Cryptography Research Center

CryptographicEstimators

A Software Library for Cryptographic Hardness Estimation

Andre Esser, Javier Verbel, Floyd Zweydinger and Emanuele Bellini

@NIST PQC Seminar, Nov. 2023



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- Hardness Estimation
- CryptographicEstimators
- 2. Theoretical Considerations
 - Time and Memory
 - Memory Access Costs
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- 7. Future Developments

Introduction

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Cryptographic Hardness Estimation

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Cryptographic Hardness Estimation

Estimation of required time to solve a (cryptographic) problem

Security guarantees

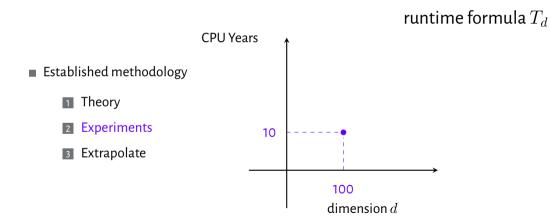
- Parameter selection
 - Example: RSA keysize recommendations

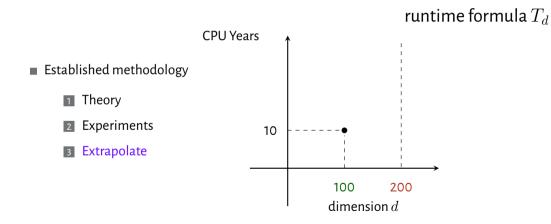
Estimates change over time: adaptive process

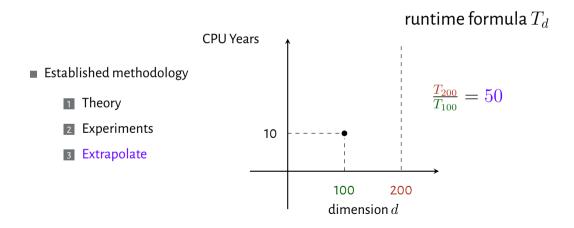
Established methodology

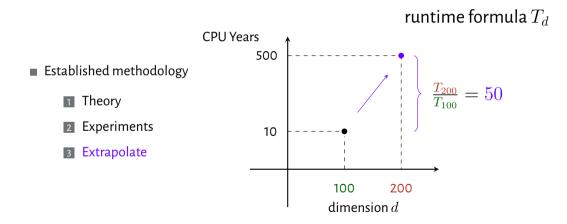
- Established methodology
 - 1 Theory
 - 2 Experiments
 - 3 Extrapolate

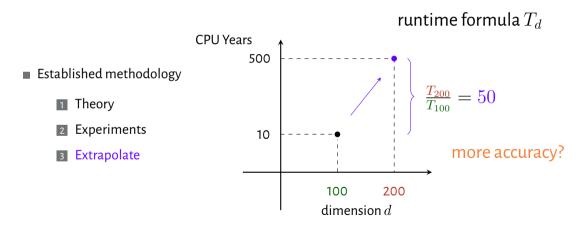
runtime formula T_d **CPU** Years Established methodology 1 Theory 2 Experiments 3 Extrapolate dimension d

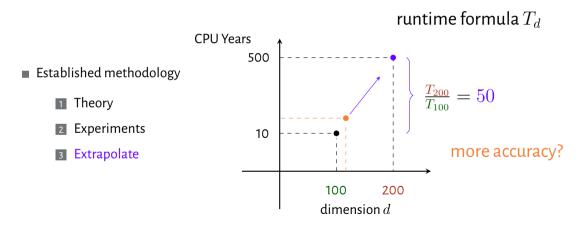


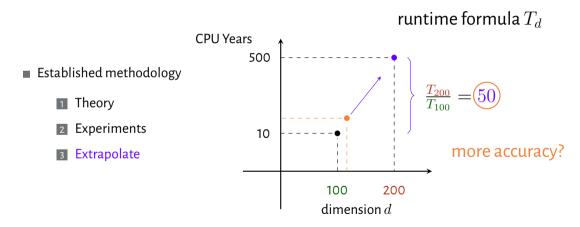


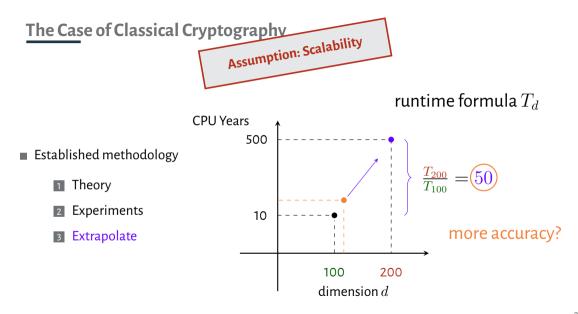














Difficult scalability (memory)





Time to solve a problem = Time of fastest known algorithm



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Hardness depends on best known algorithms

Requires estimation of time of all known algorithms



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Main Challenges



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- Main Challenges
 - Consensus

The Case of PQC

■ Difficult scalability (memory) ⇒ estimation methodology

Time to solve a problem = Time of fastest known algorithm

Hardness depends on best known algorithms

Requires estimation of time of all known algorithms

- Main Challenges
 - Consensus
 - Accessibility





Python / Sage library for estimations of cryptographic problems

CryptographicEstimators

Python / Sage library for estimations of cryptographic problems

Main goals

- State-of-the-art estimations
- Centralization of estimation efforts
- Community-driven open-source project
- Easy accessibility

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Python / Sage library for estimations of cryptographic problems

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- Community-driven open-source project
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- Current State: 6 Estimators, 32 Algorithms
 - Multivariate Quadratic (MQ)
 - Binary Syndrome Decoding (SD)
 - Syndrome Decoding over \mathbb{F}_q (SDFq)
 - Permutation Equivalence (PE)
 - Linear Equivalence (LE)
 - Permuted Kernel (PK)

Theoretical Considerations





What to Estimate?

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- Memory: Measured in basic elements el

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 - **Example MQ-Problem** : op : \mathbb{F}_q -multiplication and el : \mathbb{F}_q -element
 - Example binary SD-Problem: op : \mathbb{F}_2^n -vector addition and el : \mathbb{F}_2^n -vector

What to Estimate?

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- Problem defines op / el to bit (operation) conversion





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 - Accessing 1 bit in memory of size M takes $f(M) = \sqrt{M}, \sqrt[3]{M}, \dots$ bit operations



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- Upper bound on memory access: $T \cdot f(M)$
 - $\blacksquare \ \text{Real cost } C \text{ of the full algorithm} \colon T \leq C \leq T \cdot f(M)$

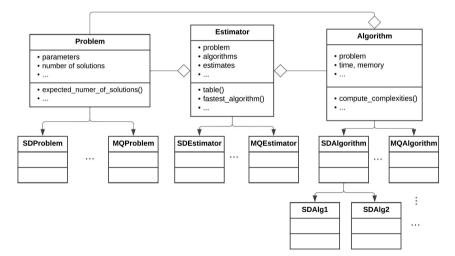
Technical Design







CryptographicEstimators: An object-oriented Python library



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Usage

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From web application ²:



■ Install -> sage ->

from cryptographic_estimators.SDEstimator import SDEstimator A = SDEstimator(n=500, k=250, w=50) A.table()

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■ Install -> sage ->

From web application ²:

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	algorithm		time		memory	I
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	BJMMdw		64.0		25.2	I
	BJMMpdw		63.9		28.0	I
	ВЈММ		63.6		29.4	I
	BJMM_plus		63.5		26.5	I
	BothMay		63.2		25.2	I
	Dumer		63.9		28.5	I
	May0zerov		62.7		34.0	I
	Prange		77.7		17.3	I
	Stern		63.8		22.9	I
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- Install -> sage ->
- Docker¹->make docker-run->same as local

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Technology Innovation Institute Crypto	ographicEstimators				
Binary Syndrom	Binary Syndrome Decoding 🛛 👻 🗙				
Code length * 🛛 🛛	500				
Code dimension * 🛛 😧	250				
Error weight * 🛛 🛛 🛛 🛛 🖓	50				
Memory limit 🛛 😧					
Insert value]				
Leave empty if no limit is desired Number of solutions	d D				
Insert value	mount of solutions				
Optional parameters					
Estimators parameters	>				



■ Install -> sage ->

Docker¹->make docker-rur

From web application ²:

Results				
Algorithms BallCollision, BJM Y	Estimate			
BallCollision, BJM +	Time, Memory, Par *			
	_		Estimate	
Algorithm	Time	Memory	Parameters	+
BallCollision	64	22	r: 6 p: 2 pl: 0 l: 12	+
BJMM	63	29	r: 6 depth: 2 p: 4 p1: 2 t: 26	+
BJMM_plus	63	26	r:6 p:4 pl:2 l:26 ll:9	+
BothMay	63	25	r:6 p:4 wE0 w2:0 p1:2 l:10	+
Dumer	63	28	r:6 t:19 p:3	+
MayOzerov	62	33	r:6 depthc2 p:6 p1:4 t:25	+
Prange	77	17	n 6	+
Stern	63	22	r: 6 p: 2 l: 13	+
			Let Export configuration	🗄 Download table.tex

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<pre>sage: from cryptographic_estimators.MQEstimator import MQEstimator</pre>
sage: E = MQEstimator(n=24, m=24, q=16)
sage: E.algorithm_names()
['BooleanSolveFXL',
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Access single algorithms: complexity of crossbred -> sage: from cryptographic_estimators.MQEstimator import MQEstimator sage: E = MQEstimator(n=24, m=24, q=16) sage: E.crossbred.time_complexity() 70.8336959616846

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Access single algorithms: optimal parameters -> age: from cryptographic_estimators.MQEstimator import MQEstimator age: E = MQEstimator(n=24, m=24, q=16) sage: E.crossbred.optimal_parameters() ('D': 9, 'd': 1, 'k': 11}

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• Access single algorithms: complex. for $(D, d, k) = (6, 1, 3) \rightarrow$ sage: from cryptographic_estimators.HQEstimator import HQEstimator sage: E = MQEstimator(n=24, m=24, q=16) sage: E.crossbred time_complexity(k=3, D=6, d=1) 98.20496250072115

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A full user guide is <u>available</u>.

Contributing

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³https://github.com/Crypto-TII/CryptographicEstimators.



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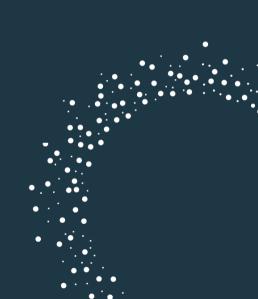
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- 2 Participate in the discussion (within the repository).

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NIST PQC Signatures







- Current included estimators:
 - Multivariate Quadratic (MQ)



■ Multivariate Quadratic (MQ) —-> $\mathbf{y} = \mathcal{P}(\mathbf{x})$, with \mathcal{P} a quadratic map.



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Coverage

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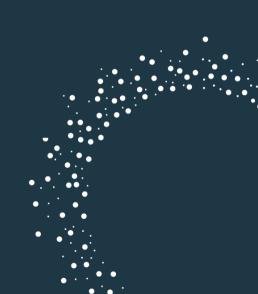
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Best known attacks of 8 / 30 (remaining) submissions fall into this scope

Estimation of NIST Candidates

Estimates for NIST Category I parameter sets

Scheme	Hardness Assumption	Est. Time	Est. Memory
SDitH	SDFq	147.0	26.9
LESS	LE	136.6	39.0
PERK	РК	155.5	154.4
MQOM	MQ	142.8	51.8
TUOV / UOV	MQ	144.5	59.6
VOX	MQ	153.0	59.8
PROV	MQ	150.1	62.3



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 - Scheme vs. Problem estimators
 - Advanced memory access

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 - MinRank -> NIST, TII and University of Limoges
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 - Regular Syndrome Decoding -> Marche Polytechnic University, TII, and potential other collaborators.

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