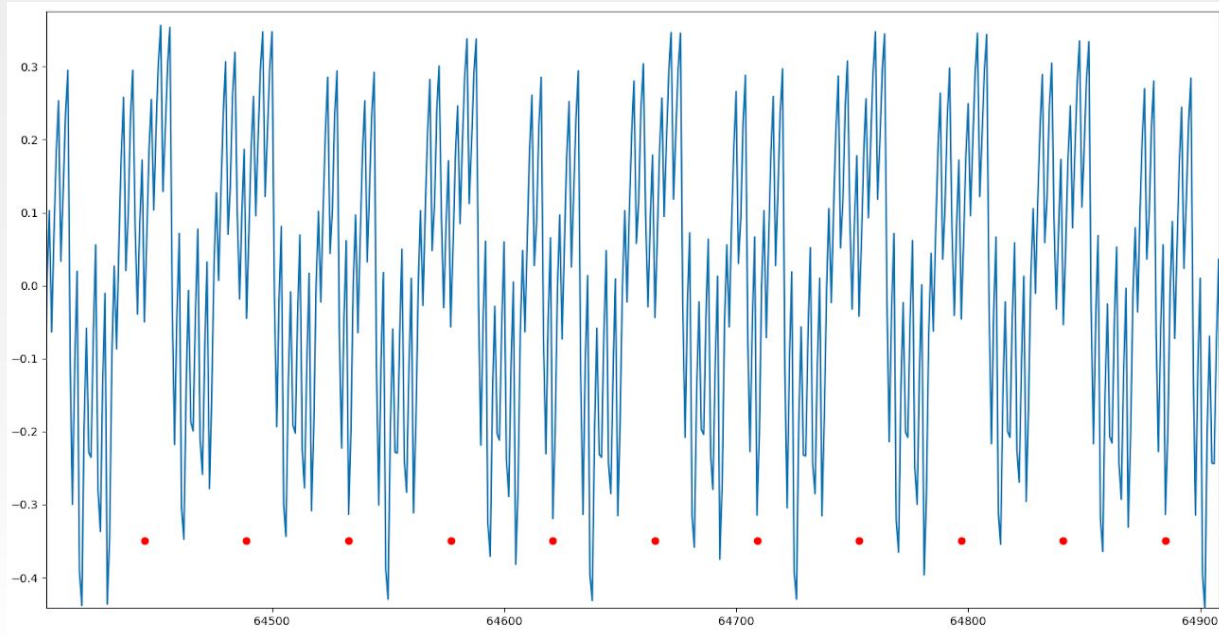
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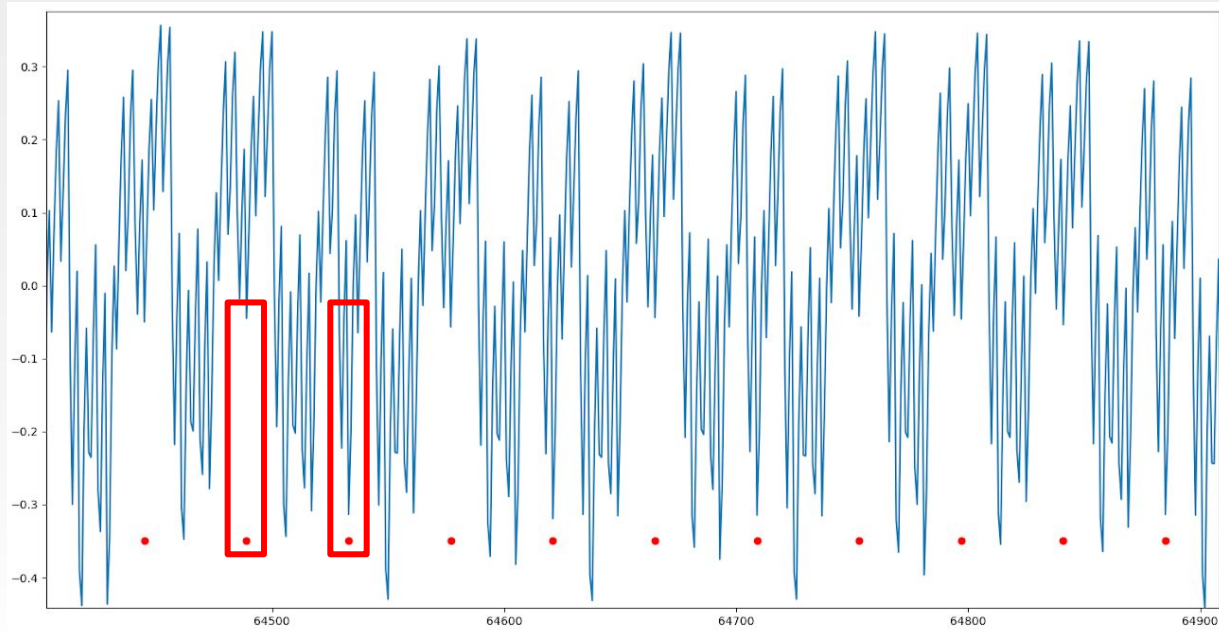
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Power Measurement



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Packing Coefficients

- Five consecutive coefficients packed as $b = \sum_{i=0}^4 f_i \cdot 3^i$
 - [2, 1, 0, 0, 1] -> 86
- Unpacking in two steps
 - 86 -> [86, 28, 9, 3, 1]
 - [86, 28, 9, 3, 1] -> [2, 1, 0, 0, 1]



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Modulo 3

- Highlighted intermediate result:
 - “...000” if $a \equiv 0 \pmod{3}$ and $a \neq 0$
 - “...000” if
 - “...111” otherwise

```
static uint16_t mod3(uint16_t a) {
    uint16_t r;
    int16_t t, c;
    r = (a >> 8) + (a & 0xff); // r mod 255 == a mod 255
    r = (r >> 4) + (r & 0xf); // r' mod 15 == r mod 15
    r = (r >> 2) + (r & 0x3); // r' mod 3 == r mod 3
    r = (r >> 2) + (r & 0x3); // r' mod 3 == r mod 3
    t = r - 3;
    c = t >> 15;
    return (c&r) ^ (~c&t);
}
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$S = \{79, 94, 109, 124, 127, 139, 142, 154, 157, 169, 172, 175, 184, 187, 190, 199, 202, 205, 214, 217, 220, 223, 229, 232, 235, 238\}$

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Partial Key Recovery

- Iterate 3^5 candidates for quintuples
 - Discard those not matching measurements
- On average we recover 75% of f



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Full Key Recovery

- Apply lattice reduction

$$\mathbf{B} = \begin{pmatrix} \mathbf{I} & \mathbf{H} \\ \mathbf{0} & q \cdot \mathbf{I} \end{pmatrix}$$

$$(\mathbf{f}, \mathbf{k}) \mathbf{B} = (\mathbf{f}, \mathbf{g})$$



Some Remarks

- Relies only on very strong leakages
 - Robust and single trace
- Leakage can be reduced
- Implementation does not claim SCA protection



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Questions?

