

Post-Quantum Protocols for Banking Applications

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IDEMIA

Context

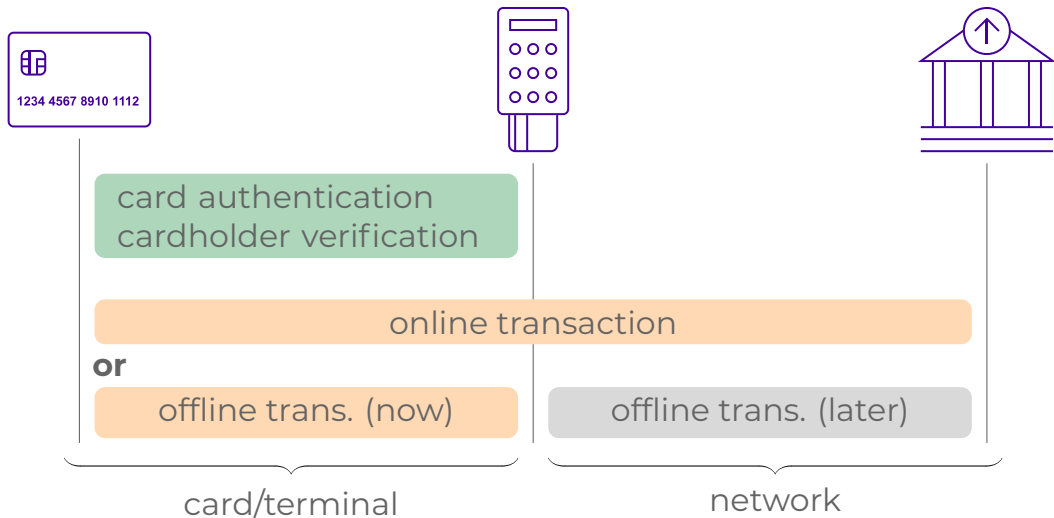
Payment Applications

- › Implemented in smartcards or smartphones
- › Contact or contactless payment
- › Typical actors: Cardholder+Card, Merchant+Terminal, Issuer+Server

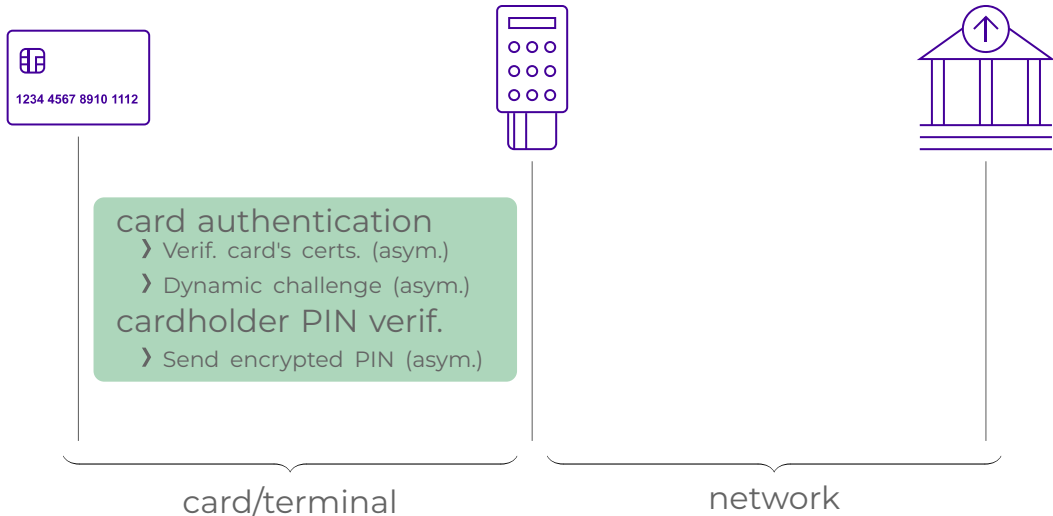
EMVCo: a “standard” for card-based payment

- › EMVCo: global technical body formed in 1999
- › Overseen by six member organizations – American Express, Discover, JCB, Mastercard, UnionPay and Visa
- › Manages and evolves **EMV Specifications**
- › More than 90% transactions worldwide

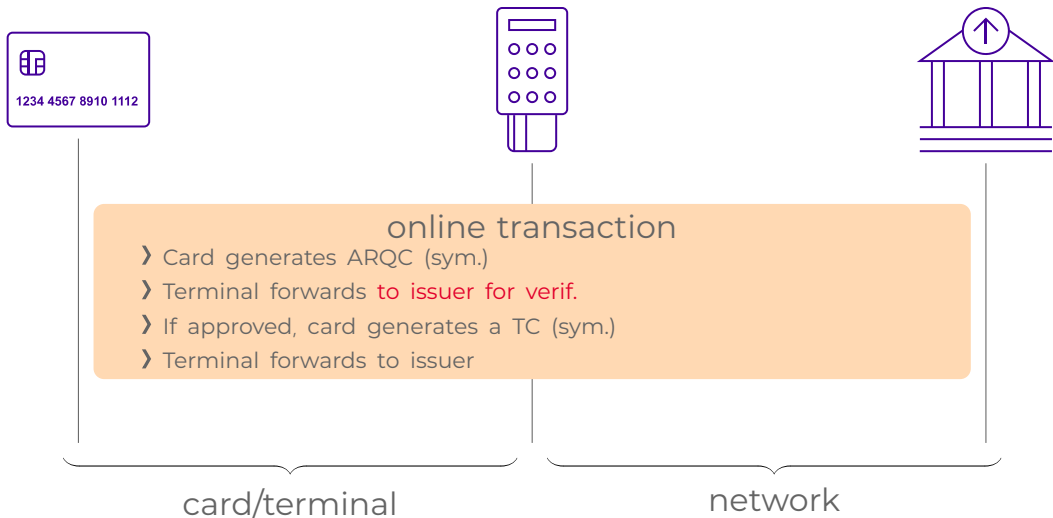
Card-Based Transactions



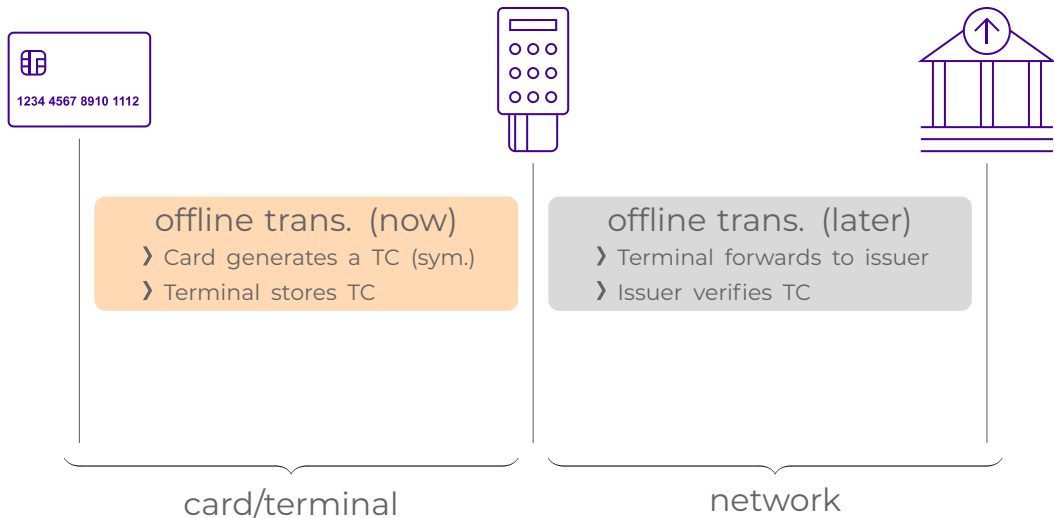
Common Part



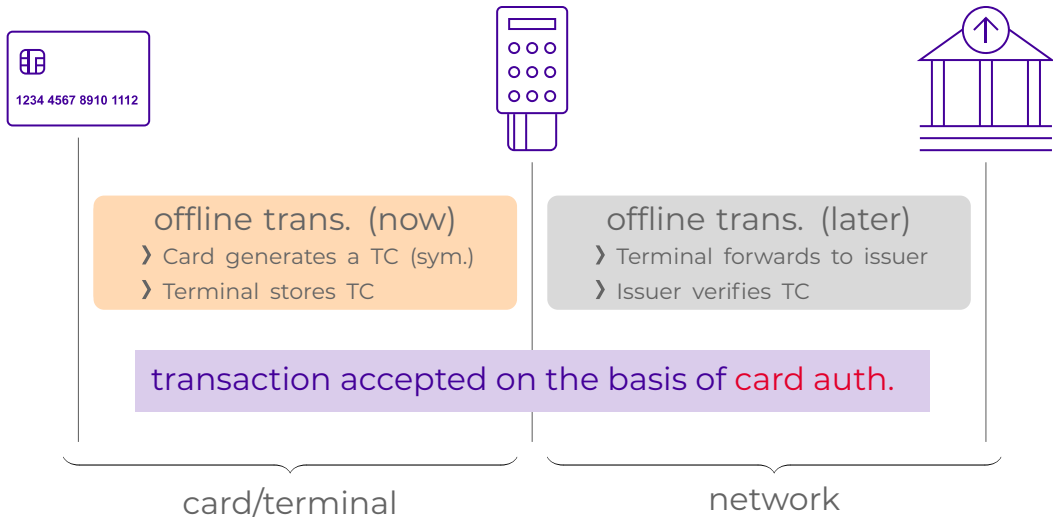
Online Transaction



Offline Transaction



Offline Transaction



Quantum Threat

All current asymmetric crypto. would be **broken by quantum computers**.

Online transactions

- › symmetric is quantum-safe (larger keys)
- › no real threat

Offline transactions

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- › if card authentication is broken \rightsquigarrow **too late!**

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Wait-and-see policy?

- › no “record now, harvest later” risk
- › we can still **force online** transactions!

Why bother now?

Offline transactions are crucial in certain cases

- › high volume/high speed payment scenarios
- › cost of bandwidth for small merchants
- › locations with low connectivity

Forbidding them would have a significant impact on the business

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Viable quantum-safe offline solution will be a long process:

- › Cryptographic migration in the payment industry is long and complex
- › HW changes could be required to keep fast transactions

start experiments now → identify obstacles → facilitate the effort

Outline

Contributions

- › Post-Quantum, efficient versions of current protocols
- › Implementation of hybrid versions using NIST PQ finalists
- › Performance analysis

Outline

1 › Existing Protocols

2 › Post-Quantum Migration

3 › Implementation Results

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Existing Protocols

Already Deployed (with RSA)

- › Static Data Authentication (SDA)
- › Dynamic Data Authentication (DDA)
- › Combined Data Authentication (CDA)

Privacy Preserving Protocol (with ECC)

- › Blinded Diffie-Hellman (BDH)

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In this presentation: we only detail CDA migration.

EMV CDA protocol (1/4)



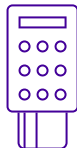
EMV CDA protocol (2/4)



EMV CDA protocol (3/4)



PIN, MK_{AC} , sk_C , $cert_C$, $cert_{IS}$



pk_{CA}

Get PIN from user: $uPIN$

$c \leftarrow \text{RSA.Enc}(pk_C, uPIN|nonce_C)$

VerifyPIN(c)

$uPIN|nonce \leftarrow \text{RSA.Dec}(sk_C, c)$

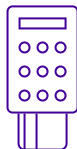
Check $uPIN$ and $nonce$

Success/Failed

EMV CDA protocol (4/4)



PIN, MK_{AC} , sk_C , $cert_C$, $cert_{IS}$



pk_{CA}

$nonce_T \leftarrow \text{random}()$

GenerateAC($nonce_T$)

$TC \leftarrow \text{MAC}(MK_{AC}, \text{transInfo})$

$data = nonce_C | TC | nonce_T$

$SDAD \leftarrow \text{RSA.Sign}(sk_C, data)$

$TC, SDAD$

$\text{RSA.Verify}(pk_C, SDAD)$

If ok, transaction is accepted and TC stored.

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Post-Quantum Migration

How to proceed?

- › Smooth transition \rightsquigarrow minimum amount of change in the transaction flow
- › The lazy way: replace by equivalent primitives (enc/sign)
- › The efficient way: use **key encapsulation mechanism**

Post-Quantum Migration

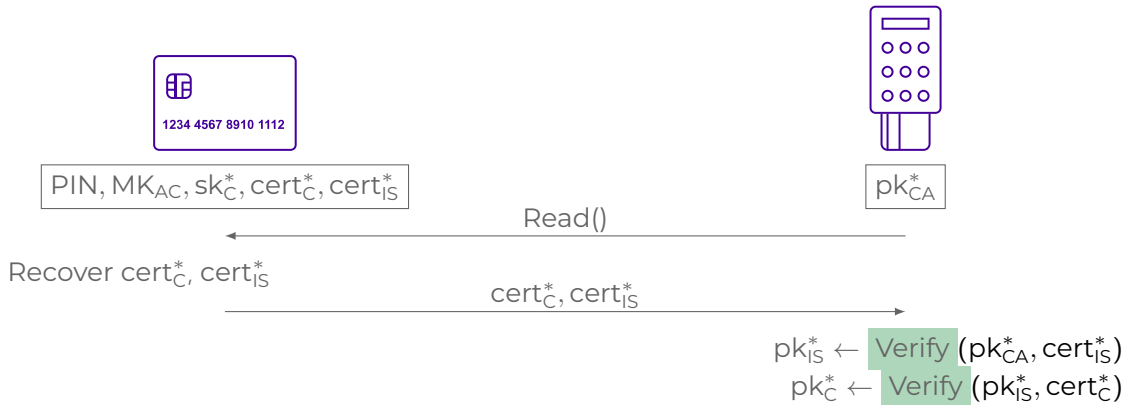
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Advantages

- › Key decaps. faster than signature with PQ algorithms
 - › Only **one asym.** operation to perform on card
 - › Ciphertexts *usually* smaller than signatures
-
- › Use **AES-256** for all symmetric primitives
 - › Use authenticated encryption for both confidentiality **and** authenticity

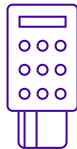
PQ analog of CDA protocol (1/4)



PQ analog of CDA protocol (2/4)



PIN, MK_{AC} , sk_C^* , $cert_C^*$, $cert_{IS}^*$



pk_{CA}^*

$ss, ct \leftarrow \text{Encaps}(pk_C^*)$

GetChallengeAndEstablishKey(ct)

$ss \leftarrow \text{Decaps}(sk_C^*, ct)$

$nonce_C \leftarrow \text{random}()$

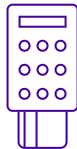
$K \leftarrow \text{KDeriv}(ss|nonce_C)$

$nonce_C$

PQ analog of CDA protocol (3/4)



PIN, MK_{AC} , sk_C^* , $cert_C^*$, $cert_{IS}^*$



pk_{CA}^*

$K \leftarrow KDeriv(ss|nonce_C)$
Get PIN from user: $uPIN$
 $c \leftarrow AEnc(K, uPIN)$

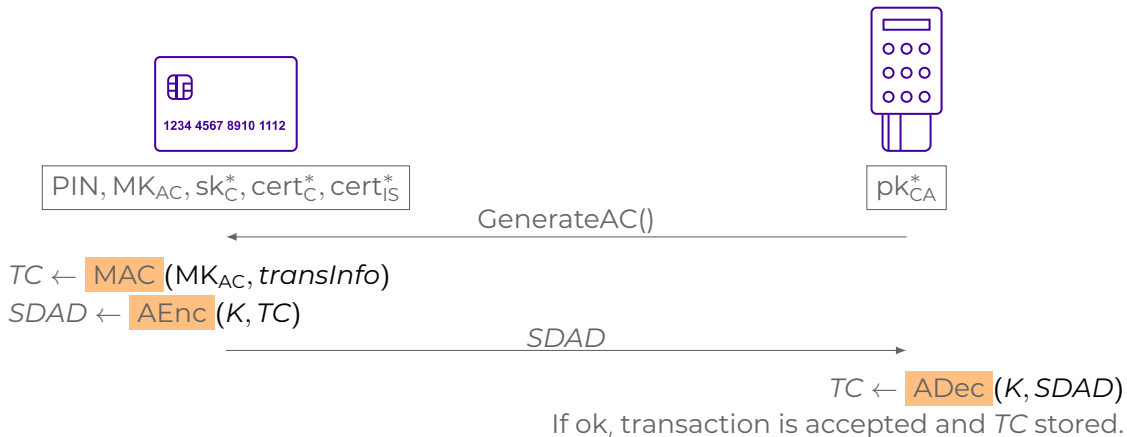
VerifyPIN(c)

$uPIN \leftarrow ADec(K, c)$

check $uPIN = PIN$

Success/Failed

PQ analog of CDA protocol (4/4)



Hybrid Protocol

Hybrid Cryptography

Combine classical and post quantum cryptography.

Why Hybrid Cryptography?

- › Proposed PQ scheme could still be broken (e.g. Rainbow, SIKE)
- › Take advantage of mature algorithms
- › Solution promoted by several agencies

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Our strategy

- › Store both classic and PQ certificates on card
- › Concatenate signatures (e.g. $\text{RSA.Sign}(sk_C, data) \parallel \text{AEnc}(K, data)$)
- › Stack encryption (e.g. $\text{RSA.Enc}(pk_C, \text{AEnc}(K, PIN))$)

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Choice of PQ algorithms

NIST Standardization Process

- › Selection not finished when this work started
- › 4 finalists for KEM
- › 3 finalists for Signature

NIST Selected KEMs

- › CRYSTALS-Kyber
- › NTRU
- › Saber
- › Classic McEliece

NIST Selected Signatures

- › CRYSTALS-Dilithium
- › Falcon
- › Rainbow

- › Only a KEM is required on smartcard: signature/verification is performed outside.
- › Signature choice impacts only the **size of the certificates**

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Personalization overhead

Figures for an ARM Cortex-M3 device running at 100 MHz, with kyber-512.

Algorithm	Overhead (bytes)		Increase (%)
	Theoretical	Measured	
Falcon-512	4661	5120	39.48
Falcon-1024	6785	7168	47.73
Dilithium-2	8584	9216	54.01
Dilithium-3	10 970	11 264	58.93
Dilithium-5	14 214	N/A	N/A

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Remarks

- › Same overhead for hybrid BDH variant
- › Difference theoretical/measured due to pagination
- › More personalization data \rightsquigarrow slower perso = less cards produced: **huge impact**

Full transaction timings

Figures for an ARM Cortex-M3 device running at 100 MHz, with kyber-512.

Protocol	Certificate Algo.	Timings (ms)			Ratio
		Comm.	Card	Total	
Hybrid CDA	Falcon-512	5267	376	5643	3.36
	Falcon-1024	7302	376	7678	4.58
	Dilithium-2	9018	384	9402	5.61
	Dilithium-3	11 297	402	11 699	6.98

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Remarks

- › Ratio w.r.t. classic protocol
- › Most of the time spent in **communication!**
- › Slightly worse for hybrid BDH variant
- › Very similar results with `nttrup52048509`

Conclusion

Summary

- › Using KEM for authentication is more efficient
- › **Communication cost** is significantly higher than computation one
- › In this context, **short signatures** should be preferred
- › Same conclusions apply equally to CRYSTALS-Kyber and NTRU
- › See the paper for BDH variant

Future Work

- › Formal proof of the protocols
- › More studies of post-quantum migration in other domains



Questions?



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