pqm4: Benchmarking NIST Additional Post-Quantum Signature Schemes on Microcontrollers

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Abstract. In July 2022, the US National Institute for Standards and Technology (NIST) announced the first set of Post-Quantum Cryptography standards: Kyber, Dilithium, Falcon, and SPHINCS+. Shortly after, NIST published a call for proposals for additional post-quantum signature schemes to complement their initial portfolio. In 2023, 50 submissions were received, and 40 were accepted as round-1 candidates for future standardization.

In this paper, we study the suitability and performance of said candidates on the popular Arm Cortex-M4 microcontroller. We integrate the suitable implementations into the benchmarking framework pqm4 and provide benchmarking results on the STM32L4R5ZI featuring 640 KB of RAM. pqm4 currently includes reference implementations for 15 submissions and M4-optimized implementations for five submissions. For the remaining candidates, we describe the reasons that hinder integration – the predominant reason being large key size or excessive memory consumption.

While the performance of reference implementations is rather meaningless and often does not correlate with the performance of well-optimized implementations, this work provides some first indication of which schemes are most promising on microcontrollers. The publicly available implementations in pqm4 also provide a good starting point for future optimization efforts.

Initially, we were hoping for a much higher code quality than for initial submissions to NIST's previous PQC project. However, we got grossly disappointed: Half of the submissions make use of dynamic memory allocations, often completely without reason; Many implementations have compiler warnings, sometimes hinting at more serious issues; Many implementations do not pass simple sanitizer tests such as using valgrind; Multiple implementations make use of static memory.

1 Introduction

While NIST's initial portfolio of post-quantum cryptography consisting of Kyber [SAB+19], Dilithium [LDK+19], Falcon [PFH+19], and SPHINCS+ [HBD+19] provides reasonable performance characteristics for most applications, there are some use-cases that benefit from different trade-offs in terms of key, ciphertext, and signature sizes. Therefore, NIST as well as the cryptography community continues to study schemes achieving different performance characteristics: In Round 4 of the NIST PQC project [NIS23b], there are three code-based key-encapsulation mechanisms (KEMs) with the most prominent candidate being Classic McEliece [BCC⁺22] achieving much smaller ciphertexts than Kyber, but requiring very large public keys. Additionally, NIST [NIS23a] has issued another call for additional post-quantum digital signatures to complement their lattice-based and hash-based signature schemes. The call is primarily for schemes not based on structured lattices, and achieving short signatures and fast verification. By the deadline in June 2023, NIST had received 50 submissions of which they accepted 40 submissions as *complete-and-proper*. These schemes are now being studied in the first round of evaluation.

The first round of evaluation in cryptographic competitions is commonly focussing on the security of the submitted schemes with numerous schemes being broken within hours of appearing publicly [AAS⁺19]. However, as submission teams were required to submit portable C reference implementations as a part of their submission package, we can also evaluate implementation performance and identify any problems with the

^{*} Date of this document: 2024-01-25. pqm4 version: 62244ef.

submitted implementations. This allows cryptographic engineers to identify promising optimization targets for a wide variety of platforms and publish numerous papers by comparing performance to the reference implementations. The performance of the resulting optimized implementations commonly plays an important role in the selection process beyond round 1 [AASA⁺20].

Besides the x64 reference platform, NIST is explicitly interested in additional performance results on constrained devices, such as smart cards and microcontrollers. NIST has designated the Arm Cortex-M4 as the primary optimization target in the class of microcontrollers [AASA⁺20]. During the evaluation of the candidates for the first NIST PQC call, the pqm4 framework [KRSS19] emerged as the de facto standard evaluation platform for PQC implementations targeting the Arm Cortex-M4 microcontroller. It provides an easy-to-use and automated evaluation framework that unifies benchmarking and implementations of underlying symmetric cryptography to enable fair and meaningful performance comparisons. The pqm4 repository includes the majority of relevant and publicly available implementations optimized for the Cortex-M4 of selected and remaining PQC candidates from the first NIST call.

In this paper, we extend this repository and evaluate *all* 40 additional signature candidates in terms of their suitability for embedded devices and provide benchmarks for multiple performance characteristics. We extended pqm4 with 15 reference implementations and five optimized implementations. For the remaining schemes there are either publicly announced severe vulnerabilities, current implementations require too much memory to be functional on any Arm Cortex-M4 platform, or would require major engineering effort to allow porting to pqm4 and embedded devices in general.

In Section 2 we briefly introduce the pqm4 framework and the major changes since the first report [KRSS19]. Following, in Section 3 we list the implementations added to pqm4 or the reasons that prohibit us from doing so for each of the 40 schemes. The performance results can be found in Section 4.

2 pqm4

The pqm4 project is a reference framework for benchmarking implementations of PQC KEM and signature schemes for Arm Cortex-M4 based embedded systems. The goal is not just to provide a common ground for comparing the various PQC schemes, but also different implementations of a single scheme. To that end, a set of tests are used to determine performance criteria important to embedded platforms of the basic operations of each scheme: key generation, de- and encapsulation for KEM schemes, or signature generation and verification for signature schemes. The performance criteria covered here are the speed in terms of CPU cycles spent on the operations, as well as static- and dynamic memory usage. Furthermore, a profiling test is used to determine the cycles spent on symmetric primitives, which many schemes use for hashing or as PRNGs for sampling uniformly random bits. The memory usage is analyzed in terms of code size, static memory used, for example, as look-up-tables, and stack-usage during execution. Aside from these performance tests, the framework also includes functional tests, which check the expected behavior of the implementations' APIs. This covers normal functional tests, tests against invalid inputs (invalid secret key or ciphertext during decryption, or invalid public key during verification), and a testvector test, which checks whether each implementation produces the same outputs as a reference implementation on a host platform.

History. The initial framework covered 10 KEM and 3 signature schemes (in multiple security levels for some schemes) of the first round of the NIST PQC standardization process. During this first round, the framework was structured as a monolithic framework, targeting only one Arm Cortex-M4 platform. For the second round, the framework was split up into pqm4 and mupq. The former included only platform specific implementations for the Cortex-M4, while the latter covered platform independent implementations suitable for 32 bit embedded platforms and further integrates implementations from the pqclean project [KSSW22]. This split facilitated the creation of pqm3⁴ and pqriscv⁵, which aim to cover Arm Cortex-M3 and RISC-V based embedded platforms. For the third round the build system was overhauled from the ground up to support multiple target boards, with a common build system integrated into mupq to be used across pqm4, pqm3, and pqriscv. New target boards supported by pqm4 allowed the use of more memory on larger microcontrollers, or even the QEMU simulator. The simulated platform enables up to 4 MiB flash- and system memory. Lacking

⁴ https://github.com/mupq/pqm3

⁵ https://github.com/mupq/pqriscv

any meaningful time-measurement, the simulator is not useful for performance tests concerning speed, but suitable for all memory measurements and functional tests.

Measurement methodology. The main tool used for measuring the cycle counts of the scheme operations is the SYSTICK timer available Arm Cortex-M4 cores. The "Data Watchpoint and Trace Unit" (DWT) in the core features 32 bit cycle counter, which is, however, insufficient to measure schemes with runtimes exceeding 2^{32} cycles. While the SYSTICK timer itself only uses a 24 bit counter, it can be used for longer measurements, as overflows are caught and counted by an interrupt service routine to measure an arbitrary number of cycles. For the speed benchmarks, the current cycle count is sampled before and after each scheme operation, and the difference taken as the result. Speed tests are repeated, to estimate a minimum-, average-, and maximum runtime for schemes that use random sampling methods. During the profiling tests, the current cycle count is further sampled on entry and exit from symmetric primitives, and the accumulated sum of cycles spent is noted. All schemes integrated in pqm4 are adapted to use the same implementation for a set of common symmetric primitives (SHA-2, SHA-3/SHAKE, AES), which are regularly updated to the current state-of-art in terms of speed. For AES, pqm4 offers a bitsliced constant-time implementation [AP20], as well as a faster variable-time t-table implementation [SS17]. To avoid an impact on the speed measurements by the timings of flash-based code memory, which may differ between manufacturers, the target microcontroller runs at a reduced clock frequency during cycle counting tests. The reduced frequency allows the execution of code without any wait-states, the maximum frequency depends on the microcontroller at hand⁶. The memory benchmarks use the compiler toolchain to determine the static memory usage (code- and static data sizes). To determine the dynamic stack usage during runtime, "stack spraying" is used, i.e., the stack is filled with a pattern and checked up to which size it is modified. As no implementations with dynamic heap memory allocations are integrated into pqm4, no measurement of heap usage is necessary.

Target Platform. Previous versions of pqm4 targeted the STM32F4DISCOVERY evaluation board, which featured the STM32F407VG microcontroller. The board was chosen for its affordable price and wide availability, and with its 1 MiB flash and 192 KiB SRAM offered a reasonably high amount of memory. For this report, however, we opted for the Nucleo-L4R5ZI evaluation board, featuring a STM32L4R5ZI microcontroller. With 2 MiB flash and 640 KiB SRAM, this new target platform offers a significant increase in memory resources, enabling us to benchmark more of the new signature schemes. The core clock frequency of the board is lower with 120 MHz compared to 168 MHz, the benchmarking clock frequency is further reduced to 20 MHz from 24 MHz. Both the old and the new platform feature a TRNG which is used as the entropy source (randombytes()) for the schemes. Our RNG implementation exhibits identical timing performance on both platforms. The memory layout between the platforms differs significantly. On the STM32F407VG, the 192 KiB is made up of three blocks of SRAM: one 112 KiB block, one 16 KiB block, and a 64 KiB core-coupled block. The core-coupled memory was unused by pqm4, the second smaller block was only used for schemes which required the additional memory. This was a deliberate choice, as the second block exhibits different memory timings, taking about 1000 cycles more to read a 4 KiB block of memory. For the new platform, we chose not to avoid any memory blocks. The memory layout of the STM32L4R5ZI consists of three blocks: one 192 KiB, one 64 KiB, and one 384 KiB block. Similar to the old platform, only the first block exhibits the fastest timings. Including the slower blocks affects benchmarking results of schemes relying more heavily on memory reads.

Choice of schemes. Not all of the 40 accepted submissions for the first round of additional post-quantum digital signatures are suitable for benchmarking in the context of pqm4. The first and foremost criteria for inclusion is the size of the key pair and signature, as they have to fit within the constrained memory resources of the chosen embedded platform. At the time of writing, the benchmarking framework is not setup for partial benchmarking of single operations, e.g., benchmarking signature verification of a signature/public key pair generated on an unconstrained system. Hence, only schemes which fit with an entire key pair and signature can be tested. The second criteria is the portability of the available implementations for a scheme. Some of the candidates rely heavily on third party libraries such as GMP or FLINT, which cannot be (easily) ported to an embedded system. Some of the other third party dependencies can be easily replaced. For example, we modified implementations to use our SHA-3/SHAKE implementations, when the original implementation was

⁶ Higher core clock frequencies require the CPU to enter wait states, to wait for the flash memory to respond. In our case up to six clock cycles, resulting in $2-3\times$ higher cycles counts. The effective execution time is, however, amortized by the higher clock speeds.

making use of the eXtended Keccak Code Package⁷ or OpenSSL. The last criteria is the absence of dynamic memory allocations, e.g., using the malloc function. While heap-based dynamic memory allocation can be implemented on embedded systems, their use is generally avoided in favor of static or stack-based allocations. As such, the pqm4 framework only includes implementations that use stack-based allocations.

Porting implementations. Integrating a scheme implementation into pqm4 usually follows similar steps. Some scheme authors already provide suitable implementations, which are integrated easily. In all other cases, we used the reference or optimized implementations included in the submissions accepted by NIST as a startoff point. Wherever necessary, the APIs are adapted to use the correct type system. Next, all calls to symmetric primitives are replaced to those provided by pqm4. If dynamic memory allocations are present, they are replaced by stack allocations. In some cases this is a simple replacement of a pointer with an array of appropriate size, e.g., if the allocated memory is not returned from the function. In a few cases, it may also involve allocating the memory in the calling function and passing a pointer. Implementations using more complex memory allocations were not ported to the framework. Some implementations used lookup tables that are computed on-the-fly and cached for future calls. As these caches are allocated statically, they break the reentrancy of the implementation, and complicate performance measurement. For that reason, implementations were modified to allocate the the lookup tables on the stack and compute them for every invocation of the scheme. Ideally, implementers should only rely on precomputed lookup tables, allocated to flash memory.⁸

When the implementation can be adapted with the previous steps, it is then integrated and tested using the QEMU simulator. Due to the up to 4 MiB of memory available on the simulated mps2-an386 platform, all but the largest schemes can be functionally tested and the stack-usage estimated. The result is a list of all integrated implementations and their respective memory requirements. Implementations that fit the target platform are then tested on real hardware, checked whether they pass the functional- and testvector tests, and finally benchmarked for their execution speed. For code- and stack-size benchmarks, the simulated results using QEMU can be considered, as they are identical to those gathered from real hardware.

3 New signature schemes included in pqm4

Table 1 gives an overview of all schemes submitted to NIST to the call for additional signature schemes. For each scheme, we either list which implementations have been included in pqm4 (and reference the respective pull requests), or list the reasons why they cannot be included. We exclude 9 schemes against which (convincing) attacks have been publicly announced on the NIST pqcforum. Note that we do not aim to provide a complete picture of the state of cryptanalysis, but instead want to focus our engineering efforts on the schemes that are most likely going to advance the next round. For some schemes, the submission teams provided updated versions addressing vulnerabilities. For four schemes the public keys (for all parameter sets) itself are too large to fit into the memory of our target platform (640 KB). For further seven schemes, current implementations use too much memory to fit onto the target platform. Five schemes require external libraries that are not suitable for the Cortex-M4. One schemes does not have portable C-code. 20 out of 40 schemes have reference implementations that make use of dynamic memory allocations. We tried our best to eliminate those dynamic memory allocations and in the vast majority of cases they can be easily replaced by statically-sized buffers or variable-length arrays. Only for MIRA [ABB⁺23c], RYDE [ABB⁺23b], and FAEST [BBdSG⁺23], the dynamic memory allocations were the main reason for exclusion from pqm4 as eliminating those would require significant rewriting of the implementations. Overall, we currently have reference implementations for 15 schemes and M4-optimized implementations for five schemes in pgm4.

In the following we give the details for each scheme.

3.1 Code-based Signatures

- CROSS [BBB+23b]: The CROSS reference implementation from the submission package has been merged into pqm4 in #309. For some parameter sets (e.g., cross-sha3-r-sdpg-1-small) test vectors are inconsistent between the host and the Cortex-M4 unless the signed message is initialized to 0 at the

⁷ https://github.com/XKCP/XKCP

⁸ Memory marked **const** is allocated to flash on embedded systems, depending on the linker script.

				po	qm4			reaso	on(s) for a	exclusio	n	
		issue	\mathbf{PR}	ref	m4f	vuln	pk	mem	not port	ext lib	dyn mem	params
CROSS	$[BBB^+23b]$	#265	#309	1								12/24
Enhanced pqsigRM	$[CNL^+23]$	#270					X				X	0/1
FuLeeca	$[RBK^+23]$	#272				X						0/3
LESS	$[BBB^+23a]$	#278						X			X	0/7
MEDS	$[CNP^+23]$	#280	#324	1								2/6
Wave	$[BCC^+23a]$	#298					X				X	0/3
SQIsign	$[CSSF^+23]$	#293								X	X	0/3
EagleSign	[SHDS23]	#267				X						0/4
EHTv3 and EHTv4	[SF23]	#268				X					X	0/5
HAETAE	$[CCD^+23b]$	#273	#313	1	1							3/3
HAWK	$[BBD^+23]$	#274	#305	1								3/3
HuFu	$[YJL^+23]$	#276					X				×	0/3
Raccoon	$[dPEK^+23]$	#288							X			0/18
SQUIRRELS	[ENST23]	#294					X			X		0/5
Biscuit	[BKPV23]	#264	#314	1								3/6
MIRA	$[ABB^+23c]$	#281									X	0/6
MiRitH	$[ARZV^+23]$	#282	#315	1	1							16/32
MQOM	[FR23]	#283	#322	1							(🗙)	2/12
PERK	$[ABB^+23a]$	#284	#318	1	1							12/12
RYDE	$[ABB^+23b]$	#289									X	0/6
SDitH	$[MFG^+23]$	#290						X			×	0/12
3WISE	[Rod23a]	#260				X				X		0/3
DME-Sign	[LA23]	#266				X						0/3
HPPC	[Rod23b]	#275								X		0/3
MAYO	$[BCC^+23b]$	#279	#302	1	1							3/4
PROV	$[GCF^+23]$	#286						X			×	0/3
QR-UOV	$[FIH^+23]$	#287						X			X	0/12
SNOVA	$[WCD^+23]$	#291	#311	1								7/18
TUOV	$[DGG^+23]$	#295	#327	1				X			×	0/12
UOV	$[BCD^+23]$	#296	#300	1	1							3/12
VOX	$[PCF^+23]$	#297				X					×	0/3
AIMer	$[\mathrm{KCC}^+23]$	#261	#323	1							(X)	3/12
Ascon-Sign	$[SGJ^+23]$	#263	#308	1								8/8
FAEST	[BBdSG ⁺ 23]	#271									×	0/12
SPHINCS-alpha	[YCZ23]	#292	#312	1								6/24
ALTEQ	$[BDN^+23]$	#262						X			×	0/6
eMLE-Sig 2.0	[LZ23]	#269				X					×	0/3
KAZ-SIGN	$[AAC^+23]$	#277				X				X	×	0/3
Preon	$\left[\mathrm{CCC}^{+23}\right]$	#285						X			×	0/9
Xifrat1-Sign.I	[NP23]	#299				X						0/1
				15	5	0	4	7	1	5	20	02/225

Image: Table 1. Overview of the 40 candidates submitted to the NIST call for additional signatures and their status ofinclusion in pqm4. Reasons for exclusion include publicly announced severe vulnerabilities, public key sizes exceedingthe RAM size of the target device (640 KB), implementations exceeding 640 KB of memory, non-portable code,external dependencies incompatible with the Cortex-M4, and dynamic memory allocations. The last column gives thenumber of parameter sets for which reference implementations have been integrated into pqm4 and that are functionalwith 640 KB.

beginning of the signing procedure. We added the initialization, but this hints at missing initialization and possibly unused parts of the signature. For both the sha2 and sha3 variants, spdg-1-fast and the two respective sets with security level 3 and 5 in addition to spdg-1-small, spd-1-fast and spd-2-fast are running on our physical evaluation platform. This makes in total 12 out of the 24 parameter sets.

- Enhanced pqsigRM [CNL⁺23]: The single parameter set of Enhanced pqsigRM (Enh-pqsigRM-613) has public keys of 2.00 MB and signatures of 1023 bytes. 2 MB exceeds the available memory on our target M4 platform, and we hence do not include Enhanced pqsigRM in pqm4.
- FuLeeca [RBK⁺23]: FuLeeca has been shown vulnerable by Hörmann and van Woerden⁹ and the FuLeeca submission team acknowledged the vulnerability. At the time, no updated version of FuLeeca is available and we, hence, do not include FuLeeca in pqm4. Additionally, when studying the implementation of FuLeeca, we found that it is incorrectly implementing SHAKE (mixing calls to an incremental and non-incremental API). We reported this problem to the submission team and they published a patch resolving this problem.¹⁰ This fix changes KATs.
- LESS [BBB⁺23a]: The LESS reference implementation makes use of dynamic memory allocations. Additionally, signing and verification require a buffer (of type normalized_IS_t) of $T \cdot K \cdot (N - K)$ bytes totaling 3.1 MB (LESS-1s) to 102 MB (LESS-5b). This is far beyond the reach of any microcontroller and, hence, we cannot include LESS into pqm4. We also noticed that compute_digest and compute_digest_array do not hash the tree_salt (due to passing a too short length to the hash function). This appears to be incompatible with the specification and fixing the issue will result in different KATs.
- MEDS [CNP⁺23]: The MEDS reference implementation from the submission package has been merged into pqm4 in #324. Two (meds13220 and meds55604) of the six parameter sets are running on the Nucleo-L4R5ZI board.
- Wave [BCC⁺23a]: Wave uses public keys of 3.7 MB (Wave822) to 13.6 MB (Wave1644) and, hence, can not be included in pqm4.

3.2 Isogeny Signatures

- SQIsign [CSSF+23]: Currently all available implementations of SQIsign make heavy use of dynamic memory allocations. This is, in part, due to variable-sized integer arithmetic. Hence, at the moment, we cannot include SQIsign in pqm4. We believe that SQIsign without dynamic memory allocations is possible, but requires significant engineering efforts. Preliminary experiments (on x86) using the reference implementation in the submission package show that the total memory (stack + heap) of signing of SQISign-NIST-I is slightly above 300 KB. Verification requires only about 12 KB of memory and uses significantly less dynamic memory. The implementation relies on the GMP library.

3.3 Lattice-based Signatures

- EagleSign [SHDS23]: EagleSign in the version submitted to NIST has been shown to be vulnerable by Tibouchi.¹¹ Tibouchi published code practically recovering an EagleSign secret key from a sufficient number of signatures.¹² We do not include EagleSign in pqm4.
- EHTv3 and EHTv4 [SF23]: Practical attacks against EHT (as submitted to NIST) were presented by both Postlethwaite and van Woerden¹³ as well as Ryan and Suhl.¹⁴ We do not include EHTv3 and EHTv4 in pqm4.
- HAETAE [CCD⁺23b]: HAETAE was merged into pqm4 in #313. We include both the reference implementation¹⁵(v2.0) and the M4F-optimized version described in [CCD⁺23a]. However, these implementations are not compatible with the original version submitted to NIST (v1.0). Small increases

⁹ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/KvIege2EbuM/m/oPrvAPLaBQAJ

¹⁰ https://gitlab.lrz.de/tueisec/fuleeca-signature/-/commit/3fc5835ea2e833efd01830944137a8dc0f4d0e58

¹¹ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/zas5PLiBe6A/m/A2KSHtqUAgAJ

¹² https://github.com/mti/attack_eaglesign

¹³ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/mFl_5Rq6-RU/m/2511f9lLAAAJ

¹⁴ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/bkJKBFq3TDY/m/lTCum6zgBQAJ

¹⁵ https://kpqc.cryptolab.co.kr/haetae

in the signature size, resulting from bug fixes and differences in the encoding and packing to address implementation vulnerabilities demonstrated by Saarinen¹⁶¹⁷ make this version incompatible with the original one. All three parameter sets are able to run on our evaluation platform.

- HAWK [BBD⁺23]: The reference implementations of HAWK from the submission package have been merged into pqm4 in #305 and all three parameter sets are running.
- HuFu [YJL⁺23]: HuFu uses public keys of 1059 KB (security level 1) to 3573 KB (security level 3). This exceeds the available memory available on our target platform and we, thus, do not include HuFu in pqm4.
- Raccoon [dPEK⁺23]: The reference implementation of Raccoon is making use of the __int128 datatype which is not portable to our platform. The Raccoon reference implementation offers the option (POLYR_Q32) to switch to 32-bit NTTs (mod 16515073 and mod 33292289) rather than 64-bit NTTs (mod 16515073 · 33292289) for better support of 32-bit platforms. However, turning on this option does not eliminate all instances of __int128. We have contacted the submission team and learned that a fully portable implementation is work in progress. We cannot include Raccoon in pqm4 at this moment.
- SQUIRRELS [ENST23]: SQUIRRELS requires public keys of 666 KB (Squirrels-I) to 2721 KB (Squirrels-V) which is too large for our target platform. Furthermore, the reference implementation of SQUIRRELS depends on multiple external libraries. We do not include SQUIRRELS in pqm4.

3.4 MPC-in-the-Head Signatures

- Biscuit [BKPV23]: The reference implementation of Biscuit from the submission package has been merged into pqm4 in #314. Of the six parameter sets, three (biscuit128f, biscuit192f, biscuit256f) meet the constraints of our evaluation board.
- MIRA [ABB+23c]: The MIRA reference implementation makes heavy use of dynamic memory that would have to be eliminated prior to merging it into pqm4. MIRA is re-using the same data structures (e.g., gfqm_vec or gf16_mat) for multiple sizes. This requires significant refactoring for eliminating dynamic memory allocations. We do not include MIRA in pqm4 for now.
- MiRitH [ARZV⁺23]: The MiRitH reference implementation from the official repository¹⁸ has been added to pqm4 in #315. The MiRitH team also provides an implementation optimized for the Cortex-M4 in the same repository. We have merged the optimized implementation in #325. We have reported multiple small issues with those implementations to the submission team which have been fixed in the official repository by now. All fast parameter sets of the non-hypercube variant are functional on our testing platform. From the hypercube parameter sets, the fast and short sets for security level 1 and 3 and mirith_hypercube_Va_fast and mirith_hypercube_Vb_fast are running.
- MQOM [FR23]: The MQOM reference implementation from the submission package has been merged into pqm4 in #322. We have eliminated a large number of dynamic memory allocations from the reference implementation. Luckily, the vast majority of dynamic memory allocations actually had a static size and could easily be replaced. A small number of variable-sized buffers have been replaced by variable-length arrays (VLAs). Those could be replaced with static buffers of worst-case length. Only the two parameter sets mqom_cat1_gf251_fast and mqom_cat1_gf31_fast of the 12 available sets are suitable for our evaluation board.
- PERK [ABB⁺23a]: The PERK team has contributed a reference implementation and a M4-optimized implementation compatible with the specification (v1.1) in #318. The M4 implementation is presented in [BBB⁺24]. This version is, however, incompatible with the one in the NIST submission package. Also note that compared to the official reference implementation (v1.1), the PERK team has replaced the GMP dependency with standalone arithmetic. All parameter sets are running on our evaluation board.
- RYDE [ABB⁺23b]: Similar as MIRA, RYDE uses numerous (>50) dynamic memory allocations re-using the same data-structures for differently sized buffers. Including RYDE into pqm4 would require significant refactoring effort for removing dynamic memory allocations. We do not include it for now.
- SDitH [MFG⁺23]: The SDitH reference implementation makes light use of dynamic memory allocations, but those could be eliminated without too much effort. However, the overall memory footprint ranges from

¹⁶ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/ImcSqGLFdoo/m/G86jtgDtBQAJ

¹⁷ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/Hq-wRFDbIaU/m/iLZctTiLAgAJ

¹⁸ https://github.com/Crypto-TII/mirith_nist_submission/commit/f27b540b77215dd17b10417726c6c6f7ccd41aa5

800 KB (for SDitH-L1-gf256) to 2.5 MB (for SDitH-L5-gf256) which exceeds the available memory of our target platform. Getting SDitH to work on the Cortex-M4 would require a stack-optimized implementation. We do not include SDitH in pqm4 for now.

3.5 Multivariate Signatures

- 3WISE [Rod23a]: Smith-Tone [ST23] presented a polynomial time attack on 3WISE, therefore, we did not include 3WISE in pqm4. Furthermore, the dependency on the FLINT library currently prevents the reference code from being included in pqm4.
- DME-Sign [LA23]: Briaud, Bros, Perlner, and Smith-Tone¹⁹ presented a key recovery attack on DME-Sign, which has been acknowledged by the DME-Sign team. Thus, we currently do no consider adding DME-Sign to pqm4.
- HPPC [Rod23b]: The reference implementation of HPPC unfortunately depends on the external libraries FLINT and M4RI. These external dependencies prohibit a straightforward integration into pqm4.
- MAYO [BCC⁺23b]: The MAYO reference implementation from the official repository²⁰ as well as the M4F-optimized implementation described in [BCC⁺23c] have been merged into pqm4 in #302. Note that [BCC⁺23c] also proposes a change to the MAYO specification by switching to a nibble representation rather than the bitsliced representation. We merged the bitsliced version that is compatible with the round-1 specification. Only the parameter set for the highest security level does not fit on the evaluation board.
- PROV [GCF⁺23]: Even for the smallest parameter set (PROV-I), PROV requires more than the available 640 KB for generating a signature. The current implementation requires 428 536 bytes for the expanded public key which together with the compressed public key (68 326 bytes) and the secret key (203 688 bytes) already exceeds the available memory.
- QR-UOV [FIH⁺23]: QR-UOV has not been included in pqm4 because the reference implementation allocates huge arrays for signature computation that exceed the available memory resources of our target platform.
- SNOVA [WCD⁺23]: We merged the SNOVA reference implementation from the submission package into pqm4 in #311. The current implementation of SNOVA implementation is using a pre-computed static table S. In the reference implementation, this table is computed dynamically and cached. The computation is done outside of the core function and, thus, not reflected in the benchmarks. For a quick integration into pqm4, we re-compute this table at the beginning of key generation, signing, and verification, but maintain the static array. This results in fairer benchmarks while not structurally changing the code. A better solution would be to pre-compute the constants and placing them in the code. All three parameter sets targeting security level 1 run both in the esk and the ssk variant on the evaluation board. snova-37-8-16-4-ssk targeting security level three is also running and leading to seven out of 18 variants. [IA24] raised concerns about the current SNOVA parameter sets reaching the claimed security levels. The SNOVA team acknowledged these concerns and have proposed updated parameters²¹. However, as of now, no updated implementation is available.
- TUOV [DGG⁺23]: The high memory usage of the reference implementation from the submission package of TUOV prevents it from running on our evaluation platform. The parameter set with the lowest memory consumption (tuov-Ip) requires around 750 KB of RAM. The implementation is making use of dynamic memory allocations. However, these can be easily eliminated. We included TUOV in pqm4 in #327, but none of the parameter sets are functional on the target board.
- UOV [BCD⁺23]: The UOV implementation from the official repository as well as the M4F-optimized implementation described in [BCH⁺23] have been merged in #300. Note that only the uov-Ip parameter sets fit on our target platform. In the paper, the authors are also able to evaluate the uov-Is parameters on the same target by writing public and secret keys to flash memory. Writing to flash memory is not supported by the pqm4 framework. The three security level 1 parameter sets require less then the 640 KiB SRAM and are thus functional on our board, the remaining nine parameter sets require more than this.

¹⁹ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/aoXpl4TlNh4/m/Eal1YHw0BAAJ

²⁰ https://github.com/PQCMayo/MAYO-C

²¹ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/m11kg20sTyU/m/cLkGIDaiBAAJ

- VOX [PCF⁺23]: Furue and Ikematsu [FI23] raised concerns that the security of the current VOX parameters has been severely overestimated.²². The VOX team has acknowledged the attack and proposes new parameters in [MRPC⁺23]. However, at this time no reference implementations of the new parameter sets have been published and we, hence, do not include VOX in pqm4.

3.6 Symmetric-based Signatures

- AIMer [KCC⁺23]: The AIMer reference implementation from the submission package contains a large number of dynamic memory allocations. All of those can be easily converted into stack allocations. We eliminated the dynamic memory allocations and merged AIMer into pqm4 in #323. However, for some of the parameter sets, that these buffers exceed the 4 MiB stack memory available on qemu's mps2-an386 (or even the 8 MiB default stack size on Linux). We only include parameter sets that we managed to successfully test using qemu. aimer-l1-param1, aimer-l1-param2 and aimer-l3-param1 are running on our target platform, nine further parameter sets for AIMer are not.
- Ascon-Sign [SGJ+23]: We have merged the Ascon-Sign reference implementation from the submission package into pqm4 in #263. Since pqm4 does not support Ascon, we use the Ascon version shipped in the submission package. All parameter sets are running on our evaluation board.
- **FAEST** [**BBdSG**⁺23]: The FAEST reference implementation in the submission package uses a large number of dynamic memory allocations (>150). All of the dynamic memory allocations in faest_aes.c and vole.c can easily be converted into VLAs. They could also be converted to fixed-sized buffers if defining parameters statically. However, some other dynamic memory allocations are harder to eliminate: The structs vec_com_t, vec_com_rec_t, tree_t hold pointers to buffers that have varying size even for a single parameter set. Those would have to be duplicated for each size needed, or alternatively, the worst case size needs to be used potentially increasing the memory footprint. We do not include FAEST in pqm4.
- SPHINCS-alpha [YCZ23]: The SPHINCS-alpha reference implementation from the submission package has been merged into pqm4 in #312. The implementation used static memory that contains a large (280 KB for sphincs-a-shake-128f) lookup table that is computed during first use (i.e., key generation) and re-used throughout the computation including signing and verification. This table computation requires significant time (around 15 million clock cycles for sphincs-a-shake-128f). As this leads to unfair benchmarking results in signing and verification, we instead compute the table once in the beginning of each of key generation, signing, and verification. We also move the table to the stack. We were able to make 6 out of 24 variants functional for pqm4: 128f, 128s and 192f each in both the sha2 and the shake version.

3.7 Other Signatures

- ALTEQ [BDN⁺23]: The ALTEQ implementation available in the NIST submission package makes heavy use of dynamic memory allocations. Additionally, the memory footprint is too large to fit on our target platform (alteq-shortsig-I requires around 1 MB, alteq-balanced-I requires around 2 MB). We, thus, do not include ALTEQ in pqm4.
- eMLE-Sig 2.0 [LZ23]: eMLE-Sig 2.0 has been shown vulnerable by Tibouchi.²³ An implementation of the attack is available.²⁴ We do not include eMLE-Sig 2.0 in pqm4.
- KAZ-SIGN [AAC⁺23]: Bernstein demonstrated a signature forgery attack against KAZ-SIGN,²⁵ we therefore did not include the scheme in pqm4. The KAZ-SIGN team has published updated versions four times (with the latest iteration being KAZ-SIGN 1.4²⁶), each time being broken by Bernstein within one day.
- Preon [CCC⁺23]: The reference implementation of the smallest parameter set of Preon (Preon128A) currently requires around 200 MB of memory for signing. Additionally, it has more than 250 dynamic memory allocations. We do not include Preon in pqm4.

²² https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/icHfTrzkfw4/m/Zj7GrnjMAQAJ

²³ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/zas5PLiBe6A/m/aOnAlT6cAQAJ

²⁴ https://github.com/mti/attack_emle

²⁵ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/2ljDcgtawFw/m/61PiLt6WAgAJ

²⁶ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/jv72ZzYwAZQ/m/ayNzr7U1GQAJ

- Xifrat1-Sign.I [NP23]: Xifrat1-Sign.I was practically broken by Panny²⁷. We do not include it.

4 Results

In this section we summarize the benchmarking results at the time of writing. As the pqm4 framework is under constant development, the numbers may change over time. The largest change for now is the use of a new target platform, which produces different results, due to its different memory timings. This change is accompanied by a newer compiler version. At the time of writing, we use the version 13.2 of the GNU C Compiler toolchain provided by Arm²⁸. In the future, implementations for the schemes may be replaced by faster versions, or implementations with other goals (e.g., lower memory requirements) are added. The pqm4 GitHub repository contains continually updated listing of the results.

Table 2 presents the measured execution speed of each implementation in terms of CPU cycles in thousands (i.e., kilocycles). We measured ten executions per scheme and list the average value, with the exception of the dilithium, haetae, and hawk schemes, which were executed 100 times, due to their significant variance in execution time. The table lists the average cycle counts of all executions, with the difference to the minimum and maximum shown in the super- and subscript. The percentage of cycles spent in symmetric primitives is shown in parentheses. As reference, we included numbers for dilithium implementing the third round specification.

Table 3 presents the memory requirements of each implementation. Listed are the sizes of the text (i.e., compiled code), data, and BSS sections produced by the implementations source, as well as the required stacksize (measured in KiB) of each operation, excluding the key, message, and signature. While the code/data sizes can be determined statically with the compiler tools, the stack size was determined using the QEMU simulator. As the QEMU simulator produces the same results for the memory metrics as real hardware platforms and provides more resources, it allows us to test more schemes and security levels. Some of the largest schemes, however, are still too big for the simulated platform.

Scheme	impl.	keygen		sign		verify	
dilithium2	clean	$1874 \stackrel{+41}{_{-35}}$	(62%)	$7283 \begin{array}{c} +13672 \\ -3962 \end{array}$	(37%)	$2062 \ ^{+0}_{-0}$	(53%)
	m4f	$1426 \ ^{+40}_{-47}$	(80%)	$3815 \ _{-2001}^{+7908}$	(67%)	$1417 \ ^{+0}_{-0}$	(77%)
dilithium3	clean	$3205 \ ^{+2}_{-1}$	(65%)	$12893 \begin{array}{c} +52247 \\ -7796 \end{array}$	(40%)	$3376 \ ^{+0}_{-0}$	(57%)
	m4f	$2516 \ ^{+1}_{-1}$	(82%)	$6374 \ ^{+11353}_{-3439}$	(69%)	$2411 \ _{-0}^{+0}$	(79%)
dilithium5	clean	$5340 \stackrel{+66}{_{-53}}$	(67%)	$15533 \begin{array}{c} +35954 \\ -7581 \end{array}$	(45%)	$5610 \ ^{+0}_{-0}$	(61%)
	m4f	$4277 {}^{+41}_{-46}$	(84%)	$8473 \ _{-3591}^{+16493}$	(74%)	$4185 \ ^{+0}_{-0}$	(82%)
haetae2	ref	$9265 \begin{array}{c} +49825 \\ -7549 \end{array}$	(25%)	$32068 \stackrel{+15301}{_{-25792}}$	8 (43%)	$1154 \ ^{+450}_{-50}$	(45%)
	m4f	$9184 \ ^{+34372}_{-7629}$	(27%)	$26104 \ {}^{+95950}_{-21385}$	(57%)	918 $^{+0}_{-0}$	(54%)
haetae3	ref	$17553 \begin{array}{c} +59078 \\ -14530 \end{array}$	(30%)	$44320 \begin{array}{c} +11618 \\ -34537 \end{array}$	3 (43%)	$2097 \stackrel{+890}{_{-99}}$	(50%)
	m4f	$14630 \ ^{+63266}_{-11877}$	(33%)	$30588 \stackrel{+15933}{_{-23135}}$	4 (57%)	$1761 \ ^{+0}_{-0}$	(57%)

Table 2: Average execution speed for key generation, signature generation, and signature verification for each scheme implementation, as measured on the Nucleo-L4R5ZI evaluation board. Execution speed is shown in thousands of cycles, with the difference to the minimum and maximum shown in the super- and subscript respectively. Cycles spent on symmetric cryptography shown in parentheses.

²⁷ https://groups.google.com/a/list.nist.gov/g/pqc-forum/c/9FXtBZKWueA/m/DojbRt6ZAgAJ ²⁸ https://developer.arm.com/Tools%20and%20Software/GNU%20Toolchain

Scheme	impl.	keygen		sign		verify	
haetae5	ref	$19940 \ ^{+84658}_{-16076}$	(31%)	$55087 \stackrel{+20754}{-4309}$	$^{42}_{7}(44\%)$	$2593 \ ^{+1186}_{-132}$	(54%)
	m4f	$19447 \ ^{+92871}_{-15916}$	(34%)	$42365 \begin{array}{c} +16212 \\ -33103 \end{array}$	$^{29}_{3}$ (57%)	$2324 \ ^{+0}_{-0}$	(58%)
hawk256	ref	$16846 \stackrel{+22553}{_{-5306}}$	(51%)	$1116 \ ^{+1848}_{-161}$	(62%)	$628 \ ^{+0}_{-0}$	(11%)
hawk512	ref	$53382 \begin{array}{c} +48360 \\ -8110 \end{array}$	(12%)	$1972 \stackrel{+0}{_{-0}}$	(49%)	$1294 \ ^{+1}_{-1}$	(9%)
hawk1024	ref	$231721 \begin{array}{c} +308769 \\ -47959 \end{array}$	5 (5%)	$4310 \ ^{+0}_{-0}$	(49%)	$2782 \stackrel{+2}{_{-2}}$	(8%)
biscuit128f	ref	$1055 \ ^{+0}_{-0}$	(54%)	$274072 \stackrel{+0}{_{-0}}$	(10%)	254371 $^{+0}_{-0}$	(9%)
biscuit192f	ref	$1886 \ ^{+0}_{-0}$	(54%)	765314 $^{+0}_{-0}$	(6%)	713413 $^{+0}_{-0}$	(6%)
biscuit256f	ref	$3302 \stackrel{+0}{_{-0}}$	(54%)	$1747188 \stackrel{+0}{_{-0}}$	(4%)	$1678999 \stackrel{+0}{_{-0}}$	(4%)
mayo1	ref	$7977 \stackrel{+0}{_{-0}}$	(39%)	$18005 \stackrel{+0}{_{-0}}$	(18%)	$6294 \ ^{+0}_{-0}$	(50%)
	m4f	$5242 \begin{array}{c} +0 \\ -0 \end{array}$	(60%)	9101 $^{+0}_{-0}$	(35%)	$4953 \ _{-0}^{+0}$	(63%)
mayo2	ref	$18433 \stackrel{+0}{_{-0}}$	(23%)	23547 $^{+0}_{-0}$	(18%)	5494 $^{+0}_{-0}$	(76%)
	m4f	$11918 \ _{-0}^{+0}$	(35%)	$11980 \ ^{+0}_{-0}$	(36%)	5130 $^{+0}_{-0}$	(81%)
mayo3	m4f	$18947 \stackrel{+0}{_{-0}}$	(55%)	$32477 \stackrel{+0}{_{-0}}$	(33%)	$16853 \ ^{+0}_{-0}$	(62%)
ov-Ip	ref	$350784 \stackrel{+0}{_{-0}}$	(3%)	$6479 \ ^{+0}_{-0}$	(0%)	$1301 \ ^{+0}_{-0}$	(1%)
	m4f	139186 $^{+0}_{-0}$	(8%)	$2705 \ ^{+2157}_{-240}$	(1%)	994 $^{+3}_{-15}$	(1%)
ov-Ip-pkc	ref	$375130 \ ^{+0}_{-0}$	(3%)	$6924 \ ^{+0}_{-0}$	(0%)	$11430 \ ^{+0}_{-0}$	(80%)
	m4fspeed	$175417 \ ^{+0}_{-0}$	(6%)	$2484 \ ^{+0}_{-0}$	(1%)	$11200 \ ^{+3}_{-16}$	(82%)
	m4fstack	$175417 \ _{-0}^{+0}$	(6%)	$2484 \ ^{+0}_{-0}$	(1%)	$12043 \ ^{+6}_{-21}$	(81%)
ov-Ip-pkc-skc	ref	$375130 \ ^{+0}_{-0}$	(3%)	$241521 \stackrel{+0}{_{-0}}$	(5%)	$12161 \ ^{+0}_{-0}$	(81%)
	m4fspeed	$175417 \ ^{+0}_{-0}$	(6%)	$89193 \ ^{+0}_{-0}$	(13%)	11987 $^{+4}_{-13}$	(82%)
	m4fstack	175417 $^{+0}_{-0}$	(6%)	$89193 \ _{-0}^{+0}$	(13%)	$12037 \stackrel{+12}{_{-33}}$	(82%)
snova-24-5-16-4-esk	ref	$24841 \stackrel{+1}{_{-1}}$	(12%)	139248 $^{+9}_{-13}$	(0%)	$88454 \ ^{+1}_{-1}$	(3%)
snova-24-5-16-4-ssk	ref	$24772 \ ^{+1}_{-1}$	(12%)	$174091 \stackrel{+12510}{_{-13926}}$	$_{5}^{6}$ (2%)	$88454 \ ^{+1}_{-1}$	(3%)
snova-25-8-16-3-esk	ref	$35281 \stackrel{+0}{_{-0}}$	(10%)	$65183 \ ^{+10}_{-10}$	(0%)	$42543 \ ^{+0}_{-0}$	(8%)
snova-25-8-16-3-ssk	ref	$35195 \ ^{+0}_{-0}$	(10%)	93185 $^{+13}_{-11}$	(4%)	$42543 \ _{-0}^{+0}$	(8%)
snova-28-17-16-2-esk	ref	51178 $^{+0}_{-0}$	(11%)	$21283 \ ^{+11}_{-15}$	(0%)	$19180 \ _{-0}^{+0}$	(28%)
snova-28-17-16-2-ssk	ref	$51065 \ ^{+0}_{-0}$	(11%)	$50759 \ ^{+9}_{-13}$	(11%)	19180 $^{+0}_{-0}$	(28%)
snova-37-8-16-4-ssk	ref	$122024 \stackrel{+1}{_{-0}}$	(9%)	$576130 \begin{array}{c} +38757 \\ -43107 \end{array}$	7^{9} (2%)	335807 $^{+1}_{-0}$	(3%)
cross-sha2-r-sdp-1-fast	ref	$5615 \begin{array}{c} +36 \\ -25 \end{array}$	(90%)	$216566 \ ^{+830}_{-822}$	(86%)	$142974 \ _{-1350}^{+782}$	(91%)
cross-sha2-r-sdp-3-fast	ref	$8201 \ ^{+7}_{-53}$	(87%)	$241882 \ _{-65}^{+166}$	(80%)	$123737 \ ^{+585}_{-782}$	(82%)
cross-sha2-r-sdpg-1-fast	ref	$2151 \stackrel{+24}{_{-35}}$	(93%)	$116163 \ _{-43}^{+20}$	(91%)	$87579 \begin{array}{c} +373 \\ -365 \end{array}$	(92%)
cross-sha2-r-sdpg-1-small	ref	$2151 \stackrel{+24}{_{-35}}$	(93%)	$391735 \ ^{+39}_{-69}$	(91%)	$368474 \begin{array}{c} +379 \\ -658 \end{array}$	(93%)

Table 2: Average execution speed for key generation, signature generation, and signature verification for each scheme implementation (cont.)

Scheme	impl.	keygen		sign		verify	
cross-sha2-r-sdpg-3-fast	ref	$3049 \ ^{+6}_{-52}$	(90%)	$136249 \ ^{+47}_{-13}$	(87%)	$84235 \ _{-266}^{+426}$	(87%)
cross-sha2-r-sdpg-5-fast	ref	$4302 {}^{+6}_{-52}$	(88%)	$221766 \ ^{+56}_{-65}$	(83%)	$134737 \ ^{+422}_{-591}$	(82%)
cross-sha3-r-sdp-1-fast	ref	968 $^{+11}_{-4}$	(67%)	$58864 \ ^{+13}_{-11}$	(61%)	$30641 \stackrel{+15}{_{-11}}$	(72%)
cross-sha3-r-sdp-3-fast	ref	$2195 \ ^{+10}_{-4}$	(68%)	90063 $^{+18}_{-32}$	(58%)	$48560 \stackrel{+9}{_{-24}}$	(62%)
cross-sha3-r-sdpg-1-fast	ref	$290 \ ^{+8}_{-2}$	(72%)	29964 $^{+10}_{-4}$	(75%)	$20095 \stackrel{+7}{_{-8}}$	(77%)
cross-sha3-r-sdpg-1-small	ref	$290 \ ^{+8}_{-2}$	(72%)	$102854 \begin{array}{c} +9 \\ -6 \end{array}$	(75%)	75138 $^{+16}_{-8}$	(78%)
cross-sha3-r-sdpg-3-fast	ref	$628 \ ^{+10}_{-2}$	(72%)	$43573 \begin{array}{c} +11 \\ -8 \end{array}$	(68%)	27513 $^{+10}_{-7}$	(69%)
cross-sha3-r-sdpg-5-fast	ref	$1146 \ ^{+8}_{-4}$	(71%)	93558 $^{+6}_{-9}$	(66%)	$59963 \begin{array}{c} +81 \\ -85 \end{array}$	(67%)
meds13220	ref	$47801 \stackrel{+4}{_{-4}}$	(2%)	$1773022 \stackrel{+26}{_{-23}}$	(5%)	$1766410 \ ^{+66}_{-58}$	(5%)
meds55604	ref	$253604 \stackrel{+8}{_{-18}}$	(2%)	$8009980 \stackrel{+43}{_{-115}}$	(3%)	$8320807 \stackrel{+52}{_{-96}}$	(3%)
aimer-l1-param1	ref	$393 \ ^{+0}_{-0}$	(72%)	$32386 \stackrel{+0}{_{-0}}$	(47%)	$31112 \ ^{+0}_{-0}$	(47%)
aimer-l1-param2	ref	$393 \ ^{+0}_{-0}$	(72%)	79451 $^{+1}_{-0}$	(49%)	$78428 \ ^{+7}_{-6}$	(49%)
aimer-13-param1	ref	981 $^{+0}_{-0}$	(77%)	90954 $^{+0}_{-0}$	(42%)	$88351 \ _{-0}^{+0}$	(42%)
mqom_cat1_gf251_fast	ref	7790 $^{+6}_{-10}$	(74%)	$149074 \stackrel{+10}{_{-6}}$	(33%)	$136748 \ _{-57}^{+44}$	(33%)
mqom_cat1_gf31_fast	ref	$8473 \ ^{+1}_{-1}$	(65%)	$243805 \begin{array}{c} +31 \\ -15 \end{array}$	(20%)	$244375 \ _{-53}^{+71}$	(19%)
mirith_Ia_fast	ref	$1304 \ ^{+0}_{-0}$	(61%)	$296733 \ _{-2}^{+2}$	(10%)	$276068 \ _{-22}^{+21}$	(10%)
mirith_Ib_fast	ref	$2515 \ ^{+0}_{-0}$	(61%)	565780 $^{+4}_{-2}$	(7%)	$528405 \begin{array}{c} +5 \\ -4 \end{array}$	(6%)
mirith_IIIa_fast	ref	$3009 \ ^{+0}_{-0}$	(64%)	$891195 \stackrel{+3}{_{-2}}$	(7%)	$831720 \ _{-15}^{+22}$	(7%)
mirith_IIIb_fast	ref	$4565 \ ^{+0}_{-0}$	(65%)	$1298812 \stackrel{+2}{_{-3}}$	(5%)	$1214256 \ _{-22}^{+14}$	(5%)
mirith_Va_fast	ref	$6255 \ ^{+0}_{-0}$	(65%)	2373351 $^{+6}_{-6}$	(4%)	$2233948 \ ^{+47}_{-55}$	(4%)
mirith_Vb_fast	ref	$8808 \stackrel{+0}{_{-0}}$	(65%)	$3406389 \ _{-4}^{+4}$	(3%)	$3205121 \ ^{+46}_{-67}$	(3%)
mirith_hypercube_Ia_fast	ref	$1304 \ ^{+0}_{-0}$	(61%)	$116967 \stackrel{+35}{_{-16}}$	(27%)	$111503 \stackrel{+14}{_{-30}}$	(24%)
	opt	996 $^{+0}_{-0}$	(78%)	$58998 \stackrel{+25}{_{-17}}$	(52%)	53603 $^{+17}_{-21}$	(49%)
mirith_hypercube_Ia_short	ref	$1304 \ ^{+0}_{-0}$	(61%)	$371003 \stackrel{+20}{_{-3}}$	(67%)	$364003 \ _{-5}^{+19}$	(67%)
mirith_hypercube_Ib_fast	ref	$2515 \ _{-0}^{+0}$	(61%)	$205513 \ _{-25}^{+25}$	(19%)	199773 $^{+5}_{-7}$	(17%)
	opt	$1880 \stackrel{+0}{_{-0}}$	(79%)	$83818 \begin{array}{c} +36 \\ -19 \end{array}$	(45%)	$78142 \begin{array}{c} +9 \\ -7 \end{array}$	(42%)
mirith_hypercube_Ib_short	ref	$2515 \ ^{+0}_{-0}$	(61%)	$447353 \begin{array}{c} +8 \\ -1 \end{array}$	(57%)	$440282 \begin{array}{c} +21 \\ -3 \end{array}$	(57%)
mirith_hypercube_IIIa_fast	ref	$3009 \ ^{+0}_{-0}$	(64%)	$323231 \stackrel{+58}{_{-42}}$	(19%)	$313484 \stackrel{+16}{_{-22}}$	(17%)
mirith_hypercube_IIIa_short	ref	$3009 \ ^{+0}_{-0}$	(64%)	718710 $^{+15}_{-3}$	(55%)	730959 $^{+6}_{-9}$	(56%)
mirith_hypercube_IIIb_fast	ref	$4565 \ ^{+0}_{-0}$	(65%)	$450968 \stackrel{+24}{_{-36}}$	(14%)	$441592 \ _{-9}^{+14}$	(12%)
mirith_hypercube_IIIb_short	ref	$4565 \ ^{+0}_{-0}$	(65%)	$845099 \stackrel{+17}{_{-4}}$	(47%)	$860322 \ _{-9}^{+13}$	(48%)
mirith_hypercube_Va_fast	ref	$6255 \ ^{+0}_{-0}$	(65%)	$808047 \stackrel{+140}{_{-99}}$	(13%)	$792074 \ _{-84}^{+59}$	(12%)

Table 2: Average execution speed for key generation, signature generation, and signature verification for each scheme implementation (cont.)

Scheme	impl.	keygen		sign		, , , , , , , , , , , , , , , , , , ,	verify	
<pre>mirith_hypercube_Vb_fast</pre>	ref	$8808 \ ^{+0}_{-0}$	(65%)	$1134255 \begin{array}{c} +113 \\ -75 \end{array}$	(11%)	1117091	$^{+19}_{-34}$	(10%)
perk-128-fast-3	ref	$698 \ ^{+7}_{-4}$	(59%)	$217643 \ _{-309}^{+409}$	(28%)	96371	$^{+1308}_{-400}$	(62%)
	m4	$595 \ ^{+4}_{-7}$	(69%)	$175927 \ _{-4014}^{+4141}$	(67%)	81048	$^{+511}_{-272}$	(72%)
perk-128-fast-5	ref	911 $^{+43}_{-2}$	(51%)	$215520 \ _{-481}^{+719}$	(27%)	93954	$^{+1388}_{-483}$	(60%)
	m4	733 $^{+28}_{-2}$	(63%)	$169043 \ _{-4737}^{+6630}$	(66%)	78118	$^{+509}_{-438}$	(71%)
perk-128-short-3	m4	$595 \ ^{+4}_{-7}$	(69%)	$976147 \ _{-26439}^{+41536}$	(64%)	477798	$^{+18823}_{-19046}$	(65%)
perk-128-short-5	m4	733 $^{+28}_{-2}$	(63%)	$902935 \begin{array}{c} ^{+44104}_{-39106}$	(63%)	445552	$^{+18846}_{-19329}$	(64%)
perk-192-fast-3	m4	$1422 {}^{+43}_{-5}$	(73%)	$420038 \ ^{+7735}_{-7607}$	(68%)	194755	$^{+1390}_{-1202}$	(72%)
perk-192-fast-5	m4	$1679 \ ^{+46}_{-14}$	(68%)	$400224 \begin{array}{c} +9757 \\ -11504 \end{array}$	(67%)	187024	$^{+993}_{-768}$	(71%)
perk-192-short-3	m4	$1422 \begin{array}{c} +43 \\ -5 \end{array}$	(73%)	$2405700 \begin{array}{c} +92909 \\ -108003 \end{array}$	(63%)	1203963	$^{+29542}_{-29447}$	(63%)
perk-192-short-5	m4	$1686 \stackrel{+39}{_{-22}}$	(68%)	$2222672 \begin{array}{c} +90220 \\ -70349 \end{array}$	(62%)	1120718	$^{+23752}_{-24792}$	(62%)
perk-256-fast-3	m4	$2482 \stackrel{+50}{_{-9}}$	(74%)	$872567 \ ^{+15040}_{-14028}$	(70%)	417918	$^{+2037}_{-2912}$	(75%)
perk-256-fast-5	m4	$2888 {+52 \atop -19}$	(68%)	$834202 \begin{array}{c} ^{+18719}_{-15315}$	(70%)	398767	$^{+1968}_{-1190}$	(74%)
perk-256-short-3	m4	$2482 \stackrel{+50}{_{-9}}$	(74%)	$5076941 \ {}^{+110828}_{-61663}$	(65%)	2650317	$^{+15593}_{-20683}$	(65%)
perk-256-short-5	m4	$2882 \stackrel{+57}{_{-3}}$	(68%)	$4682541 \begin{array}{c} ^{+89928}_{-119145}$	(64%)	2454136	$^{+8774}_{-8640}$	(64%)
ascon-sign-128f-robust	ref	$122506 \stackrel{+0}{_{-0}}$	(0%)	$2855798 \ _{-0}^{+1}$	(0%)	177864	$^{+4568}_{-3738}$	(0%)
ascon-sign-128f-simple	ref	$69377 \ _{-0}^{+0}$	(0%)	$1629111 \ ^{+0}_{-0}$	(0%)	96768	$^{+5151}_{-3091}$	(0%)
ascon-sign-128s-robust	ref	$7842367 \ _{-0}^{+0}$	(0%)	$59267553 \ ^{+14}_{-16}$	(0%)	61063	$^{+2616}_{-4028}$	(0%)
ascon-sign-128s-simple	ref	$4441129 \ ^{+0}_{-0}$	(0%)	$33877716 \ ^{+16}_{-22}$	(0%)	34009	$^{+1442}_{-1534}$	(0%)
ascon-sign-192f-robust	ref	$222614 \stackrel{+0}{_{-0}}$	(0%)	5712961 $^{+3}_{-2}$	(0%)	320251	$^{+4393}_{-4910}$	(0%)
ascon-sign-192f-simple	ref	$128167 \stackrel{+0}{_{-0}}$	(0%)	$3345488 \stackrel{+7}{-4}$	(0%)	178458	$^{+10555}_{-5104}$	(0%)
ascon-sign-192s-robust	ref	$14249839 \stackrel{+1}{_{-2}}$	(0%)	$126646611 \ ^{+24}_{-19}$	(0%)	109281	$^{+6202}_{-7236}$	(0%)
ascon-sign-192s-simple	ref	$8204011 \stackrel{+0}{_{-0}}$	(0%)	$74760764 \ _{-38}^{+23}$	(0%)	61709	$^{+2204}_{-1276}$	(0%)
sphincs-a-sha2-128f	ref	$30279 \stackrel{+0}{_{-0}}$	(43%)	$382271 \stackrel{+1}{_{-1}}$	(87%)	35696	$^{+2}_{-1}$	(48%)
sphincs-a-sha2-128s	ref	$814837 \stackrel{+1}{_{-1}}$	(73%)	$6981930 \ _{-6}^{+5}$	(88%)	187091	$^{+2}_{-3}$	(12%)
sphincs-a-sha2-192f	ref	$45931 \stackrel{+0}{_{-0}}$	(58%)	$634374 \stackrel{+2}{_{-1}}$	(88%)	35146	$^{+1}_{-1}$	(46%)
sphincs-a-shake-128f	ref	$61578 \stackrel{+0}{_{-0}}$	(72%)	1188147 $^{+1}_{-1}$	(96%)	76330	$^{+2}_{-2}$	(76%)
sphincs-a-shake-128s	ref	$2342299 \stackrel{+1}{_{-2}}$	(91%)	$22926755 \ _{-3}^{+4}$	(97%)	241835	$^{+2}_{-3}$	(32%)
sphincs-a-shake-192f	ref	$110028 \stackrel{+0}{_{-0}}$	(83%)	$1814954 \stackrel{+1}{_{-1}}$	(97%)	64022	$^{+1}_{-1}$	(71%)

Table 2: Average execution speed for key generation, signature generation, and signature verification for each scheme implementation (cont.)

			0			,	/
		libra	ry siz	e	stack usage		
Scheme	impl.	code	data	\mathbf{bss}	keygen	sign	verify
dilithium2	clean	7996	0	0	37.4	50.7	35.3
	m4f	18516	0	0	37.4	48.2	35.3
dilithium3	clean	7496	0	0	59.4	77.7	56.3
	m4f	20004	0	0	59.4	67.2	56.3
dilithium5	clean	7784	0	0	95.4	119.7	90.6
	m4f	18312	0	0	95.4	113.2	90.6
haetae2	ref	25568	0	0	25.5	53.1	29.0
	m4f	35708	0	0	19.3	54.3	22.8
haetae3	ref	25980	0	0	42.5	78.5	47.3
	m4f	35936	0	0	28.8	81.4	31.0
haetae5	ref	25688	0	0	53.5	97.8	60.6
	m4f	35692	0	0	33.3	101.4	36.3
hawk256	ref	102015	0	0	7.7	3.2	3.6
hawk512	ref	102027	0	0	14.2	4.7	6.1
hawk1024	ref	102031	0	0	27.2	7.8	11.2
biscuit128f	ref	7580	0	0	0.6	134.0	14.2
biscuit128s	ref	7696	0	0	0.6	1067.6	81.7
biscuit192f	ref	7780	0	0	0.6	259.6	20.2
biscuit192s	ref	7904	0	0	0.6	2193.0	104.7
biscuit256f	ref	8216	0	0	0.7	466.2	32.2
biscuit256s	ref	8248	0	0	0.7	3889.6	144.7
mayo1	ref	26436	8	0	72.7	213.5	390.0
	m4f	19300	8	0	72.7	110.8	430.3
mayo2	ref	24404	8	0	108.9	232.7	263.9
	m4f	17292	8	0	108.9	121.7	271.8
mayo3	ref	30828	8	0	239.0	699.2	1110.7
	m4f	23612	8	0	239.0	332.7	458.8
ov-Ip	ref	29215	0	0	15.2	12.4	6.0
	m4f	118939	0	0	15.2	5.1	2.5
ov-Ip-pkc	ref	29423	0	0	15.2	12.4	277.9
	m4fspeed	119131	0	0	138.8	5.1	274.4
	m4fstack	119059	0	0	138.8	5.1	6.3

Table 3: Memory requirements for each scheme implementation. Code, data and BSS size listed are in bytes, stack usage in 2^{10} byte (i.e., KiB).

Scheme	impl.	code	data	bss	keygen	sign	verify	
ov-Ip-pkc-skc	ref	29467	0	0	247.6	247.5	277.9	
	m4fspeed	119175	0	0	371.1	237.4	274.4	
	m4fstack	119103	0	0	371.1	237.4	6.3	
tuov_ip	ref	92760	0	0	3201.4	3517.0	1764.9	
tuov_ip_pkc	ref	88400	0	0	502.6	15.7	401.6	
tuov_ip_pkc_skc	ref	88448	0	0	736.4	800.5	401.6	
tuov_is	ref	49316	0	0	337.6	10.6	1.6	
tuov_is_pkc	ref	49504	0	0	740.1	10.6	595.6	
tuov_is_pkc_skc	ref	49548	0	0	1082.1	1179.2	595.6	
tuov_iii	ref	92576	0	0	981.0	24.0	5.8	
tuov_iii_pkc	ref	92712	0	0	2177.7	24.0	1764.9	
tuov_iii_pkc_skc	ref	92760	0	0	3201.4	3517.0	1764.9	
tuov_v_pkc	ref	86304	0	0	una	unable to test		
tuov_v_pkc_skc	ref	86352	0	0	una	est		
snova-24-5-16-4-esk	ref	52132	0	336	165.0	87.9	115.4	
snova-24-5-16-4-ssk	ref	52132	0	336	165.0	165.1	115.4	
snova-25-8-16-3-esk	ref	28240	0	299	186.2	85.4	119.8	
snova-25-8-16-3-ssk	ref	28240	0	299	186.2	186.3	119.8	
snova-28-17-16-2-esk	ref	16436	0	280	302.0	124.2	195.0	
snova-28-17-16-2-ssk	ref	16436	0	280	302.0	302.1	195.0	
snova-37-8-16-4-esk	ref	52440	0	336	625.3	289.1	401.8	
snova-37-8-16-4-ssk	ref	52440	0	336	625.3	625.4	401.8	
snova-43-25-16-2-esk	ref	15060	0	280	1015.1	407.9	650.6	
snova-43-25-16-2-ssk	ref	15060	0	280	1015.1	1015.2	650.6	
snova-49-11-16-3-esk	ref	26912	0	299	852.6	361.6	528.0	
snova-49-11-16-3-ssk	ref	26912	0	299	852.6	852.7	528.0	
snova-60-10-16-4-esk	ref	52588	0	336	1897.8	820.3	1179.7	
snova-60-10-16-4-ssk	ref	52588	0	336	1897.8	1897.9	1179.7	
snova-61-33-16-2-esk	ref	15020	0	280	2581.5	1027.9	1643.0	
snova-61-33-16-2-ssk	ref	15020	0	280	2581.5	2581.6	1643.0	
snova-66-15-16-3-esk	ref	27624	0	299	2117.3	876.1	1297.9	
snova-66-15-16-3-ssk	ref	27624	0	299	2117.3	2117.4	1297.9	

Table 3: Memory requirements for each scheme implemenation. (cont.)

Scheme	impl.	code	data	\mathbf{bss}	keygen	sign	verify
cross-sha2-r-sdp-1-fast	ref	14244	0	104	5.1	213.2	103.1
cross-sha2-r-sdp-1-small	ref	15285	0	104	5.1	691.1	314.9
cross-sha2-r-sdp-3-fast	ref	14744	0	128	9.7	317.2	154.4
cross-sha2-r-sdp-3-small	ref	14921	0	128	9.7	1238.9	584.9
cross-sha2-r-sdp-5-fast	ref	14580	0	152	16.4	839.3	401.1
cross-sha2-r-sdp-5-small	ref	14657	0	152	16.4	1661.8	784.1
cross-sha2-r-sdpg-1-fast	ref	18409	0	104	2.7	127.8	67.9
cross-sha2-r-sdpg-1-small	ref	18674	0	104	2.7	455.4	239.7
cross-sha2-r-sdpg-3-fast	ref	19609	0	128	4.4	200.3	105.9
cross-sha2-r-sdpg-3-small	ref	19994	0	128	4.4	737.1	383.9
cross-sha2-r-sdpg-5-fast	ref	18669	0	152	7.2	389.2	208.7
cross-sha2-r-sdpg-5-small	ref	18818	0	152	7.2	1004.5	521.6
cross-sha3-r-sdp-1-fast	ref	14472	0	208	4.6	213.3	103.1
cross-sha3-r-sdp-1-small	ref	15353	0	208	4.6	691.2	314.9
cross-sha3-r-sdp-3-fast	ref	14884	0	208	9.1	317.1	154.1
cross-sha3-r-sdp-3-small	ref	14825	0	208	9.1	1238.9	584.6
cross-sha3-r-sdp-5-fast	ref	14576	0	208	15.8	839.3	400.7
cross-sha3-r-sdp-5-small	ref	14629	0	208	15.8	1661.9	783.7
cross-sha3-r-sdpg-1-fast	ref	18605	0	208	2.3	127.9	67.9
cross-sha3-r-sdpg-1-small	ref	18846	0	208	2.3	455.5	239.8
cross-sha3-r-sdpg-3-fast	ref	19689	0	208	3.9	200.3	105.6
cross-sha3-r-sdpg-3-small	ref	19846	0	208	3.9	737.2	383.6
cross-sha3-r-sdpg-5-fast	ref	18593	0	208	6.7	389.3	208.3
cross-sha3-r-sdpg-5-small	ref	18762	0	208	6.7	1004.5	521.2
meds9923	ref	16720	0	0	36.8	973.8	98.1
meds13220	ref	16844	0	0	43.0	176.1	46.6
meds134180	ref	9180	0	0	392.5	853.4	375.7
meds167717	ref	9152	0	0	448.7	567.0	420.3
meds41711	ref	8948	0	0	135.6	1268.2	172.2
meds55604	ref	9012	0	0	158.3	383.6	156.5
aimer-11-param1	ref	19302	468	0	10.7	183.2	192.1
aimer-11-param2	ref	19894	468	0	10.7	432.2	441.1

Table 3: Memory requirements for each scheme implemenation. (cont.)

Scheme	impl.	code	data	\mathbf{bss}	keygen	sign	verify
aimer-l1-param3	ref	19658	468	0	10.7	1390.4	1399.3
aimer-13-param1	ref	23398	468	0	23.3	404.1	425.4
aimer-13-param2	ref	23570	468	0	23.3	1027.7	1049.0
aimer-15-param1	ref	28142	468	0	57.0	821.4	876.1
aimer-15-param2	ref	28490	468	0	57.0	2035.7	2090.4
mqom_cat1_gf251_fast	ref	16865	0	0	180.0	390.0	246.7
mqom_cat1_gf251_short	ref	18193	0	0	180.0	649.8	371.0
mqom_cat1_gf31_fast	ref	23718	0	0	256.7	598.5	411.6
mqom_cat1_gf31_short	ref	23510	0	0	256.7	847.8	541.1
mqom_cat3_gf251_fast	ref	17074	0	0	665.0	1254.4	803.1
mqom_cat3_gf251_short	ref	18710	0	0	665.0	1840.4	1116.9
mqom_cat3_gf31_fast	ref	23881	0	0	952.0	2099.4	1500.2
mqom_cat3_gf31_short	ref	24025	0	0	952.0	2606.9	1732.4
mqom_cat5_gf251_fast	ref	20401	0	0	1662.4	3146.7	2115.1
mqom_cat5_gf251_short	ref	20137	0	0	1662.4	4019.1	2477.9
mirith_Ia_fast	ref	7610	256	0	10.0	119.6	22.2
mirith_Ia_short	ref	7635	256	0	10.0	986.7	92.1
mirith_Ib_fast	ref	7602	256	0	18.7	147.5	32.9
mirith_Ib_short	ref	7675	256	0	18.7	1156.7	117.1
mirith_IIIa_fast	ref	7720	256	0	21.4	260.2	45.0
mirith_IIIa_short	ref	7889	256	0	21.4	2130.0	139.8
mirith_IIIb_fast	ref	7760	256	0	32.1	291.2	57.7
mirith_IIIb_short	ref	7901	256	0	32.1	2314.5	162.3
mirith_Va_fast	ref	7890	256	0	44.0	473.1	83.2
mirith_Va_short	ref	7987	256	0	44.0	3701.9	202.5
mirith_Vb_fast	ref	7914	256	0	61.4	522.6	103.5
mirith_Vb_short	ref	7991	256	0	61.4	3994.6	234.9
mirith_hypercube_Ia_fast	ref	8844	256	0	10.0	75.1	20.4
	opt	10932	0	0	10.0	75.1	20.4
mirith_hypercube_Ia_short	ref	8781	256	0	10.0	212.5	30.4
mirith_hypercube_Ia_shorter	ref	8843	256	0	10.0	1728.9	211.2

Table 3: Memory requirements for each scheme implementation. (cont.)

Scheme	impl.	code	data	\mathbf{bss}	keygen	sign	verify	
mirith_hypercube_Ib_fast	ref	8820	256	0	18.7	94.7	30.5	
	opt	10908	0	0	18.7	94.7	30.5	
mirith_hypercube_Ib_short	ref	8825	256	0	18.7	231.8	40.2	
mirith_hypercube_Ib_shorter	ref	8819	256	0	18.7	1748.7	221.0	
mirith_hypercube_IIIa_fast	ref	8966	256	0	21.4	163.1	41.6	
mirith_hypercube_IIIa_short	ref	9139	256	0	21.4	475.0	55.2	
mirith_hypercube_IIIa_shorter	ref	9109	256	0	21.4	3788.1	325.9	
mirith_hypercube_IIIb_fast	ref	9002	256	0	32.1	185.1	53.8	
mirith_hypercube_IIIb_short	ref	9159	256	0	32.1	497.1	67.0	
mirith_hypercube_IIIb_shorter	ref	9093	256	0	32.1	3809.8	337.6	
mirith_hypercube_Va_fast	ref	9140	256	0	44.0	301.2	79.2	
mirith_hypercube_Va_short	ref	9241	256	0	44.0	832.8	93.8	
	ref	9203	256	0	una	ble to t	est	
	ref	9160	256	0	61.4	336.8	98.8	
	ref	9245	256	0	61.4	868.6	112.8	
mirith_hypercube_Vb_shorter	ref	9247	256	0	una	unable to test		
perk-128-fast-3	ref	11053	4	0	7.5	306.0	305.8	
	m4	13421	4	0	7.5	23.5	20.2	
perk-128-fast-5	ref	11129	4	0	8.8	298.5	298.3	
	m4	13493	4	0	8.8	24.6	21.2	
perk-128-short-3	ref	31757	4	0	7.5	1524.5	1524.3	
	m4	26313	4	0	7.5	27.1	24.6	
perk-128-short-5	ref	31905	4	0	8.8	1428.5	1428.3	
	m4	26461	4	0	8.8	27.9	25.4	
perk-192-fast-3	ref	11377	4	0	14.6	670.5	670.2	
	m4	12253	4	0	14.6	46.6	40.4	
perk-192-fast-5	ref	11365	4	0	16.5	646.7	646.4	
	m4	12277	4	0	16.5	47.6	41.4	
perk-192-short-3	ref	31789	4	0	14.6	3388.5	3388.2	
	m4	24189	4	0	14.6	50.1	45.6	
perk-192-short-5	ref	31813	4	0	16.5	3148.5	3148.2	
	m4	24913	4	0	16.5	50.7	46.1	

Table 3: Memory requirements for each scheme implemenation. (cont.)

Scheme	impl.	code	data	\mathbf{bss}	keygen	sign	verify
perk-256-fast-3	ref	11313	4	0	24.9	1163.9	1163.7
	m4	12337	4	0	24.9	78.4	68.3
perk-256-fast-5	ref	11325	4	0	27.4	1114.5	1114.1
	m4	12337	4	0	27.4	79.0	69.0
perk-256-short-3	ref	31961	4	0	una	ble to t	est
	m4	31905	4	0	24.9	80.3	73.0
perk-256-short-5	ref	31973	4	0	una	ble to t	est
	m4	32989	4	0	27.4	80.1	73.1
ascon-sign-128f-robust	ref	17664	0	0	3.1	2.7	2.9
ascon-sign-128f-simple	ref	17596	0	0	2.6	2.1	2.4
ascon-sign-128s-robust	ref	17972	0	0	3.3	2.7	2.2
ascon-sign-128s-simple	ref	17904	0	0	2.8	2.2	1.7
ascon-sign-192f-robust	ref	17960	0	0	5.9	4.8	4.4
ascon-sign-192f-simple	ref	17912	0	0	4.7	3.6	3.6
ascon-sign-192s-robust	ref	18472	0	0	6.2	4.9	4.2
ascon-sign-192s-simple	ref	18424	0	0	5.0	3.7	3.0
sphincs-a-sha2-128f	ref	6343	0	0	274.6	274.1	274.1
sphincs-a-sha2-128s	ref	6899	0	0	572.2	571.9	571.6
sphincs-a-sha2-192f	ref	6867	0	0	492.3	490.7	490.6
sphincs-a-sha2-192s	ref	7423	0	0	1259.2	1258.4	1258.0
sphincs-a-sha2-256f	ref	7203	0	0	1046.7	1044.6	1044.3
sphincs-a-sha2-256s	ref	7511	0	0	2208.4	2207.1	2206.7
sphincs-a-shake-128f	ref	5552	0	0	274.5	274.0	274.0
sphincs-a-shake-128s	ref	6108	0	0	572.2	571.8	571.5
sphincs-a-shake-192f	ref	5916	0	0	491.7	490.2	490.1
sphincs-a-shake-192s	ref	6428	0	0	1258.6	1257.8	1257.5
sphincs-a-shake-256f	ref	6188	0	0	1046.2	1044.1	1043.8
sphincs-a-shake-256s	ref	6456	0	0	2207.9	2206.5	2206.2

Table 3: Memory requirements for each scheme implementation. (cont.)

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