

Identity based Authentication in Session Initiation

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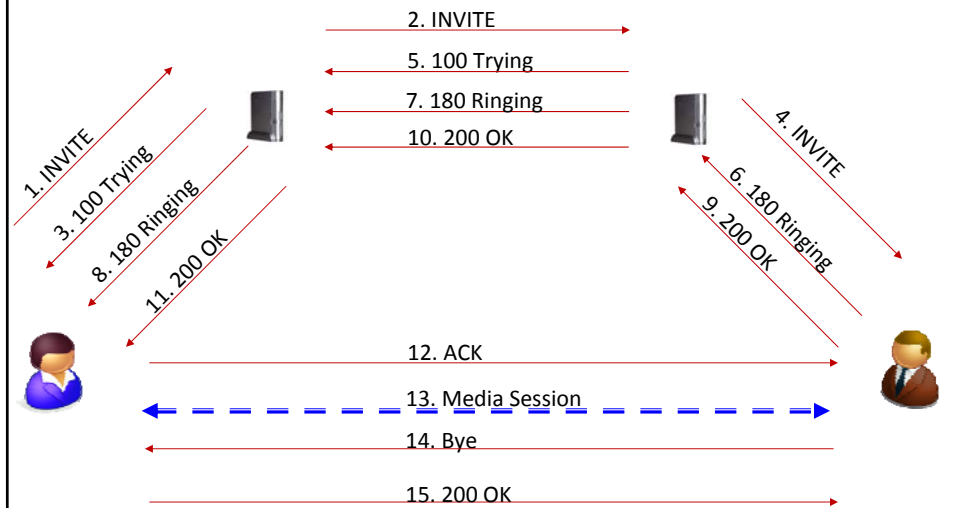
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Samsung Telecommunications

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Session Initiation Protocol



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INVITE message in SIP

INVITE sip:bob@biloxi.com SIP/2.0

Via: SIP/2.0/UDP pc33.atlanta.com;branch=z9hG4bK776asdhdhs

Max-Forwards: 70

To: Bob <sip:bob@biloxi.com>

Spooled — From: Alice <sip:alice@atlanta.com>;tag=1928301774

Call-ID: a84b4c76e66710@pc33.atlanta.com

CSeq: 314159 INVITE Contact: <sip:alice@pc33.atlanta.com>

Content-Type: application/sdp Content-Length: 142

(Alice's SDP not shown)

RFC 4474

INVITE sip:bob@biloxi.example.org SIP/2.0

Contact: <sip:alice@pc33.atlanta.example.com>

Identity:

"ZYNBbHC00VMZr2kZt6VmCvPonWJMGvQTBdqghoWeLxJfzB2a1pxAr3VgrB0SsAa
ifsRdiOPoQZYoy2wrVghuhcsMbHWUSFxI6p6q5TOQXHMmz6uEo3svJsSH49thyGn
FVcnyaZ++yRIBYYQLqWzJ+KVhPKbfU/pryhVn9Yc6U="

Identity-Info: <https://atlanta.example.com/atlanta.cer>;alg=rsa-sha1

Content-Type: application/sdp

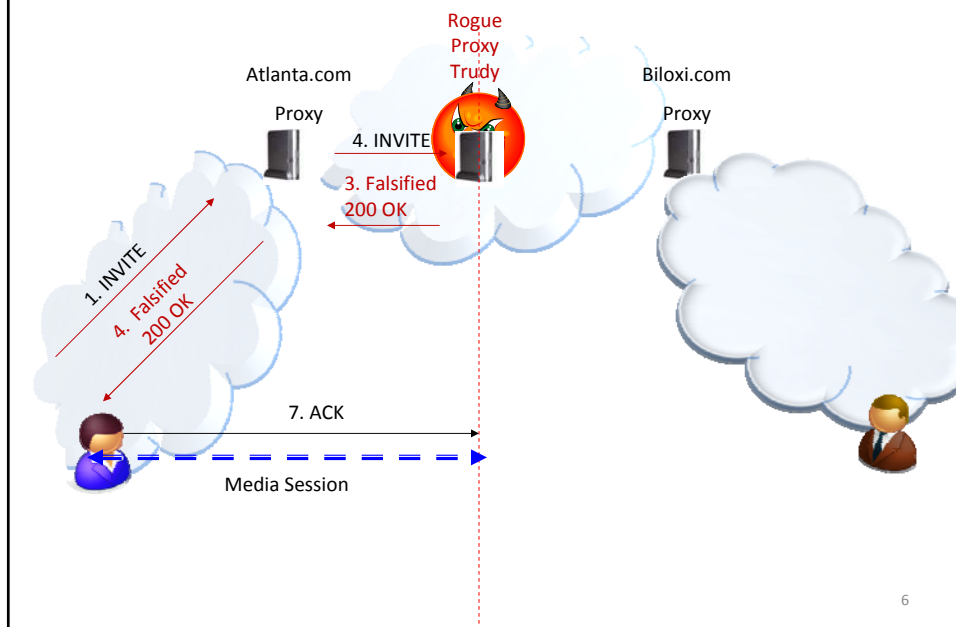
Content-Length: 147

(Alice's SDP not shown)

We cannot apply RFC 4474 to SIP responses

- Response messages cannot be challenged.
- SIP response messages may not encode the identity of the responder

Rogue Proxy sends a falsified 200 OK



Trudy sends a falsified 200 OK message to Alice

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SIP/2.0 200 OK
Via: SIP/2.0/UDP server10.biloxi.com;branch=z9hG4bK4b43c2ff8.1 ;received=192.0.2.3
Via: SIP/2.0/UDP bigbox3.site3.atlanta.com;branch=z9hG4bK77ef4c2312983.1
;received=192.0.2.2
Via: SIP/2.0/UDP pc33.atlanta.com;branch=z9hG4bKnashds8 ;received=192.0.2.1
To: Bob <sip:bob@biloxi.com>;tag=a6c85cf
From: Alice <sip:alice@atlanta.com>;tag=1928301774
Call-ID: a84b4c76e66710
CSeq: 314159 INVITE
Contact: <IP address of the Rogue Proxy>
Content-Type: application/sdp Content-Length: 131
v=0
o=bob 2890844527 2890844527 IN IP4 client.biloxi.example.com
s=-
c=IN IP4 < IP Address of the Rogue Proxy >
t=0 0
m=audio 49172 RTP/AVP 0
a=rtpmap:0 PCMU/8000
    
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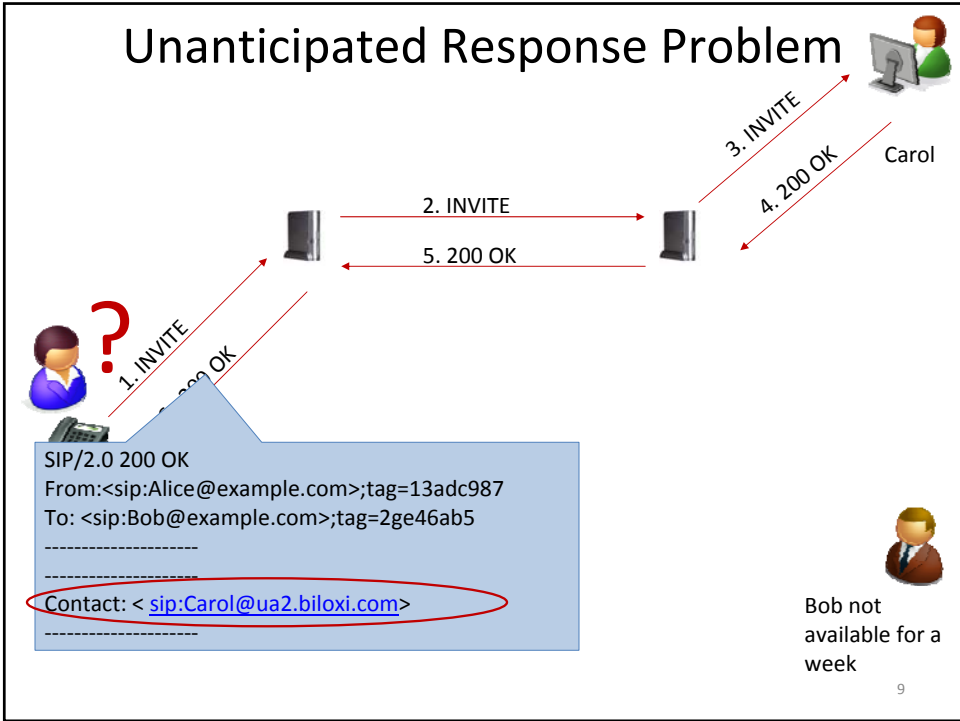
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Approach: Transform Response Identity problem into Connected Identity Problem



Messages 1-3 convey Alice's Identity to Bob
 Messages 7-9 convey Bob's Identity to Alice

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Drawbacks

INVITE
with the
Identity field

→

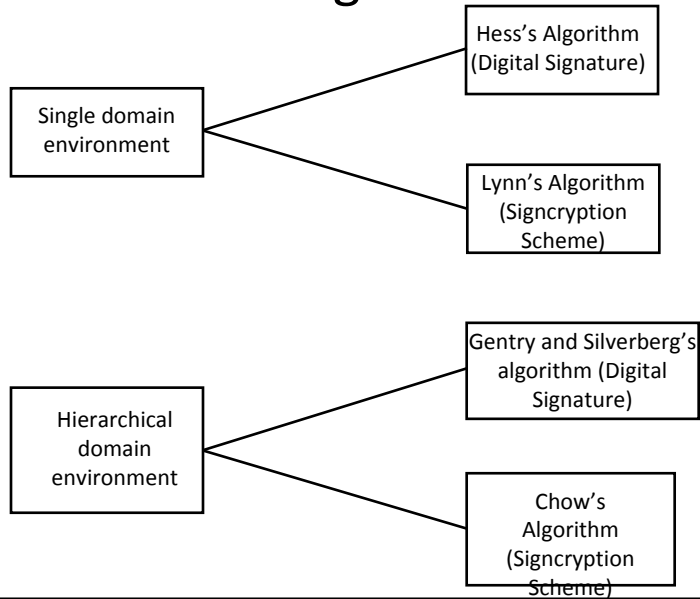
How do we verify the Identity field ?

- Self signed certificates
- Dependence on PKI
 - Discovery of a public key certificate [Linn,Brancaud04]
 - Complex path construction process

Di. Berbecaru, A. Lioy and M. Marian "On the Complexity of Public-Key Certificate Validation, " in Proceedings of the 4th International Conference on Information Security, Lecture Notes in Computer Science, Springer-Verlag, Vol 2200 pages 183-203, 2001

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Identity based signature algorithms



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Signature and Key Size

Criteria	RFC 4474	IBS Schemes
Signature Size	175 bytes(sig) + 512 bytes + CA certs	Hess's algorithm 511 bytes
		Lynn's 434 bytes
		Gentry & Silverberg algorithm 434 bytes
		Chow et al's algorithm 434 bytes
Key size	1024 bits	160 bits

H. Kupwade Patil and D. Willis, " Identity based authentication in SIP", IETF draft 2008.

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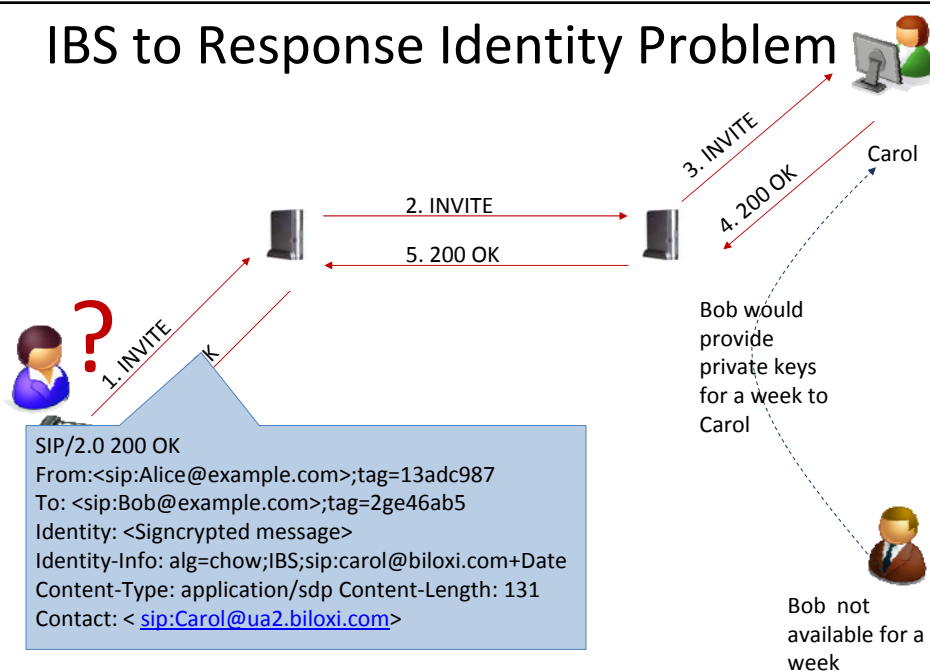
Computational Time

Scheme	Generation time in sec	Verification time in sec
Open SSL (RFC 4474)	0.109s	0.110s
PBC library Hess's algorithm	0.078s	0.051s
PBC library Lynn's algorithm	0.269s	0.238s
PBC library Gentry and Silverberg's algorithm	0.093s	0.063s
PBC library Chow et. al's algorithm	0.160s	0.162s

H. Kupwade Patil and D. Willis, "Identity based authentication in SIP", IETF draft 2008.

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IBS to Response Identity Problem



H. Kupwade Patil and D. Willis, "Identity based signcryption scheme to the connected identity problem in the SIP"

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Key Distribution in IBE

[B. Lee et. al. 2004]

$$e : G_1 \times G_1 \rightarrow G_2$$

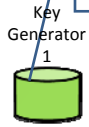
$$H_3 : \{G_2\} \rightarrow F_p$$

$$X = xP \text{ where } x \in Z \text{ and } P \in G_1$$

- 2. $Q_{ID_A} = H_1(ID_A)$
- 3. $Q_{bl_A} = H_3[e(s_0 X, P_0)]s_0 Q_{ID_A}$
- 4. $Sig(Q_{bl_A}) = s_0 Q_{bl_A}$



1. Request for partial private key X, ID



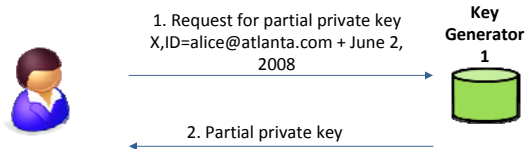
5. Partial private key $Sig(Q_{bl_A})$
 Q_{bl_A}

$$D_{ID} = \frac{Q_0}{H_3[e(P_0, P_0)^X]}$$

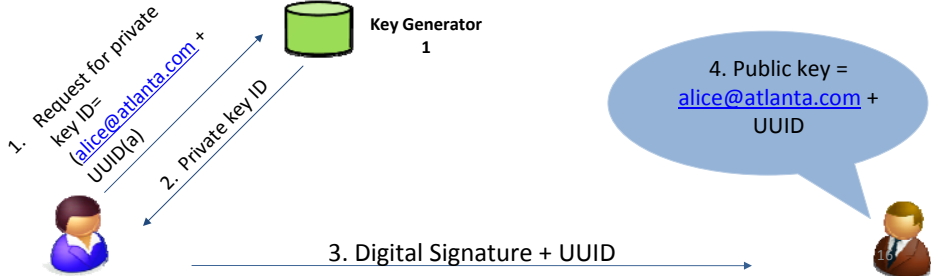
[B. Lee et. al 2004] B. Lee, C. Boyd, E. Dawson, K. Kim, J. Yang and S. Yoo, "Secure Key Issuing in ID-based Cryptography," in *Conferences in Research and Practice in Information Technology, 2004*, vol. 32, pp. 69-74.

Revocation issues

- Expiration



- Using Universally Unique User ID (UUID)



Conclusion

- Identity based signature/sigcrypton schemes
 - Reduces the complex path construction process used by PKI
 - Faster processing speed compared to the RSA based schemes (RFC 4474)

Future Work

- Identity based authentication in a peer to peer SIP environment

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Thank You

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