

# SPHINCS: practical stateless hash-based signatures

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## Traditional hash-based signatures: security vs. usability

*There's still a critical flaw in this! Recall that using a one-time signature twice **completely breaks the system**. Because of that, the signature scheme is “stateful”. This means that, when signing, the signer absolutely must record that a one-time key has been used so that they never use it again. If the private key was copied onto another computer and used there, then the whole system is broken.*

*That limitation might be ok in some situations, and actually means that one can build forward-secure signature schemes: schemes where signatures prior to a key compromise can still be trusted. Perhaps for a CA where the key is in an HSM that might be useful. However, for most environments it's a **huge foot-cannon**.*

—Adam Langley, “Hash based signatures”, 2013 (emphasis added)

## Beyond hash-based: factoring-based foot-cannons!

Modern understanding of essential structure of  
“provably secure” 1984 Goldwasser–Micali–Rivest:

One-time signature scheme based on factoring.

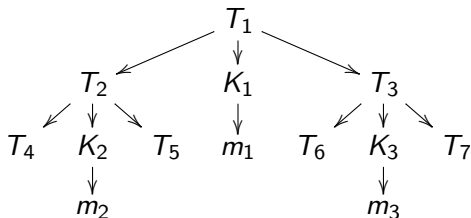
Stateful many-time signature scheme:

use one-time public key  $K_i$  to sign  $i$ th message;

use one-time public key  $T_1$  to sign  $(T_2, T_3, K_1)$ ;

use one-time public key  $T_2$  to sign  $(T_4, T_5, K_2)$ ;

use one-time public key  $T_3$  to sign  $(T_6, T_7, K_3)$ ; etc.



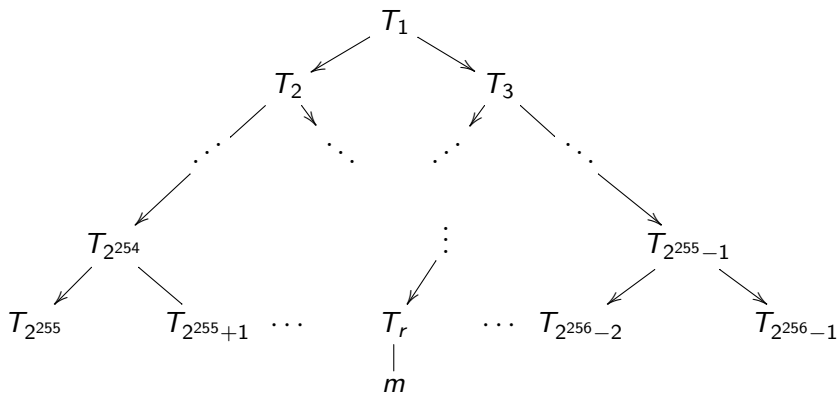
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## Huge trees (1987 Goldreich), keys on demand (Levin)

Signer chooses random  $r \in \{2^{255}, 2^{255} + 1, \dots, 2^{256} - 1\}$ ,  
uses one-time public key  $T_r$  to sign message;  
uses one-time public key  $T_i$  to sign  $(T_{2i}, T_{2i+1})$  for  $i < 2^{255}$ .  
Generates  $i$ th secret key as  $H_k(i)$  where  $k$  is master secret.



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Would dominate traffic in typical applications, and add user-visible latency on typical network connections.

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Debian operating system is designed for frequent upgrades.

At least one new signature for each upgrade.

Typical upgrade: one package or just a few packages.

1.2 MB average package size.

0.08 MB median package size.

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Example:

HTTPS typically sends multiple signatures per page.

1.8 MB average web page in Alexa Top 1000000.



# New: SPHINCS-256

## Reasonable sizes.

0.041 MB signature.

0.001 MB public key.

0.001 MB private key.

## Reasonable speeds.

**Benchmarks** of our public-domain software on Haswell:

51.1 million cycles to sign. (RSA-3072: 14.2 million.)

1.5 million cycles to verify. (RSA-3072: 0.1 million.)

3.2 million cycles for keygen. (RSA-3072: 950 million.)

## Designed for $2^{128}$ post-quantum security,

even for a user signing more than  $2^{50}$  messages:

$2^{20}$  messages/second continuously for more than 30 years.

Yes, we did the analysis of quantum attacks.

# Ingredients of SPHINCS (and SPHINCS-256)

Drastically reduce tree height (to 60).

Replace one-time leaves with few-time leaves.

Optimize few-time signature size *plus* key size.

New few-time HORST, improving upon HORS.

Use hyper-trees (12 layers), as in GMSS.

Use masks, as in XMSS and XMSS<sup>MT</sup>,  
for standard-model security proofs.

Optimize short-input (256-bit) hashing speed.

Use sponge hash (with ChaCha12 permutation).

Use fast stream cipher (again ChaCha12).

Vectorize hash software and cipher software.

See paper for details: [sphincs.cr.yp.to](http://sphincs.cr.yp.to)

