

# Cryptanalysis of Internal *Keyed* Permutation of FLEXAEAD

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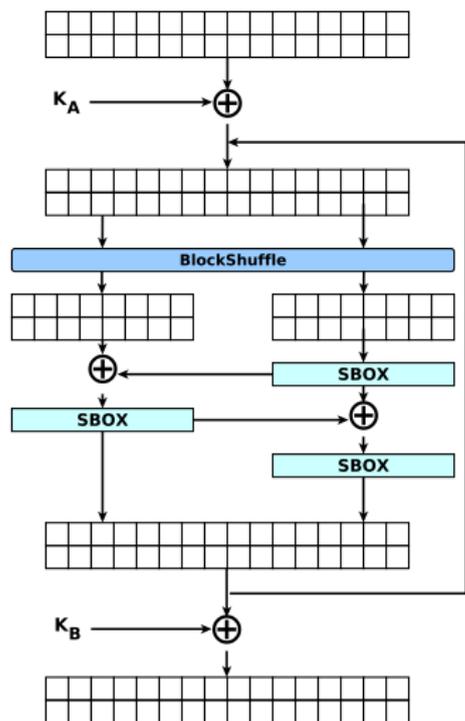


**Lightweight Cryptography Workshop 2019**

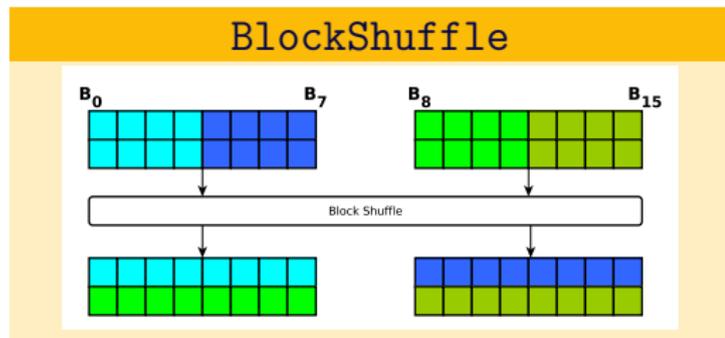
# Introduction

- ▶ FLEXAEAD is round 1 candidate of NIST LWC
- ▶ The underlying Blockcipher is *Internal Keyed Permutation*
- ▶ Block Size can be 64-bit, 128-bit or 256-bit
- ▶ Reported Key Recovery Attack for each variant
- ▶ The attacks are of two type
  1. Iterated Truncated Differential
  2. Yoyo Attacks

# Internal *Keyed* Permutation of FLEXAEAD

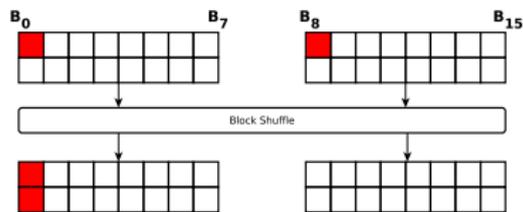


1.  $x$ -bit Flex state is called FLEX- $x$
2. FLEX-128 round function
3. State Bifurcation
4. AES Sbox is used
5. Repeated several times



# Key Observations

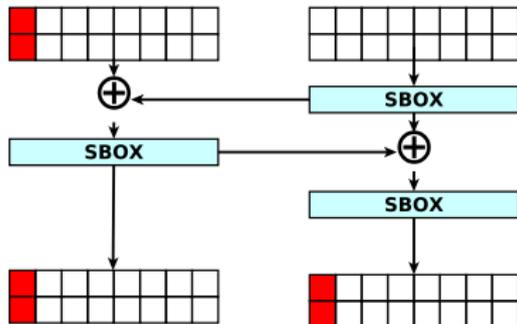
## Effect of BlockShuffle



- ▶ Same Nibble in “Symmetric Bytes” transits to a single byte
- ▶ Number of active bytes can be decreased from two to one

# Key Observations

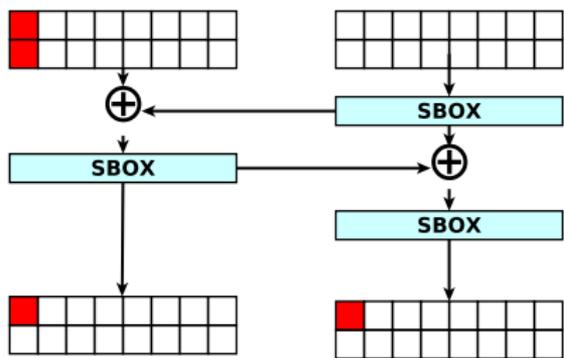
## Effect of SBoxes



- ▶ Due to the effect of XOR, one active byte activates two bytes
- ▶ A pair of “Symmetric Byte” activates a pair of “Symmetric Byte”

# Key Observations

## Effect of SBoxes: Byte to Nibble Transition



- ▶ Only upper or lower nibbles of “Symmetric Bytes” are activated
- ▶ If initially a pair of “Symmetric Bytes” are active, this event occurs with equal probability

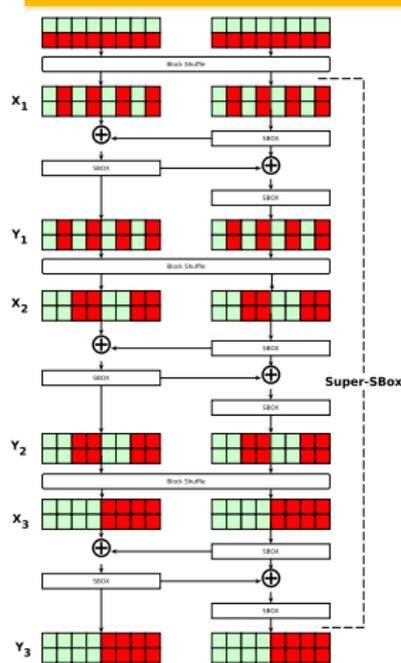
## Exploiting AES Sbox

$$\left| \begin{array}{l} \{(x_1, x_2) \mid (S(x_1) \oplus S(x_2)) \& 0xf0 = 0, \forall x_1, x_2 \in \mathbb{F}_{2^8}\} \\ \{(x_1, x_2) \mid (S(x_1) \oplus S(x_2)) \& 0x0f = 0, \forall x_1, x_2 \in \mathbb{F}_{2^8}\} \end{array} \right| = 4096$$
$$\left| \begin{array}{l} \{(x_1, x_2) \mid (S(x_1) \oplus S(x_2)) \& 0xf0 = 0, \forall x_1, x_2 \in \mathbb{F}_{2^8}\} \\ \{(x_1, x_2) \mid (S(x_1) \oplus S(x_2)) \& 0x0f = 0, \forall x_1, x_2 \in \mathbb{F}_{2^8}\} \end{array} \right| = 4096$$

With probability  $2^{-7}$  two bytes transits to either upper or lower nibble

# Key Observations

## SuperSBox

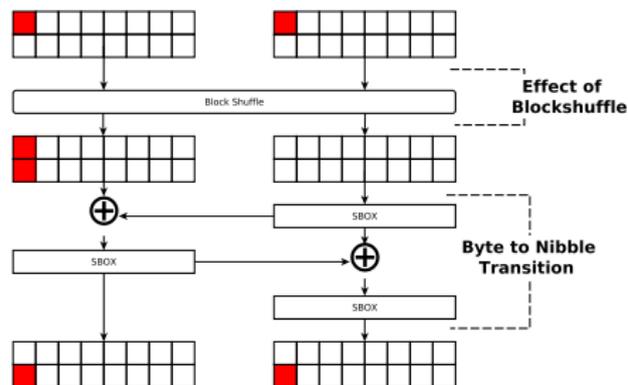


- ▶ Two Super-Sbox exists in  $F_{\text{LEX-128}}$
- ▶ Initial BlockShuffle Layer is not considered in the Super-Sbox
- ▶ Super-Sbox spans over 2.5 round
- ▶ Each Super-Sbox is of 64-bit
- ▶ Super-Sbox in  $F_{\text{LEX-64}}$  and  $F_{\text{LEX-256}}$  spans over 1.5 and 3.5 round respectively

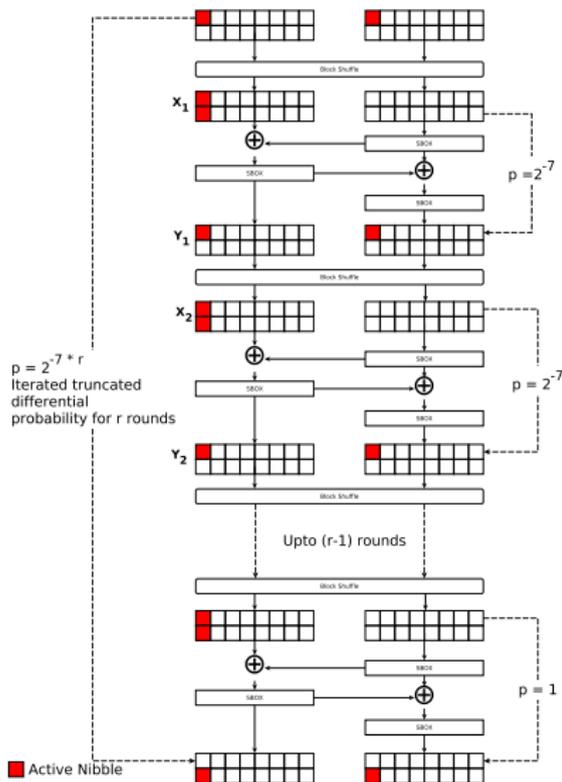
# Iterated Truncated Differential

# One Round Truncated Differential

- ▶ *Effect of BlockShuffle and Byte to Nibble Transition is Combined*
- ▶ The active nibbles in initial state and final state are in same position at the cost of  $2^{-7}$

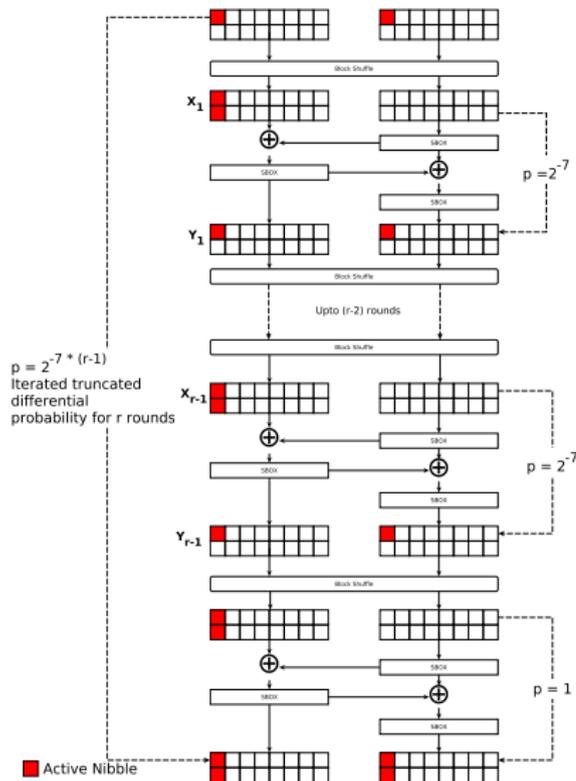


# Iterated Truncated Differential



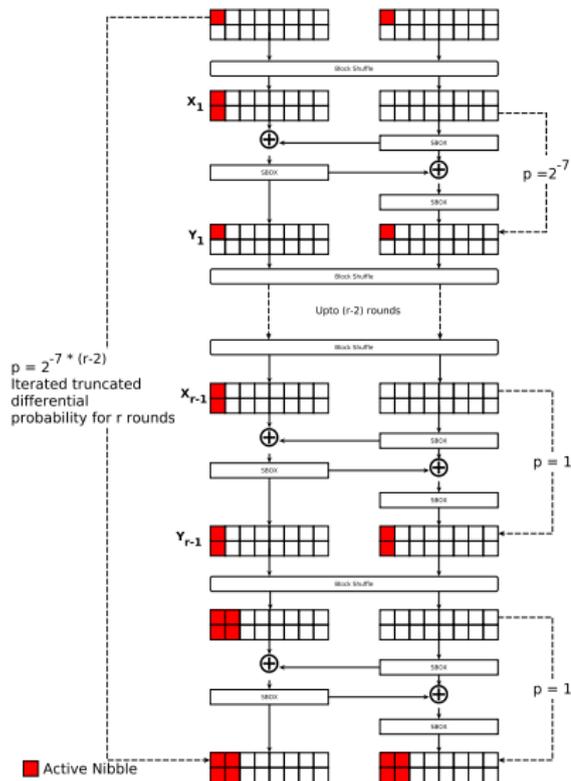
- ▶ The truncated differential can be iterated for  $r$  rounds
- ▶ Paying probability for  $r$  rounds
- ▶ Cost of the trail is  $2^{-7*r}$
- ▶ Some rounds at the end can be made free

# Iterated Truncated Differential: Free Rounds=1



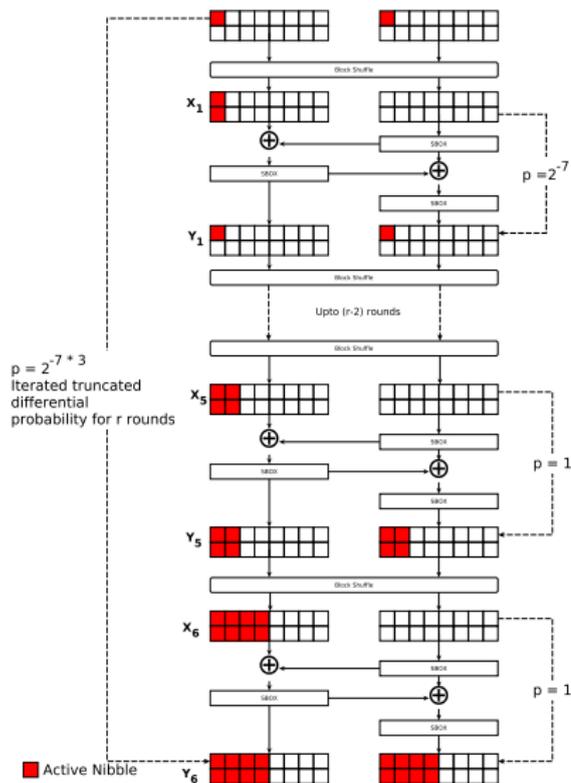
- ▶ 2 bytes are fully active
- ▶ Paying probability for  $r - 1$  rounds
- ▶ Cost of the trail is  $2^{-7*(r-1)}$

# Iterated Truncated Differential: Free Rounds=2



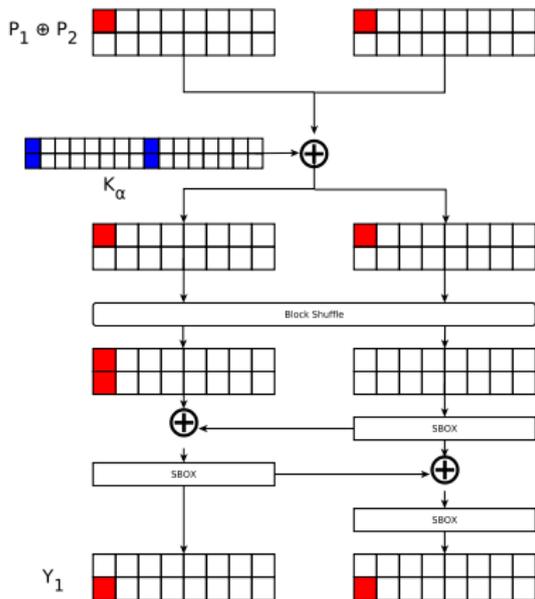
- ▶ 4 bytes are fully active
- ▶ Paying probability for  $r - 2$  rounds
- ▶ Cost of the trail is  $2^{-7*(r-2)}$

# Iterated Truncated Differential: Distinguisher



- ▶ Number of free rounds is 3
- ▶ Probability of 6-round FLEX-128 distinguisher is  $2^{-7*3}$
- ▶ In similar way, number of free rounds in 5-round FLEX-64 and 7-round FLEX-256 is 2 and 4 respectively

# Iterated Truncated Differential: Key Recovery



- ▶ Find a right pair  $(P_1, P_2)$ , such that difference is in byte 0 and 8
- ▶ Guess Key byte 0 and 8 ( $2^{16}$  possible guesses)
- ▶ Run one round encryption and check whether same of byte 0 and 8 are active or not in  $Y_1$  ( $2^9$  key candidates remain)
- ▶ Use two more right pairs to reduce key candidates to 1
- ▶ Repeat the procedure for 8 more byte pairs

# Iterated Truncated Differential Attacks: Summary

| Block Size | #rounds | Data Complexity |      | Time Complexity | Memory Complexity |
|------------|---------|-----------------|------|-----------------|-------------------|
|            |         | Encs            | Decs | MA              |                   |
| 64         | 7       | $2^{30.5}$      |      | $2^{34.5}$      | $2^{18.5}$        |
| 128        | 16      | $2^{93.5}$      |      | $2^{108.5}$     | $2^{20.5}$        |
| 256        | 21      | $2^{109.5}$     |      | $2^{125.5}$     | $2^{22.5}$        |

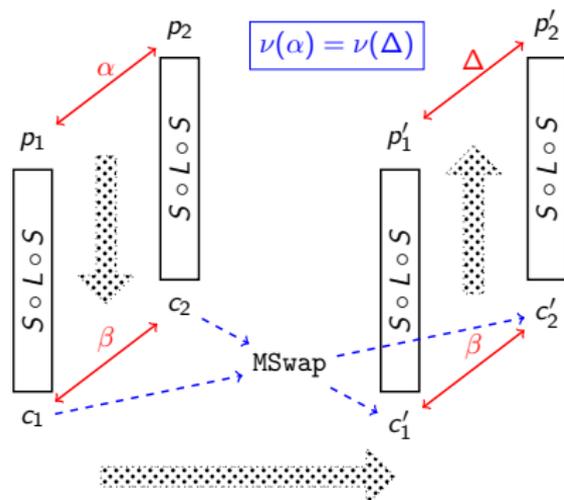
# Yoyo Attacks



$$G'_2 = L \circ S \circ L \circ S$$

Two full generic Rounds

$$G_2 = S \circ L \circ S \quad \leftarrow \text{Dropping final linear layer (to simplify)}$$



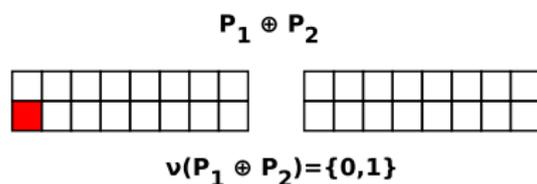
- ▶  $\nu$  is the Zero Difference Pattern

Applied to AES

- ▶ First key-independent Yoyo distinguishers of AES
- ▶ 5-round Key Recovery

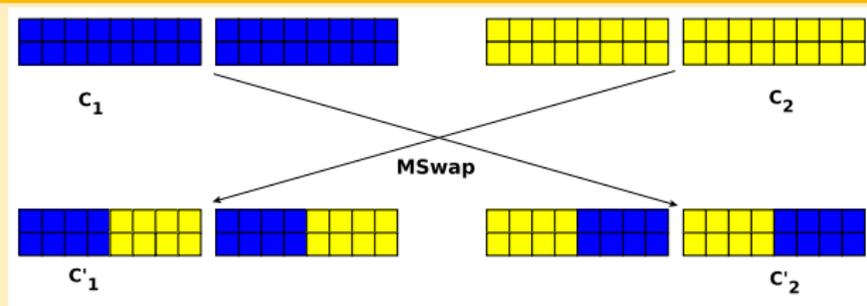
# The Yoyo Trick

## Zero Difference Pattern



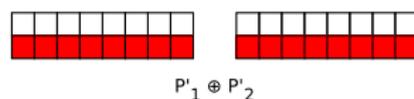
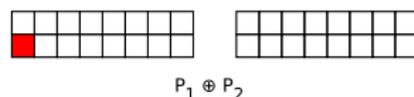
- ▶ Two Super-Sbox in FLEX-128 state
- ▶ A fully inactive Super-Sbox is denoted by 1; otherwise, 0

## MSwap



- ▶ Bytes are swapped between two texts according Super-Sbox output

# Yoyo Attacks: Deterministic Distinguisher

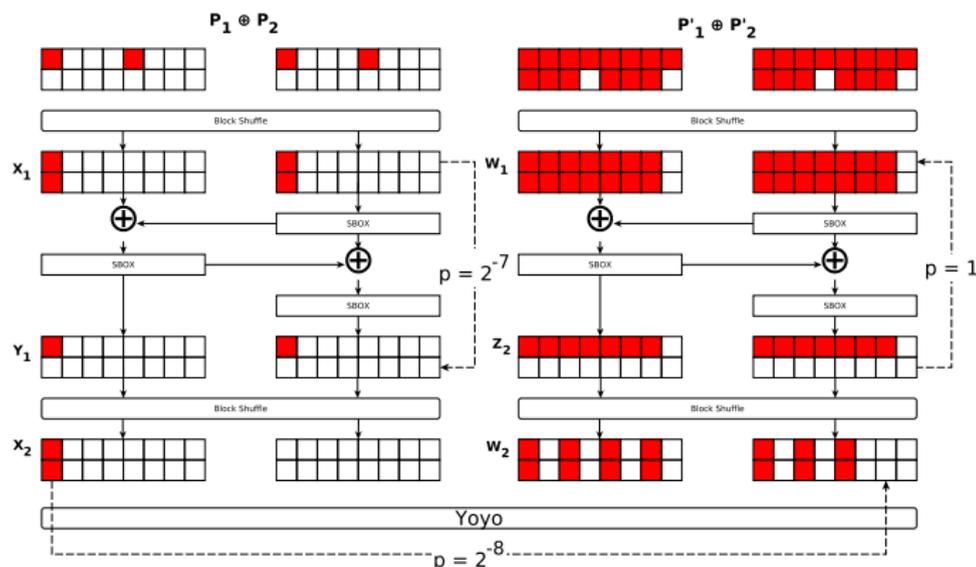


- ▶ Super-Sbox and BlockShuffle are considered as  $S$  and  $L$  layer respectively
- ▶ FLEX-128 Super-Sbox spans over 2.5 rounds
- ▶ 6-round FLEX-128 Deterministic Distinguisher
- ▶ Apply Yoyo game
  1.  $P_1, P_2 \xrightarrow{ENC} C_1, C_2$
  2.  $C_1, C_2 \xrightarrow{MSwap} C'_1, C'_2$
  3.  $C'_1, C'_2 \xrightarrow{DEC} P'_1, P'_2$

# Yoyo Attacks: Key Recovery

- ▶ 6-round Deterministic Distinguisher is the building block of 7-round FLEX-128 Key Recovery attack
- ▶ Byte to Nibble Transition is used to extend for 1 round
- ▶ Similar kinds of attacks exist for FLEX-64 and FLEX-256

# Yoyo Attacks: Key Recovery



- ▶ Choose  $P_1, P_2$  and encrypt them to obtain  $C_1, C_2$
- ▶ Apply  $MSwap$  on  $C_1, C_2$  and decrypt them to get  $P'_1, P'_2$
- ▶ Any one of the 8 active Bytes in  $W_2$  can be zero w.p.  $2^{-5}$
- ▶ Trail probability is  $2^{-12}$
- ▶ Key Recovery part is same as Iterated Truncated Differential

# Yoyo Attacks: Summary

| Block Size | #rounds | Data Complexity |            | Time Complexity | Memory Complexity |
|------------|---------|-----------------|------------|-----------------|-------------------|
|            |         | Encs            | Decs       | MAAs            |                   |
| 64         | 5       | $2^{10}$        | $2^{16.5}$ | $2^{15.5}$      | $2^{10}$          |
| 128        | 7       | $2^{10.5}$      | $2^{16.5}$ | $2^{16.5}$      | $2^{11.5}$        |
| 256        | 9       | $2^{11}$        | $2^{16.5}$ | $2^{17.5}$      | $2^{13}$          |

# Attacks Presented in this Work

| Block Size | #rounds | Data Complexity |            | Time Complexity | Memory Complexity | Attack Type                     |
|------------|---------|-----------------|------------|-----------------|-------------------|---------------------------------|
|            |         | Encs            | Decs       | MA's            |                   |                                 |
| 64         | 7       | $2^{30.5}$      |            | $2^{34.5}$      | $2^{18.5}$        | Iterated Truncated Differential |
|            | 5       | $2^{10}$        | $2^{16.5}$ | $2^{15.5}$      | $2^{10}$          | Yoyo Attack                     |
| 128        | 16      | $2^{93.5}$      |            | $2^{108.5}$     | $2^{20.5}$        | Iterated Truncated Differential |
|            | 7       | $2^{10.5}$      | $2^{16.5}$ | $2^{16.5}$      | $2^{11.5}$        | Yoyo Attack                     |
| 256        | 21      | $2^{109.5}$     |            | $2^{125.5}$     | $2^{22.5}$        | Iterated Truncated Differential |
|            | 9       | $2^{11}$        | $2^{16.5}$ | $2^{17.5}$      | $2^{13}$          | Yoyo Attack                     |

# Conclusion

1. Reported Iterated Truncated Differential which exploits AES Sbox and BlockShuffle operation
2. Generalized Yoyo Distinguishing Attack is applicable
3. All attacks are exploited to recover subkeys
4. Practical ones are experimentally verified
5. FLEXAEAD is out of 2nd round

Thank You