

Toward a New Block Cipher Mode Standard: Reasoning about Requirements Featuring the NECST Framework

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June 20, 2024

Announcements

- NIST Workshop on Formal Methods within Certification Programs (FMCP 2024)
 - July 23-25, This Room, \$205

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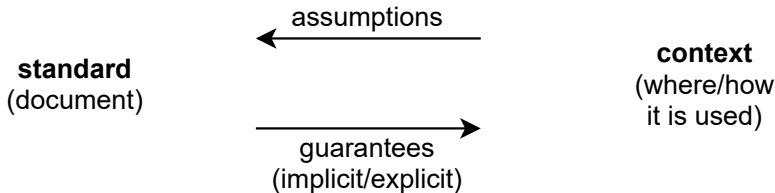
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- Information Security Conference (ISC 2024)
 - October 23-25, Arlington, VA

- NIST commitment:
 - Periodical review of standards
- First review: AES (FIPS 197)
 - NISTIR 8319: Review of the AES (July 2021)
- Next: Modes of operation (SP 800-38 Series)
 - Draft NISTIR 8459: Modes Report (March 2023)

- Where is the standard used?
- What security properties are required there?
- Does failure of security properties lead to **attacks**?



Security Models

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 - Purpose of encryption if attacker already knows the plaintext?
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- Again, focus on **attacks**...
 - Exhaustive search is infeasible, but...
 - ...side-channel attacks recover key in minutes?
 - Block cipher secure against related-key attacks, or...
 - ...disallow generating keys with known relations?
 - Vulnerabilities (CVE numbers)?

NIST SP 800-38 Series

- NIST SP 800-38A + Add: ECB, CBC, CFB, OFB, CTR
 -
- NIST SP 800-38B: CMAC
 -
- NIST SP 800-38C: CCM
 -
- NIST SP 800-38D: GCM
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- NIST SP 800-38E: XTS-AES
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- NIST SP 800-38F: KW, KWP, TKW
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- NIST SP 800-38G: FF1, FF3
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 - CPA security (except ECB)
- NIST SP 800-38B: CMAC
 - MAC security
- NIST SP 800-38C: CCM
 - CPA + MAC = CCA security
- NIST SP 800-38D: GCM
 - CPA + MAC = CCA security
- NIST SP 800-38E: XTS-AES
 - CCA up-to-block (without MAC)
- NIST SP 800-38F: KW, KWP, TKW
 - CCA (without MAC)
- NIST SP 800-38G: FF1, FF3
 - CCA (without MAC)

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 - CPA security (except ECB), IV (except ECB)
- NIST SP 800-38B: CMAC
 - MAC security
- NIST SP 800-38C: CCM
 - CPA + MAC = CCA security, nonce
- NIST SP 800-38D: GCM
 - CPA + MAC = CCA security, nonce
- NIST SP 800-38E: XTS-AES
 - CCA up-to-block (without MAC), tweak
- NIST SP 800-38F: KW, KWP, TKW
 - CCA (without MAC), no tweak
- NIST SP 800-38G: FF1, FF3
 - CCA (without MAC), tweak

Where Are 38A Modes Used?

- Length-preserving encryption
 - XTS (38E): only for storage devices
(but 38A modes “continue to be approved for such devices”)
 - Applications exist where ciphertext cannot be expanded...
- Building block for AEAD
 - Authentication-only mode: CMAC (38B), HMAC (FIPS 198-1)
 - Generic AEAD: e.g., CBC + HMAC
 - CCM (38C) and GCM (38D): AEAD based on CTR!

Always Use AEAD Modes?

- GCM (38C) and CCM (38D)
 - Based on **CTR**
- Nonce reuse:
 - Deduce plaintext from ciphertext difference!
- Short tag / no tag:
 - Control plaintext through ciphertext difference!

Alternatives?

- Nonce reuse: KW/KWP/TKW (38F) or FF1/FF3 (38G)
 - “Misuse resistant” AEAD
 - Most suitable for key wrapping (38F) or formatted data (38G)
 - Very slow for general use...
- No tag: XTS (38E)
 - ECB-like mode: independently encrypts every block...
 - Only for storage devices
 - CBC (38A): sometimes preferable?

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 - Lightweight encryption?
 - SHA-3-based encryption?
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- Why?
 - IV reuse, short or no tags, no key commitment, release of unverified plaintext,...
 - Many applications: disk encryption, packet encryption, message franking,...

- Third Block Cipher Modes Workshop (BCM3)
 - “NIST is particularly interested in discussing the possibility of standardizing a tweakable wide block encryption technique that could support a large range of input lengths”
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Where We Are Now

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 - Mode with provable security based on underlying block cipher

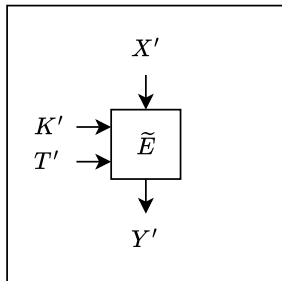
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- This Workshop
 - Mode with provable security based on underlying block cipher
- Next
 - Requirements for mode of operation?
 - Block cipher: AES or Rijndael? Related keys? Tweak input?

Hierarchy

AEAD, Key Wrapping, Disk Encryption, etc. (derived functions) ← only encoding here

Accordion Cipher (mode of operation) ← all crypto here

(Tweakable) Block Cipher ← AES/Rijndael here



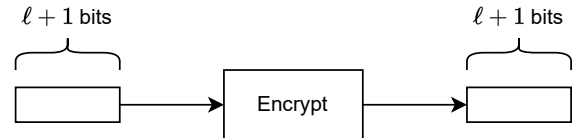
- New features / future-proofing
- Efficiency / performance
- Compatibility (can it replace existing standards?)
- Security
- Technical reason (something doesn't work / can't be done)

- Rogaway's Observation on SP 800-38 Series (BCM3)
 - SP 800-38 Series: 17 AES-based modes
 - Starting from scratch: maybe we would just have one?
 - "A mode that is like a wide block cipher"

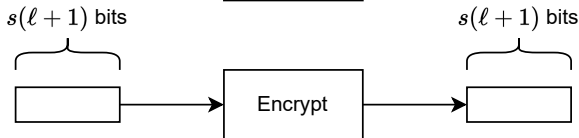
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 - "A mode that is like a wide block cipher"
- Requirements
 - (C) Can replace SP 800-38 Series modes
 - (T) Except 38G (FF1, FF3) (e.g., radix causes headaches)

- Accordion Cipher: $Y = E_{K,T}(X)$
 - Key $K \in \{0,1\}^k$
 - Tweak $T \in \{0,1\}^*$
 - Input X , Output $Y \in \{0,1\}^{m+\ell g}$, $\ell \in \{0,1,2,\dots\}$
 - Minimum input size m
 - Granularity g
 - Invertible

- CTR, OFB
($l \geq 0$)



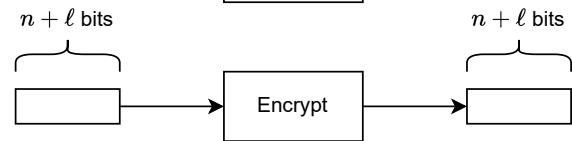
- CFB
($s =$ size of segment)



- ECB, CBC
($n =$ block size)



- CBC-CS
(CS1, CS2, CS3)



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- (T) $m = 2n$ bits may simplify designs (e.g., AEZ v5 [\[HKR17\]](#))

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or

- (S) large g helps to hide plaintext length [GJLN22]
(but more secure: length-hiding padding)
- (C,N) $g = 8$ bits sometimes needed: e.g., 509-byte for Tor

Maximum Input Size For 38-Series Compatibility

- Maximum Input Size?
 - (C) No limit specified in 38A for CBC, OFB, CFB
(CTR limit can be small in practice, e.g.: 64 KiB [RW03])
 - (S) Stay well below $2^{n/2}$ blocks [BL16]
 - (N) BCM3: consider $n = 256$ (512?) to avoid limits in practice
(Note: NIST LWC call required at least $2^{50} - 1$ bytes)

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- (E) Note: Increasing n is much more efficient than mode with beyond-birthday-bound security

Key Size for 38-Series Compatibility

- Key Size k
 - (C,S) 128, 192, 256 bits
 - (E) Fastest: 128 bits
 - (N) 512 bits? (seems not necessary: [\[Mou21\]](#))

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 - (E) Fastest: 128 bits
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- Bit Strength
 - NIST LWC call: 128 bit length, 112 bit strength
 - (E) 256 bit length, 128 bit strength?

Tweak Size for 38-Series Compatibility

- Tweak Size
 - (C) empty for Key Wrapping (38F),
1... $2^{64} - 1$ bits for GCM IV,
0... *maxTlen* bytes for FF1 (no restrictions on *maxTlen*)
 - (S) Stay well below $2^{n/2}$ blocks
 - (T) Only byte strings
 - (N) Avoid limits in practice

(N) Extra Requirements?

- Nonce-Hiding, Length-Hiding, Short-Tag, RUP-Secure MRAE?
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- Inverse-Free?
 - CMAC, GCM, CCM, FF1, etc. do not use inverse block cipher

(E) Block Cipher Requirements

```
rijndael256_new : bash — Konsole
File Edit View Bookmarks Plugins Settings Help
e136: 62 32 3d 48 dc c7 vaesenc %zmm23,%zmm8,%zmm8
e13c: 62 f2 7d 48 8d db vpermb %zmm3,%zmm0,%zmm3
e142: 62 b2 65 48 dc dd vaesenc %zmm21,%zmm3,%zmm3
e148: 62 f2 7d 48 8d ed vpermb %zmm5,%zmm0,%zmm5
e14e: 62 b2 55 48 dc ee vaesenc %zmm22,%zmm5,%zmm5
e154: 62 f2 7d 48 8d ff vpermb %zmm7,%zmm0,%zmm7
e15a: 62 b2 45 48 dc ff vaesenc %zmm23,%zmm7,%zmm7
e160: 62 f2 7d 48 8d d2 vpermb %zmm2,%zmm0,%zmm2
e166: 62 b2 6d 48 dc d5 vaesenc %zmm21,%zmm2,%zmm2
e16c: 62 f2 7d 48 8d e4 vpermb %zmm4,%zmm0,%zmm4
e172: 62 b2 5d 48 dc e6 vaesenc %zmm22,%zmm4,%zmm4
e178: 62 f2 7d 48 8d f6 vpermb %zmm6,%zmm0,%zmm6
e17e: 62 b2 4d 48 dc f7 vaesenc %zmm23,%zmm6,%zmm6
e184: 62 f2 7d 48 8d c9 vpermb %zmm1,%zmm0,%zmm1
e18a: 62 52 7d 48 8d c0 vpermb %zmm8,%zmm0,%zmm8
e190: 62 b2 75 48 dc cd vaesenc %zmm21,%zmm1,%zmm1
e196: 62 52 3d 48 dd c1 vaesenclast %zmm9,%zmm8,%zmm8
e19c: 62 f2 7d 48 8d db vpermb %zmm3,%zmm0,%zmm3
e1a2: 62 b2 65 48 dc de vaesenc %zmm22,%zmm3,%zmm3
e1a8: 62 f2 7d 48 8d ed vpermb %zmm5,%zmm0,%zmm5
e1ae: 62 b2 55 48 dc ef vaesenc %zmm23,%zmm5,%zmm5
e1b4: 62 f2 7d 48 8d ff vpermb %zmm7,%zmm0,%zmm7
e1ba: 62 d2 45 48 dd f9 vaesenclast %zmm9,%zmm7,%zmm7
e1c0: 62 f2 7d 48 8d d2 vpermb %zmm2,%zmm0,%zmm2
rijndael256_new : bash x
```

(E) Block Cipher Requirements

- AES-256: one AES round = one clock cycle throughput since Sandy Bridge (2011)
 - $\frac{14 \text{ rounds} \times 1 \text{ cycle/round}}{16 \text{ bytes/round}} \approx 0.88 \text{ cycles/byte}$

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- Rijndael ($n = 256$, also $n = 512$) on same platform?
 - Same as AES would be: 0.22 cycles/byte
 - Rijndael benchmarks: 0.25 cycles/byte
 - Benchmark speed independent of key schedule!
 - (S) fix related-key attacks and (N) possibly add tweak input?
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 - (E) ... but many keys/tweaks will have impact on performance
- (E) Note: New mode needs to be parallel (so that bottleneck is throughput, not latency)

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