The attached DRAFT document (provided here for historical purposes) has been superseded by the following publication:

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Title: Interfaces for Personal Identity Verification

Publication Date: May 2015 (updated 2/8/2016)

• Final Publication: <a href="https://doi.org/10.6028/NIST.SP.800-73-4">https://doi.org/10.6028/NIST.SP.800-73-4</a> (which links to <a href="http://nvlpubs.nist.gov/nistpubs/SpecialPublications/NIST.SP.800-73-4.pdf">http://nvlpubs.nist.gov/nistpubs/SpecialPublications/NIST.SP.800-73-4.pdf</a>).

• Information on other NIST Computer Security Division publications and programs can be found at: <a href="http://csrc.nist.gov/">http://csrc.nist.gov/</a>

The following information was posted with the attached DRAFT document:

May 13, 2013

SP 800-73-4

DRAFT Interfaces for Personal Identity Verification (3 Parts)

Part 1- PIV Card Application Namespace, Data Model and Representation

Part 2- PIV Card Application Card Command Interface

Part 3- PIV Client Application Programming Interface

NIST announces that *Draft Special Publication (SP) 800-73-4*, *Interfaces for Personal Identity Verification*, has been released for public comment. The Draft SP 800-73-4 is updated to align with Candidate Final FIPS 201-2. Major changes in Draft SP 800-73-4 include:

- Removal of Part 4, The PIV Transitional Data Model and Interfaces:
- The addition of specifications for secure messaging and the virtual contact interface, both of which are optional to implement;
- The specification of an optional Cardholder Universally Unique Identifier (UUID) as a unique identifier for a cardholder;
- The specification of an optional on-card biometric comparison mechanism, which may be used as a means of performing card activation and as a PIV authentication mechanism; and
- The addition of a requirement for the PIV Card Application to enforce a minimum PIN length of six digits.

Except for minor editorial changes, all changes can be reviewed with the track-change version (See Track Change file for Part 1-3 below) of Draft SP 800-73-4.

NIST requests comments on Draft SP 800-73-4 by 5:00pm EDT on *June 14, 2013*. Please submit your comments, using the comment template form (see last link for this draft below) to piv\_comments@nist.gov with "Comments on Public Draft SP 800-73-4" in the subject line.

Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface

> Ramaswamy Chandramouli David Cooper Hildegard Ferraiolo Salvatore Francomacaro Ketan Mehta Jason Mohler

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## COMPUTER SECURITY

29	Draft NIST Special Publication 800-73-4
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31	<b>Interfaces for Personal Identity</b>
32	Verification – Part 2: PIV Card
33	<b>Application Card Command</b>
34	Interface
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National Institute of Standards and Technology Patrick D. Gallagher, Under Secretary of Commerce for Standards and Technology and Director

## Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface

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120	
121	Abstract
122 123	EIDS 201 defines the requirements and characteristics of a government wide interconcreble identity.
123	FIPS 201 defines the requirements and characteristics of a government-wide interoperable identity
124	credential. FIPS 201 also specifies that this identity credential must be stored on a smart card. This document, SP 800-73, contains the technical specifications to interface with the smart card to retrieve
125	and use the PIV identity credentials. The specifications reflect the design goals of interoperability and
127	PIV Card functions. The goals are addressed by specifying a PIV data model, card edge interface, and
128	application programming interface. Moreover, this document enumerates requirements where the
129	international integrated circuit card standards [ISO7816] include options and branches. The
130	specifications go further by constraining implementers' interpretations of the normative standards. Such
131	restrictions are designed to ease implementation, facilitate interoperability, and ensure performance, in a
132	manner tailored for PIV applications.
133	
134	Keywords
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137	Personal Identity Verification (PIV); physical access control; smart cards; secure messaging
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## 1. Introduction

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- 242 Homeland Security Presidential Directive-12 (HSPD-12) called for a common identification standard to
- be adopted governing the interoperable use of identity credentials to allow physical and logical access to
- 244 Federally controlled facilities and information systems. Personal Identity Verification (PIV) of Federal
- Employees and Contractors, Federal Information Processing Standard 201 (FIPS 201) [FIPS 201] was
- developed to establish standards for identity credentials. Special Publication 800-73-4 (SP 800-73-4)
- 247 contains technical specifications to interface with the smart card (PIV Card<sup>1</sup>) to retrieve and use the
- 248 identity credentials.

## 1.1 Purpose

- 250 FIPS 201 defines procedures for the PIV lifecycle activities including identity proofing, registration, PIV
- 251 Card issuance, and PIV Card usage. FIPS 201 also specifies that the identity credentials must be stored
- on a smart card. SP 800-73-4 contains the technical specifications to interface with the smart card to
- 253 retrieve and use the identity credentials. The specifications reflect the design goals of interoperability and
- 254 PIV Card functions. The goals are addressed by specifying a PIV data model, card edge interface, and
- application programming interface. Moreover, SP 800-73-4 enumerates requirements where the
- international integrated circuit card (ICC) standards [ISO7816] include options and branches. The
- specifications go further by constraining implementers' interpretations of the normative standards. Such
- restrictions are designed to ease implementation, facilitate interoperability, and ensure performance, in a
- 259 manner tailored for PIV applications.

## 260 **1.2 Scope**

- SP 800-73-4 specifies the PIV data model, application programming interface (API), and card interface
- requirements necessary to comply with the use cases, as defined in Section 6 of FIPS 201 and further
- described in Appendix B of SP 800-73-4 Part 1. Interoperability is defined as the use of PIV identity
- credentials such that client-application programs, compliant card applications, and compliant ICCs can be
- used interchangeably by all information processing systems across Federal agencies. SP 800-73-4 defines
- the PIV data elements' identifiers, structure, and format. SP 800-73-4 also describes the client application
- programming interface and card command interface for use with the PIV Card.
- 268 This part, SP 800-73-4 Part 2 PIV Card Application Card Command Interface, contains the technical
- specifications of the PIV Card command interface to the PIV Card. The specification defines the set of
- commands surfaced by the PIV Card Application at the card edge of the ICC.

#### 1.3 Audience and Assumptions

- 272 This document is targeted at Federal agencies and implementers of PIV systems. Readers are assumed to
- 273 have a working knowledge of smart card standards and applications.
- 274 Readers should also be aware of SP 800-73-4 Part 1, Section I, for the revision history of SP 800-73,
- 275 Section II, which details configuration management recommendations, and Section III, which specifies
- NPIVP conformance testing procedures. Section 1.3 of Part 1 specifies the effective date of SP 800-73-4.

<sup>&</sup>lt;sup>1</sup> A physical artifact (e.g., identity card, "smart" card) issued to an individual that contains a PIV Card Application which stores identity credentials (e.g., photograph, cryptographic keys, digitized fingerprint representation) so that the claimed identity of the cardholder can be verified against the stored credentials by another person (human readable and verifiable) or an automated process (computer readable and verifiable).

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## Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface

## 277 **1.4 Content and Organization**

- All sections in this document are *normative* (i.e., mandatory for compliance) unless specified as *informative* (i.e., non-mandatory). Following is the structure of Part 2:
- 280 + Section 1, *Introduction*, provides the purpose, scope, audience, and assumptions of the document and outlines its structure.
- + Section 2, *Overview: Concepts and Constructs*, describes the model of computation of the PIV Card Application and the PIV client application programming interface including information processing concepts and data representation constructs.
- + Section 3, *PIV Card Application Card Command Interface*, describes the set of commands accessible by the PIV Middleware to communicate with the PIV Card Application.
- + Section 4, *Secure Messaging*, describes the secure messaging protocol that is used to enable data confidentiality and integrity.
- + Appendix A, *Examples of the Use of the GENERAL AUTHENTICATE Command*, demonstrates the GENERAL AUHTENTICATE command. This section is *informative*.
- 4 Appendix B, *Terms, Acronyms, and Notation*, contains the list of terms and acronyms used in this document and explains the notation in use. This section is *informative*.
- + Appendix C, *References*, contains the lists of documents used as references by this document.
  This section is *informative*.

295	2. Overview: Concepts and Constructs		
296 297	SP 800-73-4 Parts 2 and 3 define two interfaces to an ICC that contains the PIV Card Application: a low-level card command interface (Part 2) and a high-level client API (Part 3).		
298 299 300	The information processing concepts and data constructs on both interfaces are identical and may be referred to generically as the information processing concepts and data constructs on the <i>PIV interfaces</i> without specific reference to the client API or the card command interface.		
301 302 303 304	The client API provides task-specific programmatic access to these concepts and constructs and the card command interface provides communication access to concepts and constructs. The client API is used by client applications using the PIV Card Application. The card command interface is used by software implementing the client API (middleware).		
305 306 307 308	The client API is thought of as being at a higher level than the card command interface because access to a single entry point on the client API may cause multiple card commands to traverse the card command interface. In other words, it may require more than one card command on the card command interface to accomplish the task represented by a single call on an entry point of the client API.		
309 310 311 312	The client API is a program execution, call/return style interface whereas the card command interface is a communication protocol, command/response style interface. Because of this difference, the representation of the PIV concepts and constructs as bits and bytes on the client API may be different from the representation of these same concepts and constructs on the card command interface.		
313	2.1.1 Platform Requirements		
314 315	The following are the requirements that the PIV Card Application places on the ICC platform on which it is implemented or installed:		
316	+ global security status that includes the security status of a global cardholder PIN		
317	+ application selection using a truncated Application Identifier (AID)		
318	+ ability to reset the security status of an individual application		
319 320	+ indication to applications as to which physical communication interface – contact versus contactless – is in use		
321	+ support for the default selection of an application upon warm or cold reset		
322 323	2.2 Namespaces of the PIV Card Application		
324 325 326 327	AID, names, Tag-Length-Value (BER-TLV) tags [ISO8825], ASN.1 Object Identifiers (OIDs) [ISO8824 and Proprietary Identifier eXtensions (PIXes) of the NIST Registered Application Provider IDentifier (RID) used on the PIV interfaces are specified in Part 1. Part 1 also specifies that all unspecified names, BER-TLV tags, OIDs, and values of algorithm identifiers, key references, and cryptographic mechanism identifiers, are received for future use.		

## Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface

329	2.3 Card Applications
330 331 332	Each command that appears on the card command interface shall be implemented by a <i>card application</i> that is resident on the ICC. The card command enables operations on and with the data objects to which the card application has access.
333 334 335 336 337	Each card application shall have a globally unique name called its Application Identifier (AID) [ISO7816, Part 4]. Except for the default applications, access to the card commands and data objects of a card application shall be gained by selecting the card application using its application identifier <sup>2</sup> . The PIX of the AID shall contain an encoding of the version of the card application. The AID of the PIV Card Application is defined in Part 1.
338 339	The card application whose commands are currently being used is called the <i>currently selected application</i> .
340	2.3.1 Default Selected Card Application
341 342 343 344	The card platform shall support a default selected card application. In other words, there shall be a currently selected application immediately after a cold or warm reset. This card application is the default selected card application. The default card application may be the PIV Card Application, or it may be another card application.
345	2.4 Security Architecture
346 347	The security architecture of an ICC is the means by which the security policies governing access to each data object stored on the card are represented within the card.
348 349	These security policy representations are applied to all PIV card commands thereby ensuring that the prescribed data policies for the card applications are enforced.
350	The following subsections describe the security architecture of the PIV Card Application.
351	2.4.1 Access Control Rule
352 353 354	An <i>access control rule</i> shall consist of an <i>access mode</i> and a <i>security condition</i> . The access mode is an operation that can be performed on a data object. A security condition is a Boolean expression using variables called security statuses that are defined below.
	,

2.4.2 Security Status

359

Associated with each authenticable entity shall be a set of one or more Boolean variables, each called a security status indicator of the authenticable entity. Each security status indicator, in turn, is associated

<sup>&</sup>lt;sup>2</sup> Access to the default application, and its commands and objects, occurs immediately after a warm or cold card reset without an explicit SELECT command.

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- 362 with a credential that can be used to authenticate the entity. The security status indicator of an 363 authenticable entity shall be TRUE if the credentials associated with the security status indicator of the authenticable entity have been authenticated and FALSE otherwise. 364 365 A successful execution of an authentication protocol shall set the security status indicator associated with 366 the credential used in the protocol to TRUE. An aborted or failed execution of an authentication protocol 367 shall set the security status indicator associated with the credential used in the protocol to FALSE. 368 As an example, the credentials associated with three security status indicators of the cardholder might be: 369 PIN, primary fingerprint, and secondary fingerprint. Demonstration of knowledge of the PIN is the 370 authentication protocol for the first security status indicator wherein the PIN is the credential. 371 Comparison of the fingerprint template on the card with a fingerprint acquired from the cardholder is the 372 authentication protocol for the other two security status indicators wherein the fingerprint is the 373 credential. A security condition using these three security status indicators might be "PIN AND (primary fingerprint **OR** secondary fingerprint)." 374 375 A security status indicator shall be said to be a global security status indicator if it is not changed when 376 the currently selected application changes from one application to another. In essence, when changing 377 from one application to another, the global security status indicators shall remain unchanged. 378 A security status indicator is said to be an application security status indicator if it is set to FALSE when 379 the currently selected application changes from one application to another. Every security status indicator 380 is either a global security status indicator or an application security status indicator. The security status 381 indicators associated with the PIV Card Application PIN, the PIN Unblocking Key (PUK), the primary finger OCC, the secondary finger OCC, Pairing Code, and the PIV Card Application Administration Key 382 383 are application security status indicators for the PIV Card Application, whereas the security status 384 indicator associated with the Global PIN is a global security status indicator. 385 The term *global security status* refers to the set of all global security status indicators. The term 386 application security status refers to the set of all application security status indicators for a specific 387 application. 388 2.4.3 Authentication of an Individual 389 Knowledge of a PIN is the means by which an individual can be authenticated to the PIV Card Application. 390 391 The PIV Card Application PIN and the pairing code shall each be between 6 and 8 bytes in length. If the actual length of PIV Card Application PIN or pairing code is less than 8 bytes it shall be padded to 8 392 393 bytes with 'FF' when presented to the card command interface. The 'FF' padding bytes shall be appended to the actual value of the PIN. The bytes comprising the PIV Card Application PIN and pairing code shall 394 395 be limited to values 0x30 - 0x39, the ASCII values for the decimal digits '0' - '9'. For example, + Actual PIV Card Application PIN: "123456" or '31 32 33 34 35 36' 396 397
  - + Padded PIV Card Application PIN presented to the card command interface: '31 32 33 34 35 36 FF FF'
- The PIV Card Application shall enforce the minimum length requirement of six bytes for the PIV Card Application PIN and pairing code (i.e., shall verify that at least the first six bytes of the value presented to the card command interface are in the range 0x30 0x39).

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- 402 If the Global PIN is used by the PIV Card Application then the above encoding, length, padding, and
- 403 enforcement of minimum PIN length requirements for the PIV Card Application PIN shall apply to the
- 404 Global PIN.
- The PUK shall be 8 bytes in length, and the bytes comprising the PUK shall be limited to the values 0x00
- 406 0xFE (i.e., shall not include 'FF'). The PIV Card Application shall enforce the PUK length requirement
- of eight bytes (i.e., shall verify that all eight bytes of the value presented to the card command interface
- 408 are in the range 0x00 0xFE).

## 2.5 Current State of the PIV Card Application

- 410 The elements of the *current state* of the PIV Card Application when the PIV Card Application is the
- 411 currently selected application are described in Table 1.

Table 1. State of the PIV Card Application

State Name	Always Defined	Comment	Location of State
Global security status	Yes	Contains security status indicators that span all card applications on the platform.	PIV Platform
Currently selected application Yes application  The platform shall support the a card application using the ful identifier or by providing the right truncated version and there shall support the a card application to the platform shall support the a card application to the platform shall support the a card application to the platform shall support the a card application to the platform shall support the a card application using the full support the action to the platform to t		The platform shall support the selection of a card application using the full application identifier or by providing the right-truncated version and there shall always be a currently selected application.	PIV Platform
Application security status	Yes Contains security status indicators local to the PIV Card Application.		PIV Card Application

409

## 3. PIV Card Application Card Command Interface

Table 2 lists the card commands surfaced by the PIV Card Application at the card edge of the ICC when it is the currently selected card application. All PIV Card Application card commands shall be supported by

417 a PIV Card Application. Card commands indicated with a 'Yes' in the Command Chaining column shall

support command chaining for transmitting a data string too long for a single command as defined in

419 [ISO7816].

418

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### **Table 2. PIV Card Application Card Commands**

Туре	Name	Contact Interface	Contactless Interface	Security Condition for Use	Command Chaining
PIV Card Application Card	SELECT	Yes	Yes	Always	No
Commands for Data Access	GET DATA	Yes	Yes	Data Dependent. See Table 2, Part 1.	No
	VERIFY	Yes	SM or VCI	Always	Yes <sup>3</sup>
PIV Card Application Card	CHANGE REFERENCE DATA	Yes	VCI	PIN	No
Commands for Authentication	RESET RETRY COUNTER	Yes	No	PIN Unblocking Key	No
	GENERAL AUTHENTICATE	Yes	Yes (See Note)	Key Dependent. See Table 4, Part 1.	Yes
PIV Card Application Card	PUT DATA	Yes	No	PIV Card Application Administrator	Yes
Commands for Credential Initialization and Administration	GENERATE ASYMMETRIC KEY PAIR	Yes	No	PIV Card Application Administrator	Yes

The PIV Card Application shall return the status word of '6A 81' (Function not supported) when it

receives a card command on the contactless interface marked "No" in the Contactless Interface column in

424 Table 2.

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426

Note: Cryptographic protocols using private/secret keys that require the "PIN" or "OCC" security

condition shall only be used on the contactless interface after a Virtual Contact Interface (VCI) has been

427 established.

\_

<sup>&</sup>lt;sup>3</sup> The VERIFY command is only required to support command chaining if the PIV Card Application supports on-card biometric comparison (OCC).

## 428 3.1 PIV Card Application Card Commands for Data Access

#### 3.1.1 SELECT Card Command

- 430 The SELECT card command sets the currently selected application. The PIV Card Application shall be
- selected by providing its application identifier (see Part 1, Section 2.2) in the data field of the SELECT
- 432 command.

429

- 433 There shall be at most one PIV Card Application on any ICC. The PIV Card Application can also be
- made the currently selected application by providing the right-truncated version (see Part 1, Section 2.2);
- that is, without the two-byte version number in the data field of the SELECT command.
- The complete AID, including the two-byte version, of the PIV Card Application that became the currently
- selected card application upon successful execution of the SELECT command (using the full or right-
- 438 truncated PIV AID) shall be returned in the application property template.
- 439 If the currently selected application is the PIV Card Application when the SELECT command is given
- and the AID in the data field of the SELECT command is either the AID of the PIV Card Application or
- the right-truncated version thereof, then the PIV Card Application shall continue to be the currently
- selected card application and the setting of all security status indicators in the PIV Card Application shall
- be unchanged.
- 444 If the currently selected application is the PIV Card Application when the SELECT command is given
- and the AID in the data field of the SELECT command is not the PIV Card Application (or the right-
- truncated version thereof), but a valid AID supported by the ICC, then the PIV Card Application shall be
- deselected and all the PIV Card Application security status indicators in the PIV Card Application shall
- be set to FALSE.
- 449 If the currently selected application is the PIV Card Application when the SELECT command is given
- and the AID in the data field of the SELECT command is an invalid AID not supported by the ICC, then
- 451 the PIV Card Application shall remain the currently selected application and all PIV Card Application
- security status indicators shall remain unchanged.

### 453 **Command Syntax**

CLA	'00'
INS	'A4'
P1	'04'
P2	'00'
Lc	Length of application identifier
Data Field	AID of the PIV Card Application using the full AID or the right-truncated AID (See Section 2.2, Part 1)
L <sub>e</sub>	Length of application property template

## Response Syntax

Data Field	Application property template (APT). See Table 3 below
SW1-SW2	Status word

Upon selection, the PIV Card Application shall return the application property template described in Table 3.

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Table 3. Data Objects in the PIV Card Application Property Template (Tag '61')

Description	Tag	M/O/C	Comment	
Application identifier of application	'4F'	М	The PIX of the AID includes the encoding of the version of the PIV Card Application. See Section 2.2, Part 1.	
Coexistent tag allocation authority	'79'	М	Coexistent tag allocation authority template. See Table 4.	
Application label	'50'	0	Text describing the application; e.g., for use on a man-machine interface.	
Uniform resource locator	'5F50'	0	Reference to the specification describing the application.	
Cryptographic algorithms supported	'AC'	С	Cryptographic algorithm identifier template, see Table 5.	

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## Table 4. Data Objects in a Coexistent Tag Allocation Authority Template (Tag '79')

Name	Tag	M/O	Comment
Application identifier	'4F'	М	See Section 2.2, Part 1

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A PIV Card Application may use a subset of the cryptographic algorithms defined in SP 800-78. Tag

0xAC encodes the cryptographic algorithms supported by the PIV Card Application. The encoding of the
tag 0xAC shall be as specified in Table 5. Each instance of tag 0x80 shall encapsulate one algorithm.

The presence of cryptographic algorithm identifier 0x27 or 0x2B indicates the PIV Card Application
supports secure messaging. Tag 0xAC shall be present and indicate algorithm identifier 0x27 and/or

469 0x2B when the PIV Card Application supports secure messaging.

Table 5. Data Objects in a Cryptographic Algorithm Identifier Template (Tag 'AC')

Name	Tag	M/O	Comment
Cryptographic algorithm identifier	'80'	М	For values see Table 5 of Part 1 and [SP800-78, Table 6-2]
Object identifier	'06'	М	Its value is set to 0x00

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SW1	SW2	Meaning
'6A'	'82'	Application not found
'90'	'00'	Successful execution

### 475 3.1.2 GET DATA Card Command

The GET DATA card command retrieves the data content of the single data object whose tag is given in the data field.<sup>4</sup>

## Command Syntax

CLA	'00' or '0C' for secure messaging
INS	'CB'
P1	'3F'
P2	'FF'
L <sub>c</sub>	Length of data field*
Data Field	See Table 6
Le	Number of data content bytes to be retrieved.

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\* The  $L_c$  value is '05' for all PIV data objects except for the 0x7E interindustry tag (Discovery Object) and the application property template (APT), which have an  $L_c$  value of '03'.

### Table 6. Data Objects in the Data Field of the GET DATA Card Command

Name	Tag	M/O	Comment	
Tag list	'5C'	М	BER-TLV tag of the data object to be retrieved.	See Table 3, Part 1.

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## **Response Syntax**

485 For the 0x7E Discovery Object (if present):

Data Field	BER-TLV of the 0x7E Discovery data object (see Section 3.3.2, Part 1 for a description of the Discovery Object's structure returned in the data field).
SW1-SW2	Status word

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For all other PIV data objects:

Data Field	BER-TLV with the tag '53' containing in the value field of the requested data object.
SW1-SW2	Status word

SW1	SW2	Meaning	
'61'	'xx'	Successful execution where SW2 encodes the number of response data bytes still available	
'69'	'82'	Security status not satisfied	
'6A'	'82'	Data object not found	
'90'	'00'	Successful execution	

<sup>.</sup> 

<sup>&</sup>lt;sup>4</sup> The GET RESPONSE command is used in conjunction with GET DATA to accomplish the reading of larger PIV data objects. The GET RESPONSE command is illustrated in Appendix A.4.1 (Command 3).

### **Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface**

## PIV Card Application Card Commands for Authentication

489	3.2	PIV Card Application Card Commands for Authenticatio
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3.2.1 VERIFY Card Command

491	The VERIFY card command initiates the comparison in the car	d of the reference data indicated by the
<b>サノエ</b>	The VERM I card command initiates the comparison in the car	d of the reference data maleated by the

- 492 key reference with authentication data in the data field of the command.
- 493 Key reference '80' specific to the PIV Card Application (i.e., local key references) and, optionally, the
- 494 Global PIN with key reference '00', the OCC data (key references '96' and '97'), and pairing code (key
- 495 reference '98') are the only key references that may be verified by the PIV Card Application's VERIFY
- 496

- 497 Key reference '80' shall be able to be verified by the PIV Card Application VERIFY command.
- 498 If the PIV Card Application contains the Discovery Object as described in Part 1, and the first byte of the
- 499 PIN Usage Policy value is 0x60 or 0x70, then key reference '00' shall be able to be verified by the PIV
- 500 Card Application VERIFY command.
- 501 If the PIV Card Application contains the Discovery Object as described in Part 1 and the first byte of PIN
- 502 Usage Policy is 0x50 or 0x70, then key reference '98' shall be able to be verified by the PIV Card
- 503 Application VERIFY command.
- 504 If the PIV Card Application contains the Discovery Object as described in Part 1 and the Biometric
- 505 Information Template (BIT) is present, then key references '96' and/or '97' shall be able to be verified by
- 506 the PIV Card Application VERIFY command.
- 507 If the key reference is '00' or '80' and the VERIFY command is not submitted over either the contact
- 508 interface or the VCI, or if the key reference is '96', '97', or '98', and the VERIFY command is submitted
- 509 over the contactless interface without secure messaging, then the card command shall fail, and the PIV
- Card Application shall return the status word '6A 81'. The security status and the retry counter of the key 510
- 511 reference shall remain unchanged.
- 512 If the current value of the retry counter associated with the key reference is zero, then the comparison
- 513 shall not be made, and the PIV Card Application shall return the status word '69 83'.
- 514 If the key reference is '00', '80', or '98', and the authentication data in the command data field does not
- 515 satisfy the criteria in Section 2.4.3, then the card command shall fail, and the PIV Card Application shall
- 516 return the status word '6A 80'. The security status and the retry counter of the key reference shall remain
- 517 unchanged.
- 518 If the key reference is '96' or '97' and the authentication data in the command data field is not of length
- 519 3N, where N satisfies the requirements for minimum and maximum number of minutiae specified in the
- 520 BIT, then the card command shall fail, and the PIV Card Application shall return the status word '6A 80'.
- 521 The security status and the retry counter of the key reference shall remain unchanged.
- 522 If the authentication data in the command data field does not match reference data associated with the key
- 523 reference, then the card command shall fail. If the card command fails, the security status of the key
- 524 reference shall be set to FALSE and the retry counter associated with the key reference shall be
- 525 decremented by one.

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- If the card command succeeds, then the security status of the key reference shall be set to TRUE and the retry counter associated with the key reference shall be set to the reset retry value associated with the key
- reference. The initial value of the retry counter and the reset retry value associated with the key
- reference, i.e., the number of successive failures (retries) before the retry counter associated with the key
- reference reaches zero, are issuer dependent.
- The VERIFY command shall reset the security status of the key reference in P2, when the P1 parameter is
- 532 'FF' and both L<sub>c</sub> and the data field are absent. The security status of the key reference specified in P2
- shall be set to FALSE and the retry counter associated with the key reference shall remain unchanged.

## **Command Syntax**

CLA	'00' or '10' indicating command chaining '0C' or '1C' for secure messaging	
INS '20'		
P1	'00' or 'FF'	
P2 Key reference. See Part 1, Table 4.		
Absent <sup>5</sup> – for absent command data field  L <sub>c</sub> '08' – for PIV Card Application PIN, Global PIN, or pairing code  3N – for OCC data (where N is the number of minutiae)		
Data Field  Absent <sup>5</sup> , PIV Card Application PIN, Global PIN, or pairing correference data as described in Section 2.4.3, or OCC data a described in Table 9 of SP 800-76-2.		
Le	Absent	

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## **Response Syntax**

SW1	SW2	Meaning	
'63'	'CX'	Verification failed, X indicates the number of further allowed retries	
'69'	'83'	Authentication method blocked	
'6A'	'80'	Incorrect parameter in command data field	
'6A'	'81'	Function not supported	
'6A'	'88'	Key reference not found	
'90'	'00'	Successful execution	

 $<sup>^{5}</sup>$  If P1='00', and  $L_{c}$  and the command data field are absent, the command can be used to retrieve the number of further retries allowed ('63 CX'), or to check whether verification is not needed ('90 00').

### 3.2.2 CHANGE REFERENCE DATA Card Command

- 541 The CHANGE REFERENCE DATA card command initiates the comparison of the verification data with
- the current value of the reference data and if this comparison is successful, replaces the reference data
- with new reference data.

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- Only reference data associated with key references '80' and '81' specific to the PIV Card Application (i.e.,
- local key reference) and the Global PIN with key reference '00' may be changed by the PIV Card
- 546 Application CHANGE REFERENCE DATA command. If any other key reference value is specified the
- 547 PIV Card Application shall return the status word '6A 81'.
- 548 If the CHANGE REFERENCE DATA command is not submitted over either the contact interface or the
- VCI, then the card command shall fail, and the PIV Card Application shall return the status word '6A 81'.
- 550 The security status and the retry counter of the key reference shall remain unchanged.
- Key reference '80' reference data shall be changed by the PIV Card Application CHANGE REFERENCE
- DATA command. The ability to change reference data associated with key references '81' and '00' using
- the PIV Card Application CHANGE REFERENCE DATA command is optional.
- If either the current reference data or the new reference data in the command data field of the command
- does not satisfy the criteria in Section 2.4.3, the PIV Card Application shall not change the reference data
- associated with the key reference and shall return the status word '6A 80', and retry counter shall remain
- unchanged.
- 558 If the current value of the retry counter associated with the key reference is zero, then the reference data
- associated with the key reference shall not be changed and the PIV Card Application shall return the
- 560 status word '69 83'.
- If the card command succeeds, then the security status of the key reference shall be set to TRUE and the
- retry counter associated with the key reference shall be set to the reset retry value associated with the key
- reference.
- If the card command fails, then the security status of the key reference shall be set to FALSE and the retry
- counter associated with the key reference shall be decremented by one.
- The initial value of the retry counter and the reset retry value associated with the key reference, i.e., the
- number of successive failures (retries) before the retry counter associated with the key reference reaches
- zero, is issuer dependent.

## **Command Syntax**

CLA	'00' or '0C' for secure messaging		
INS	'24'		
P1	'00'		
P2	'00' (Global PIN), '80' (PIV Card Application PIN), or '81' (PUK)		
L <sub>c</sub>	'10'		
Data Field  Current PIN reference data concatenated without delimitation new PIN reference data, both PINs as described in Section 2.			
L <sub>e</sub>	Absent		

## 571 Response Syntax

SW1	SW2	Meaning	
'63'	'CX'	Reference data change failed, X indicates the number of further allowed retries or resets	
'69'	'83'	Reference data change operation blocked	
'6A'	'80'	ncorrect parameter in command data field	
'6A'	'81'	Function not supported	
'6A'	'88'	Key reference not found	
'90'	'00'	Successful execution	

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#### 3.2.3 RESET RETRY COUNTER Card Command

- The RESET RETRY COUNTER card command resets the retry counter of the PIN to its initial value and changes the reference data. The command enables recovery of the PIV Card Application PIN in the case
- 576 that the cardholder has forgotten the PIV Card Application PIN.
- 577 The only key reference allowed in the P2 parameter of the RESET RETRY COUNTER command is the
- 578 PIV Card Application PIN. Any other key references in P2 shall not be permitted and the PIV Card
- Application shall return the status word '6A 81'.
- If the reset retry counter reference data (PUK) or the new reference data (PIN) in the command data field
- of the command does not satisfy the criteria in Section 2.4.3, the PIV Card Application shall not reset the
- retry counter associated with the PIN and shall return the status word '6A 80'. The PUK's retry counter
- shall remain unchanged.
- If the current value of the PUK's retry counter is zero, then the PIN's retry counter shall not be reset, and
- the PIV Card Application shall return the status word '69 83'.
- If the card command succeeds, then the PIN's retry counter shall be set to its reset retry value.
- Optionally, the PUK's retry counter may be set to its initial reset retry value. The security status of the
- 588 PIN's key reference shall not be changed.
- If the card command fails, then the security status of the PIN's key reference shall be set to FALSE, and
- the PUK's retry counter shall be decremented by one.
- The initial retry counter associated with the PUK, i.e., the number of failures of the RESET RETRY
- 592 COUNTER command before the PUK's retry counter reaches zero, is issuer dependent.

### 593 Command Syntax

CLA	'00'	
INS	'2C'	
P1	'00'	
P2	'80' (PIV Card Application PIN).	
Lc	'10'	
Data Field  Reset retry counter reference data (PUK) concatenated without delimita the new reference data (PIN) (both PUK and PIN as described in Section		
Le	-e Absent	

## 594 Response Syntax

SW1	SW2	Meaning	
'63'	'CX'	Reset failed, X indicates the number of further allowed resets	
'69'	'83'	Reset operation blocked	
'6A'	'80'	Incorrect parameter in command data field	
'6A'	'81'	Function not supported	
'6A'	'88'	Key reference not found	
'90'	'00'	Successful execution	

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#### 3.2.4 GENERAL AUTHENTICATE Card Command

- The GENERAL AUTHENTICATE card command performs a cryptographic operation, such as an authentication protocol, using the data provided in the data field of the command and returns the result of the cryptographic operation in the response data field.<sup>6</sup>
- The GENERAL AUTHENTICATE command shall be used with the PIV authentication keys ('9A', '9B',
- 601 '9E') to authenticate the card or a card application to the client application (INTERNAL
- AUTHENTICATE), to authenticate an entity to the card (EXTERNAL AUTHENTICATE), and to
- perform a mutual authentication between the card and an entity external to the card (MUTUAL
- 604 AUTHENTICATE).
- The GENERAL AUTHENTICATE command shall be used with the digital signature key ('9C') to realize
- the signing functionality on the PIV client application programming interface. Data to be signed is
- 607 expected to be hashed off card. Appendix A.4 illustrates the use of the GENERAL AUTHENTICATE
- 608 command for signature generation.
- The GENERAL AUTHENTICATE command shall be used with the key management key ('9D') and the
- retired key management keys ('82' '95') to realize key establishment schemes specified in SP 800-78
- 611 (ECDH and RSA). Appendix A.5 illustrates the use of the GENERAL AUTHENTICATE command for
- key establishment schemes aided by the PIV Card Application.
- The GENERAL AUTHENTICATE command shall be used with the PIV Secure Messaging key ('03')
- and cryptographic algorithm identifier '27' or '2B' to establish session keys for secure messaging as
- specified in Section 4. If key reference '03' is specified in P2 then algorithm identifiers in P1 other than
- 616 '27' and '2B' shall not be permitted and the PIV Card Application shall return the status word '6A 86'.
- 617 The GENERAL AUTHENTICATE command supports command chaining to permit the uninterrupted
- transmission of long command data fields to the PIV Card Application. If a card command other than the
- 619 GENERAL AUTHENTICATICATE command is received by the PIV Card Application before the
- 620 termination of a GENERAL AUTHENTICATE chain, the PIV Card Application shall rollback to the
- state it was in immediately prior to the reception of the first command in the interrupted chain. In other
- words, an interrupted GENERAL AUTHENTICATE chain has no effect on the PIV Card Application.

<sup>&</sup>lt;sup>6</sup> For cryptographic operations with larger keys, e.g., RSA 2048, the GET RESPONSE command is used to return the complete result of the cryptographic operation. The GET RESPONSE command is illustrated in Appendix A.4.1 (Command 3).

## 624 Command Syntax

CLA '00' or '10' indicating command chaining '0C' or '1C' for secure messaging			
INS	'87'		
P1	Algorithm reference. See Table 13 and [SP800-78, Table 6-2]		
P2	Key reference. See Table 4, Part 1 for key reference values		
L <sub>c</sub>	Length of data field		
Data Field	See Table 7		
L <sub>e</sub>	Absent or length of expected response		

### Table 7. Data Objects in the Dynamic Authentication Template (Tag '7C')

Name	Tag	M/O	Description
Witness	'80'	С	Demonstration of knowledge of a fact without revealing the fact. An empty witness is a request for a witness.
Challenge	'81'	С	One or more random numbers or byte sequences to be used in the authentication protocol.
Response	'82'	С	A sequence of bytes encoding a response step in an authentication protocol.
Exponentiation	'85'	С	A parameter used in ECDH key agreement protocol.

 The data objects that appear in the dynamic authentication template (tag '7C') in the data field of the GENERAL AUTHENTICATE card command depend on the authentication protocol being executed. The Witness (tag '80') contains encrypted data (unrevealed fact). This data is decrypted by the card. The Challenge (tag '81') contains clear data (byte sequence), which is encrypted by the card. The Response (tag '82') contains either the decrypted data from tag '80' or the encrypted data from tag '81'. Note that the empty tags (i.e., tags with no data) return the same tag with content (they can be seen as "requests for requests"):

- + '80 00' Returns '80 TL < encrypted random>' (as per definition)
- + '81 00' Returns '81 TL < random>' (as per external authenticate example)

## 637 Response Syntax

Data Field	Absent, authentication-related data, signed data, shared secret, or transported key
SW1-SW2	Status word

SW1	SW2	Meaning	
'61'	'xx'	Successful execution where SW2 encodes the number of response data bytes still available	
'69'	'82'	Security status not satisfied	
'6A'	'80'	Incorrect parameter in command data field	
'6A'	'86'	Incorrect parameter in P1 or P2	
'90'	'00'	Successful execution	

## 640 3.3 PIV Card Application Card Commands for Credential Initialization and Administration

## 3.3.1 PUT DATA Card Command

The PUT DATA card command completely replaces the data content of a single data object in the PIV

644 Card Application with new content.

## Command Syntax

CLA	'00' or '10' indicating command chaining	
INS	'DB'	
P1	'3F'	
P2	'FF'	
L <sub>c</sub>	Length of data field	
Data Field	See Tables 8 and 9	
L <sub>e</sub>	Absent	

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For the 0x7E Discovery Object (if present):

Table 8. Data Field of the PUT DATA Card Command for the Discovery Object

Tag	M/O	Description
'7E'	1\/1	BER-TLV of tag '7E' as illustrated in Section 3.3.2, Part 1.

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For all other PIV Data objects:

Table 9. Data Field of the PUT DATA Card Command for all other PIV Data Objects

Name	Tag	M/O	Description
Tag list	'5C'	М	Tag of the data object whose data content is to be replaced. See Table 3, Part 1.
Data	'53'	М	Data with tag '53' as an unstructured byte sequence.

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## **Response Syntax**

Data Field	Absent
SW1-SW2	Status word

SW1	SW2	Meaning
'69'	'82'	Security status not satisfied
'6A'	'81'	Function not supported
'6A'	'84'	Not enough memory
'90'	'00'	Successful execution

### 3.3.2 GENERATE ASYMMETRIC KEY PAIR Card Command

The GENERATE ASYMMETRIC KEY PAIR card command initiates the generation and storing in the card of the reference data of an asymmetric key pair, i.e., a public key and a private key. The public key of the generated key pair is returned as the response to the command. If there is reference data currently associated with the key reference, it is replaced in full by the generated data.

## **Command Syntax**

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CLA	'00' or '10' indicating command chaining
INS	'47'
P1	'00'
P2	Key reference. See Table 4 of Part 1 for a list of the key references
L <sub>c</sub>	Length of data field
Data Field	Control reference template. See Table 10
L <sub>e</sub>	Length of public key of data object template

Table 10. Data Objects in the Template (Tag 'AC')

Name	Tag	M/O	Description
Cryptographic mechanism identifier	'80'	М	See Part 1, Table 5
Parameter	'81'		Specific to the cryptographic mechanism

## **Response Syntax**

Data Field	Data objects of public key of generated key pair. See Table 11	
SW1-SW2	Status word	

Table 11. Data Objects in the Template (Tag '7F49')

Name	Tag
Public key data objects for RSA	
Modulus	'81'
Public exponent	'82'
Public key data objects for ECDSA	
Point	'86'

The public key data object in tag '86' is encoded as follows:

Table 12. Public Key encoding for ECDSA

Tag	Length	Value
'86'	L	04    X    Y [SECG, Section 2.3.3]

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- Note: The octet '04' indicates that the X and Y coordinates of point P are encoded without the use of point compression. The length L is 65 bytes for points on Curve P-256 and 97 bytes for points on Curve 670
- 671
- 672 P-384.

SW1	SW2	Meaning
'61'	'xx'	Successful execution where SW2 encodes the number of response data bytes still available
'69'	'82'	Security status not satisfied
'6A'	'80'	Incorrect parameter in command data field; e.g. unrecognized cryptographic mechanism
'6A'	'81'	Function not supported
'6A'	'86'	Incorrect parameter P2; cryptographic mechanism of reference data to be generated different than cryptographic mechanism of reference data of given key reference
'90'	'00'	Successful execution

## 4. Secure Messaging

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- 675 If a PIV Card Application implements the optional secure messaging protocol, it shall be implemented as
- specified in this section. Secure messaging is initiated through the use of a key establishment protocol.
- The key establishment protocol defined here is a one-way authentication protocol that authenticates the
- PIV Card Application to the client applicant and establishes a set of session keys that may be
- subsequently used to protect the communication channel between the two parties.
- 680 Section 4.1 describes the key establishment protocol used to support secure messaging in the PIV Card
- Application. Section 4.2 describes the use of secure messaging to protect commands and responses sent
- between the client application and the PIV Card Application.

## 4.1 The Key Establishment Protocol

- The key establishment protocol for the PIV Card Application is based on a simplified profile of
- OPACITY with Zero Key Management [ANSI504-1], as depicted below.

### Client Application (H)

## **PIV Card Application (ICC)**

```
CB_H = 0x10
                                                                                H1
Generate an ephemeral key pair (deH; QeH) from the domain
                                                                                H2.
parameters specified in the response to the SELECT command
                                                                                 H3
Send ID_{sH} \parallel Q_{eH} \parallel CB_H
                                                                                          ID_{sH} \parallel Q_{eH} \parallel CB_H
                                                                                                                    ID_{sICC} = T_8(SHA256(C_{ICC}*))
                                                                                                                                                                                                 C1
                                                                                                                    CB_{ICC} = CB_H \& 'F0'
                                                                                                                                                                                                 C2
                                                                                                                    Check that CB<sub>ICC</sub> is 0x10
                                                                                                                                                                                                 C3
                                                                                                                    Verify that QeH is a valid public key for the domain
                                                                                                                                                                                                 C4
                                                                                                                    parameters of O<sub>sICC</sub>
                                                                                                                                                                                                 C5
                                                                                                                    Z = ECC\_CDH(d_{sICC}, Q_{eH})
                                                                                                                    Generate nonce N<sub>ICC</sub>
                                                                                                                                                                                                 C6
                                                                                                                    SK_{CFRM} \parallel SK_{MAC} \parallel SK_{ENC} \parallel SK_{RMAC} = KDF(Z, len, OtherInfo) C7
                                                                                                                    Zeroize Z
                                                                                                                                                                                                 C8
                                                                                                                    AuthCryptogram_{ICC} =
                                                                                                                                                                                                 C9
                                                                                                                            CMAC(SK_{CRFM}, "KC_1_V" || ID_{sICC} || ID_{sH} || Q_{eH})
                                                                                                                    Zeroize SK<sub>CFRM</sub>
                                                                                                                                                                                                 C10
                                                                                        CB_{ICC} \parallel N_{ICC} \parallel
                                                                                                                    EncGUID = GUID XOR AES(SK_{ENC}, IV)
                                                                                                                                                                                                 C11
                                                                                                                    Return \; CB_{ICC} \parallel N_{ICC} \parallel AuthCrytogram_{ICC} \parallel EncGUID \parallel C_{ICC} ^* \; C12
                                                                                        AuthCryptogram<sub>ICC</sub>
                                                                                        || EncGUID || CICC
Check that CB<sub>ICC</sub> is 0x10
                                                                                H4
ID_{sICC} = T_8(SHA256(C_{ICC}*))
                                                                                H5
Extract Q<sub>sICC</sub> from C<sub>ICC</sub>*
                                                                                H6
Verify that Q<sub>sICC</sub> is a valid public key for the domain
                                                                                H7
parameters specified in the response to the SELECT command
Z = ECC\_CDH(d_{eH}, Q_{sICC})
                                                                                H8
Zeroize d<sub>eH</sub>
                                                                                H9
SK_{CFRM} \parallel SK_{MAC} \parallel SK_{ENC} \parallel SK_{RMAC} = KDF(Z, len, OtherInfo)
                                                                                H10
Zeroize Z
                                                                                H11
Check that AuthCryptogram<sub>ICC</sub> equals
                                                                                H12
          CMAC(SK_{CFRM}, \text{``KC\_1\_V''} \parallel ID_{sICC} \parallel ID_{sH} \parallel Q_{eH})
Zeroize SK<sub>CFRM</sub>
                                                                                H13
GUID = EncGUID XOR AES(SK_{ENC}, IV)
                                                                                H14
Build CICC from CICC* and GUID
                                                                                H15
Verify C<sub>ICC</sub> signature
                                                                                H16
```

Sections 4.1.1 and 4.1.2 provide additional details about each of the protocol steps performed by the client application and the PIV Card Application, and Section 4.1.3 defines the notations used in the description of the protocol. Section 4.1.4 provides the details of the two cipher suites that may be supported by the PIV Card Application. Section 4.1.5 specifies the format for the card verifiable certificate (CVC) that is used to authentication the PIV Card Application. Section 4.1.6 provides additional information about the

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key derivation function (KDF) used to derive the session keys that are used during secure messaging, and Section 4.1.7 provides additional information about the computation of the authentication cryptogram for key confirmation. Section 4.1.8 demonstrates the use of the GENERAL AUTHENTICATE to perform the key establishment protocol.

## 4.1.1 Client Application Steps

Step #	Description	Comment
H1	Set CB <sub>H</sub> to 0x10	The client application's control byte is set to 0x10 to indicate the client application does not support persistent binding, wants the GUID returned in encrypted form, and wants 3 session keys to be generated.
H2	Generate an ephemeral key pair (d <sub>eH</sub> ; Q <sub>eH</sub> )	Generate an ephemeral ECC key pair for the client application. If the 0xAC tag of the application property template (APT) includes '27' then generate an ephemeral key pair over Curve P-256. If the 0xAC tag of the APT includes '2B' then generate an ephemeral key pair over Curve P-384.
H3	Send $ID_{sH} \parallel Q_{eH} \parallel CB_H$	
	r response from PIV Card Application: $CB_{ICC} \parallel N_{ICC} \parallel AuthCryptogram_{ICC} \parallel EncGUID \parallel C_{ICC}^*$	
H4	Check that CB <sub>ICC</sub> is 0x10	Verify that the card executed the protocol in accordance with the parameters specified in Step H1.
Н5	$ID_{sICC} = T_8(SHA256(C_{ICC}^*))$	$ID_{sICC}$ , the left-most 8 bytes of the SHA-256 hash of $C_{ICC}^*$ , is used as an input for session key derivation.
Н6	Extract Q <sub>sICC</sub> from C <sub>ICC</sub> *	C <sub>ICC</sub> * is a transformation of the PIV Card's CVC, C <sub>ICC</sub> (see Section 4.1.5). C <sub>ICC</sub> * is constructed from C <sub>ICC</sub> by replacing the Subject Identifier of C <sub>ICC</sub> (T=0x5F20, L=16, V=GUID) with (T=0x5F20, L=0), and leaving all other fields of the CVC unchanged, including the DigitalSignature object.
H7	Verify that $Q_{\text{sICC}}$ is a valid public key for the domain parameters specified in the response to the SELECT command	Perform public key validation of Q <sub>sICC</sub> , where the domain parameters are those of Curve P-256 if P1 is '27' and those of Curve P-384 if P1 is '2B'.
Н8	$Z = ECC\_CDH (d_{eH}, Q_{sICC})$	Compute the shared secret, Z, using the ECC CDH primitive [SP800-56A, Section 5.7.1.2].
Н9	Zeroize d <sub>eH</sub>	Destroy the ephemeral private key generated in Step H2.
H10	$SK_{CFRM} \parallel SK_{MAC} \parallel SK_{ENC} \parallel SK_{RMAC} = KDF(Z, len, OtherInfo)$	Compute the key confirmation key and the session keys. See Section 4.1.6.
H11	Zeroize Z	Destroy the shared secret generated in Step H8.

Step #	Description	Comment
H12	$Check \ that \ Auth Cryptogram_{ICC} \ equals \\ CMAC(SK_{CFRM}, "KC\_1\_V" \parallel ID_{sICC} \parallel ID_{sH} \parallel Q_{eH})$	Perform key confirmation by verifying the authentication cryptogram as described in Section 4.1.7. Return authentication error if verification fails.
H13	Zeroize SK <sub>CFRM</sub>	Destroy the key confirmation key derived in Step H10.
H14	$GUID = EncGUID \textbf{XOR} AES(SK_{ENC}, IV)$	Decrypt GUID. IV is a 16-byte constant consisting of '80' followed by 15 bytes of '00'.
H15	Build $C_{ICC}$ from $C_{ICC}^*$ and $GUID$	Replace the empty Subject Identifier $(T=0x5F20, L=0)$ in $C_{ICC}^*$ with $(T=0x5F20, L=16, V=GUID)$ to create $C_{ICC}$ .
H16	Verify C <sub>ICC</sub> signature	Verify signature on C <sub>ICC</sub> and, using standards-compliant PKI path validation, validate the content signing certificate needed to verify the signature on C <sub>ICC</sub> . Return authentication error if verification fails.

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4.1.2 PIV Card Application Protocol Steps

Step #	Description	Comment
C1	$ID_{sICC} = T_8(SHA256(C_{ICC}^*))$	$ID_{sICC}$ , the left-most 8 bytes of the SHA-256 hash of $C_{ICC}^*$ is used as an input for session key derivation. See Step H6 for construction of $C_{ICC}^*$ (Note that $ID_{sICC}$ and $C_{ICC}^*$ are static, and so may be precomputed off card.)
C2	$CB_{ICC} = CB_H \& 'F0'$	Create the PIV Card Application's control byte from client application's control byte, indicating that persistent binding has not been used in this transaction, even if CB <sub>H</sub> indicates that the client application supports it. This may be done by setting CB <sub>ICC</sub> to the value of CB <sub>H</sub> and then setting the 4 least significant bits of CB <sub>ICC</sub> to 0.
СЗ	Check that CB <sub>ICC</sub> is 0x10	Check that client application is requesting that the GUID be returned in encrypted form and that 3 session keys be generated.
C4	Verify that $Q_{\text{eH}}$ is a valid public key for the domain parameters of $Q_{\text{sICC}}$	Perform public key validity of $Q_{eH}$ , where the domain parameters are those of $Q_{sICC}$ . Also verify that P1 is '27' if the domain parameters of $Q_{sICC}$ are those of Curve P-256 or that P1 is '2B' if the domain parameters of $Q_{sICC}$ are those of Curve P-384.

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Step #	Description	Comment
C5	$Z = ECC\_CDH (d_{sICC}, Q_{eH})$	Compute the shared secret, Z, using the
		ECC CDH primitive [SP800-56A, Section
		5.7.1.2].
C6	Generate nonce N <sub>ICC</sub>	Create a random nonce, where the length is
		as specified in Table 13.
C7	$SK_{CFRM} \parallel SK_{MAC} \parallel SK_{ENC} \parallel SK_{RMAC} =$	Compute the key confirmation key and the
	KDF (Z, len, Otherinfo)	session keys. See Section 4.1.6.
C8	Zeroize Z	Destroy shared secret generated in Step C5.
C9	AuthCryptogram <sub>ICC</sub> =	Compute the authentication cryptogram for
	$CMAC(SK_{CFRM}, "KC\_1\_V" \parallel ID_{sICC} \parallel ID_{sH} \parallel Q_{eH})$	key confirmation as described in Section
		4.1.7.
C10	Zeroize SK <sub>CFRM</sub>	Destroy the key confirmation key derived
		in Step C7.
C11	EncGUID = GUID <b>XOR</b> AES( $SK_{ENC}$ , IV)	Encrypt GUID, which is the Subject
		Identifier of C <sub>ICC</sub> , the PIV Card's CVC. IV
		is a 16-byte constant consisting of '80'
		followed by 15 bytes of '00'.
C12	Return CB <sub>ICC</sub>    N <sub>ICC</sub>    AuthCryptogram <sub>ICC</sub>	
	EncGUID    C <sub>ICC</sub> *	

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## 4.1.3 Notations

Name	Comment	Format	Size (in bytes)
ICC	Integrated Circuit Card (PIV Card)	N/A	N/A
$ID_{sICC}$	Static, non-anonymous PIV Card identifier, which is	Binary	8 bytes
	the truncated hash of C <sub>ICC</sub> *	-	
GUID	Card UUID (see Section 3.4.1 of Part 1)	Binary	16 bytes
CVC or C <sub>ICC</sub>	Card verifiable certificate, which is authenticated by	CVC	
	client application. See Section 4.1.5.		
$C_{ICC}^*$	Confidential card verifiable certificate for privacy,	CVC	
	derived from $C_{ICC}$ as follows: The Subject Identifier		
	data element of $C_{ICC}$ (T=0x5F20, L=16, V=GUID) is		
	replaced with (T=0x5F20, L=0). All other data		
	elements, including the DigitalSignature object, and		
	their order are identical to those in $C_{ICC}$ .		
$ID_{sH}$	Client application identifier. This is a locally	Binary	8 bytes
	assigned identifier for the client application. If none		
	is available, it could be set to all zeros.		
$N_{ICC}$	PIV Card Application nonce. See Table 13 for the	Binary	16 or 24 bytes
	length.		
$SK_{CFRM}$	Key confirmation key used to compute authentication		16 or 32 bytes
	cryptogram.		
$SK_{MAC}$ , $SK_{RMAC}$ ,	Secure messaging session keys. See Table 13 for		16 or 32 bytes
$SK_{ENC}$	encryption or MAC session key length.		
$T_8(Data)$	Leftmost 8 bytes of <i>Data</i> .	Binary	8 bytes
$T_{16}(Data)$	Leftmost 16 bytes of <i>Data</i> .	Binary	16 bytes
KDF(Z, len,	Key Derivation Function (KDF) specified in Section	N/A	N/A
OtherInfo)	4.1.6.		

Name	Comment	Format	Size (in bytes)
ECC_CDH	Elliptic curve cryptography cofactor Diffie-Hellman	N/A	N/A
	(ECC CDH) primitive, as specified in [SP800-56A,		
	Section 5.7.1.2].		
OtherInfo	Input parameters to the KDF function. See Section	N/A	N/A
	4.1.6.		
len	The length (in bits) of the secret keying material to	N/A	N/A
	be generated using the KDF ( $len = 512$ for cipher		
	suite 2 and 1024 for cipher suite 4).		
$CB_{ICC}$	Protocol control byte returned by the PIV Card	Binary	1 byte
$CB_H$	Protocol control byte sent by client application (host)	Binary	1 byte

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## 4.1.4 Cipher Suite

This document specifies two cipher suites (see Table 13) that may be used for key establishment and secure messaging, one that provides 128 bits of channel strength and one that provides 192 bits of channel strength. If the PIV Card Application supports the VCI and either the digital signature key ('9C'), the key management key ('9D'), or one of the retired key management keys ('82' – '95') is an ECC (Curve P-384) key, then PIV Card Application shall only support cipher suite CS4. Otherwise, the PIV Card Application may support either CS2 or CS4.

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Table 13. Cipher Suite for PIV Secure Messaging

	128 bit channel strength	192 bit channel strength
Cipher Suite ID	CS2	CS4
Algorithm Identifier (P1)	'27'	'2B'
	AES 128	AES 256
C <sub>ICC</sub> signature	ECDSA with SHA-256 using an ECDSA (Curve P-256) key	ECDSA with SHA-384 using an ECDSA (Curve P-384) key
C <sub>ICC</sub> public key	ECDH (Curve P-256)	ECDH (Curve P-384)
KDF hash	SHA-256	SHA-384
Nonce (N <sub>ICC</sub> )	16 bytes	24 bytes

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#### 4.1.5 Card Verifiable Certificate

- 712 Table 14 specifies the format for the CVC,  $C_{ICC}$ .
- $C_{ICC}$  is used to authenticate the PIV Card Application. The specific data object tags and specified order must be used to allow the CVC processing within authentication protocols. The specific data object tags
- 715 for  $C_{ICC}$  are provided in Table 14.

Table 14. Card Verifiable Certificate Format

Tag	Tag	Tag	Length	Name	Value
0x7F21				Card Verifiable	
				Certificate	
	0x5F29		1	Credential Profile	0x80
				Identifier	
	0x42		8	Issuer Identification	The leftmost 8 bytes of the
				Number	subjectKeyIdentifier in the content
					signing certificate needed to verify the
					signature on $C_{ICC}$ .
	0x5F20		16	Subject Identifier	GUID (Card UUID) [In C <sub>ICC</sub> *, the length
					of the Subject Identifier is 0.]
	0x7F49		Variable	CardHolderPublicKey	
				Data Object	
		0x06	Variable	Algorithm OID	Possible values are:
					<ul> <li>0x2A8648CE3D030107 for ECDH</li> </ul>
					(Curve P-256) or
					• 0x2B81040022 for ECDH (Curve
					P-384)
		0x86	Variable	Public Key object	Coded as follows: $04 \parallel X \parallel Y$ , where X
					and Y are the coordinates of the point on
					the curve. See the "Value" column of
					Table 12.
	0x5F4C		1	Role Identifier	0x00 for card-application key CVC
	0x5F37		Variable	DigitalSignature object	signatureAlgorithm AlgorithmIdentifier, signatureValue BIT STRING  }  AlgorithmIdentifier ::= SEQUENCE {
					algorithm OBJECT IDENTIFIER, parameters ANY DEFINED BY algorithm OPTIONAL }
					algorithm is 1.2.840.10045.4.3.2 for ECDSA with SHA-256 (cipher suite 2) and 1.2.840.10045.4.3.3 for ECDSA with SHA-384 (cipher suite 4). For both algorithms, the parameters field is absent.
					signatureValue is the DER encoding of signature result ECDSA-Sig-Value defined below.
					ECDSA-Sig-Value ::= SEQUENCE {     r

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- 718 The signature of the CVC (DigitalSignature object) is calculated over the concatenation of the TLV
- encoded Credential Profile Identifier, Issuer Identification Number, Subject Identifier,
- 720 CardHolderPublicKey Data Object, and Role Identifier, i.e., { '5F29' '01' '80' } || { '42' '08' IIN } || { '5F20'
- 721 '10' GUID } || { '7F49' L1 { { '06' L2 OID } { '86' L3 '04' X Y } } } } { '5F4C' '01' '00' }.

## 4.1.6 Key Derivation

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- The session keys shall be derived in Steps C7 and H10 of the protocol using the key derivation function
- from [SP800-56A, Section 5.8.1], with the auxiliary function H being the hash function specified as the
- KDF hash in Table 13 and *OtherInfo* being constructed using the concatenation format as show below:

Cipher Suite ID	OtherInfo
CS2	$0x04 \parallel 0x09 \parallel 0x09 \parallel 0x09 \parallel 0x09 \parallel 0x09 \parallel 0x08 \parallel ID_{sH} \parallel 0x10 \parallel T_{16}(Q_{eH}) \parallel 0x08 \parallel ID_{sICC} \parallel 0x10 \parallel N_{ICC} \parallel 0x10$
CS4	$0x04 \parallel 0x0D \parallel 0x0D \parallel 0x0D \parallel 0x0D \parallel 0x0D \parallel 0x0B \parallel ID_{sH} \parallel 0x10 \parallel T_{16}(Q_{eH}) \parallel 0x08 \parallel ID_{sICC} \parallel 0x18 \parallel N_{ICC} \parallel 0x18 \parallel 0x1$

## 4.1.7 Key Confirmation

- Key confirmation shall be performed in Steps C9 and H12 of the protocol in accordance with [SP800-
- 56A] by the generation of AuthCryptogram<sub>ICC</sub>. AuthCryptogram<sub>ICC</sub> shall be computed as
- 730 CMAC(MacKey, MacLen, MacData<sub>p</sub>), where MacKey is SK<sub>CFRM</sub>, MacLen is 128 bits, and MacData<sub>p</sub> is
- "KC\_1\_V"  $\parallel$  ID<sub>sICC</sub>  $\parallel$  ID<sub>sH</sub>  $\parallel$  Q<sub>eH</sub>. For Q<sub>eH</sub>, the coordinates of the ephemeral public key are converted from
- field elements to byte strings as specified in [SP800-56A, Appendix C.2], Field-Element-to-Byte String
- Conversion, and concatenated (with x first) to form a single byte string. CMAC is cipher-based message
- authentication code from [SP800-38B], where the block cipher is AES.

#### 4.1.8 Command Interface

The following command interface shall be used for the key establishment protocol.

## 737 Command Syntax

CLA	'00'
INS	'87'
P1	Algorithm reference ('27' or '2B'), as specified in the 0xAC tag of the application property template
P2	'03' (PIV Secure Messaging key).
Lc	Length of data field
Data Field	'81' L1 { CB $_{\rm H}$    ID $_{\rm SH}$    Q $_{\rm eH}$ } '82 00', where CB $_{\rm H}$ is 0x10, ID $_{\rm SH}$ is an 8-byte client application identifier as described in Section 4.1.3, and Q $_{\rm eH}$ is an ephemeral public key encoded as 04    X    Y, as specified in the "Value" column of Table 12.
Le	Absent

#### Response Syntax

Data Field	'82' LL { CB <sub>ICC</sub>    N <sub>ICC</sub>    AuthCryptogram <sub>ICC</sub>    EncGUID    C <sub>ICC</sub> * }
SW1-SW2	Status word

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SW1	SW2	Meaning
'61'	'xx'	Successful execution where SW2 encodes the number of response data bytes still available
'6A'	'80'	Incorrect parameter in command data field
'6A'	'86'	Incorrect parameter in P1 or P2
'90'	'00'	Successful execution

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## 4.2 Secure Messaging

743 PIV secure messaging is used to protect the integrity and confidentiality of the PIV data being transmitted 744 between the card and the relying system. PIV secure messaging shall be provided using symmetric 745 session keys derived using the key establishment protocol defined Section 4.1.

Once session keys are established and the card is authenticated as specified in Section 4.1, subsequent

- communication with the card can be performed using secure messaging by setting bits b3 and b4 of the 747 748 CLA byte of the command APDU to 1, resulting in a '0C' or '1C' CLA byte. If bits b3 and b4 of the CLA 749 byte are set, then both the command and the response shall be encrypted and integrity protected as 750 described in this section. If the PIV Card Application cannot encrypt and integrity protect the response 751 (e.g., because it does not support secure messaging or no session keys have been established), the PIV
- Card Application shall return an error (see Section 4.2.7). In the case of command chaining, if bits b3 and 752 753 b4 of the CLA are set in any command in the chain then they shall be set in every command in the chain.
- 754 When secure messaging is used, the data field of the card command (or response) is encrypted first and 755 then a message authentication code (MAC) is applied to the entire command (or response). When command (or response) chaining is required, the encryption and MAC are applied to the entire message 756 757 and the result is then fragmented into separate command (or response) data fields.
- 758 In order to ensure that message reordering or replay attacks can be detected, a 16-byte MAC chaining 759 value (MCV) is used. For the first command, and for the first response, sent after successful completion 760 of the key establishment protocol the MCV consists of 16 bytes of '00'. For each subsequent command 761 the MCV is the 16-byte MAC value computed on the previous command, and for each subsequent 762 response the MCV is the 16-byte MAC value computed on the previous response. The MCV is included
- 763 as part of the message over which the MAC value for each command (or response) is computed.
- 764 The SK<sub>ENC</sub> session key shall be used to encrypt the command data field and response data field as 765 described in Section 4.2.2. The SK<sub>MAC</sub> session key shall be used to add integrity to the command as described in Section 4.2.3. The  $SK_{RMAC}$  session key shall be used to add integrity to the response as 766 described in Section 4.2.5. 767
- 768 Secure messaging specified in this section can be applied to the following commands:
- 769 + GET DATA 770
- 771 CHANGE REFERENCE DATA

VERIFY

772 GENERAL AUTHENTICATE

## 4.2.1 Secure Messaging Data Objects

The command and response messages shall be BER-TLV encoded according to Table 15.

**Table 15. Secure Messaging Data Objects** 

Tag	Description
'87'	Padding-content indicator byte followed by the encrypted data
'8E'	Cryptographic checksum (MAC)
'97'	Le
'99'	Status word

## 4.2.2 Command and Response Data Confidentiality

Under secure messaging, the PIV data is encrypted using AES in Cipher Block Chaining (CBC) mode with the  $SK_{ENC}$  session key, where  $SK_{ENC}$  is a 128-bit key for CS2 and a 256-bit key for CS4 as per Table 13. The encryption and encoding process for command data and response data shall be the same. The encryption of the command data or response data and encoding in BER-TLV format is illustrated Figure 1. The encryption shall be computed over the entire message before applying fragmentation for data transportation.

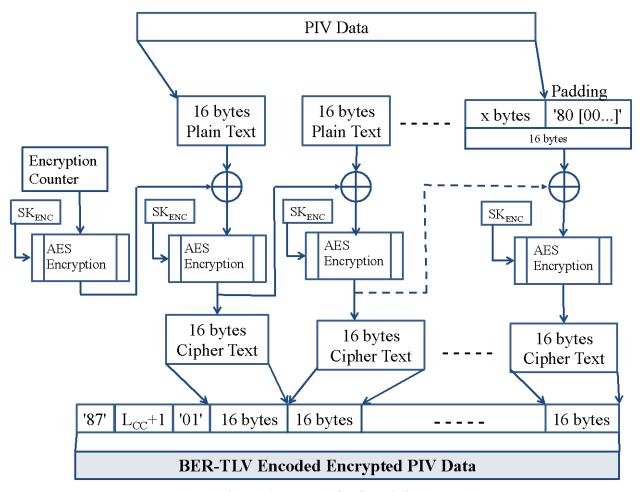


Figure 1. PIV Data Confidentiality

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- 786 Initialization Vector (IV): The IV for the AES CBC encryption of command data shall be generated by
- applying the AES block cipher to a 16-byte encryption counter. The initial value of the encryption
- counter upon successful completion of the key establishment protocol shall be '00 00 00 00 00 00 00 00 00
- 789 00 00 00 00 00 00 00 01'. The encryption counter shall be incremented by one after each creation of an
- 790 IV to encrypt command data, and it shall be reset to its initial value after each successful completion of
- the key establishment protocol. The 16-byte IV shall be created by encrypting the encryption counter
- with SK<sub>ENC</sub> using AES in the electronic codebook (ECB) mode of operation.
- The IV for the AES CBC encryption of response data shall also be generated by encrypting an encryption
- counter with  $SK_{ENC}$  using AES in the ECB mode of operation. The encryption counter value used to
- generate the IV to encrypt the response data shall be the same as the encryption counter value used to
- generate the IV to encrypt the corresponding request data, with the exception that the most significant
- byte of the 16-byte counter shall be set to '80' (i.e., the IV used to encrypt the first response after
- successful completion of the key establishment protocol shall be generated by encrypting '80 00 00 00 00
- 799 00 00 00 00 00 00 00 00 00 01' with SK<sub>ENC</sub>).
- Padding: If the length of the command or response data is not a multiple of 16 bytes then padding shall
- be added to the last block of input data. The padding shall be '80' followed by the number of zeros
- needed to make up the length of 16 byte input block. If padding is used, the first byte of the value field of
- tag '87' shall be '01'; otherwise, the first byte shall be '02'.
- As illustrated in Figure 1, the input and output of encryption is as follows:
  - Encryption input:

Plain Text

#### • Encryption output:

BER-TLV encoded encrypted message, which consists of tag '87' followed by the length of the encoded encrypted message ( $L_{cc} + 1$ ), the padding indicator byte ('01' or '02'), and then the encrypted data.  $L_{cc}$  is the length of the encrypted PIV data; it shall be a multiple of 16.

## 4.2.3 Command Integrity

- The Command MAC (C-MAC) shall be generated by applying the cipher-based MAC (CMAC)
- [SP800-38B] to the header and data field of a command using the SK<sub>MAC</sub> session key. In the case that
- fragmentation is required for data transmission, the command shall be constructed without fragmentation
- for the purposes of computing the MAC, and the CLA byte used in the computation of the MAC shall be
- 818 '0C'.

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- The data to be MACed,  $M_{C,MAC}$ , shall be constructed by concatenating the following:
- 1. The 16-byte MAC chaining value (MCV). For the first command sent after successful completion of the key establishment protocol the MCV consists of 16 bytes of '00'. For each subsequent command the MCV is the 16-byte MAC value computed for the previous command.
- 2. A 16-btye encoded header. The encoded header shall consist of the CLA byte ('0C'), the INS byte, P1, and P2, followed by twelve bytes of padding, consisting of '80' followed eleven bytes of '00'. (The length of the data field, L<sub>c</sub>, is not included in the data to be MACed.)

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- 3. The data field, which is the BER-TLV encoded encrypted message.<sup>7</sup>
- 4. L<sub>e</sub> encapsulated in BER-TLV format with tag '97', if the L<sub>e</sub> field is included in the command.

Let  $T_{C-MAC} = \text{CMAC}(\text{SK}_{\text{MAC}}, M_{C-MAC})$  as described in [SP800-38B]. The BER-TLV encoded C-MAC for the command shall be the 8 most significant bytes of  $T_{C-MAC}$  encapsulated in BER-TLV format with tag

- '8E'. The entire 16-byte value  $T_{C-MAC}$  will be the MCV for the next command.
- Figure 2 below illustrates how the C-MAC is generated for each command.

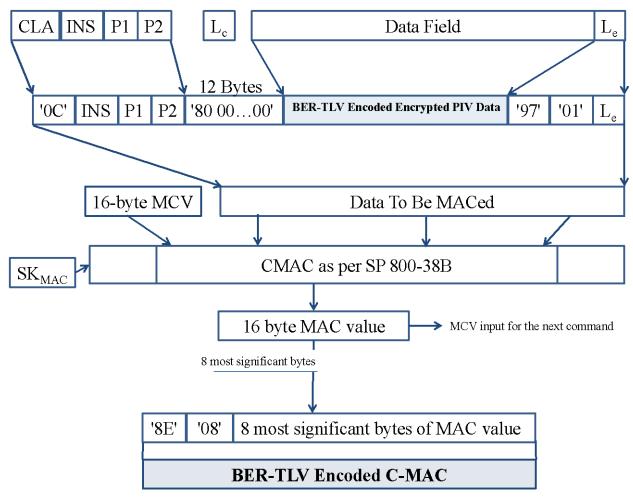


Figure 2. PIV Data Integrity of Command

## 4.2.4 Command with PIV Secure Messaging

For secure messaging, the secure messaging data field shall be constructed as the concatenation of the following: the BER-TLV encoded encrypted PIV data;  $^8$  the 3-btye BER-TLV encoded  $L_e$ , as described in Section 4.2.3, if  $L_e$  would have been included in a message sent without secure messaging; and the 10-byte BER-TLV encoded C-MAC of the command, as described in Section 4.2.3.

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<sup>&</sup>lt;sup>7</sup> The data field may be absent in the case of the VERIFY command.

<sup>&</sup>lt;sup>8</sup> The data field may be absent in the case of the VERIFY command.

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The APDU for secure messaging is shown in Figure 3 for the case in which command chaining is not required. The APDU consists of the CLA byte ('0C'), INS, P1, P2, the length of the secure messaging data field (L<sub>c</sub>'), and the secure messaging data field.

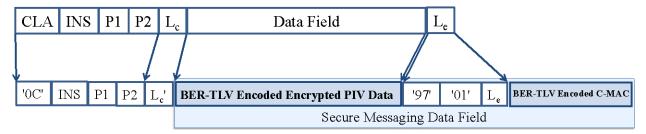


Figure 3. Single Command under Secure Messaging

If secure messaging data field to be transported is larger than 255 bytes, command chaining will be needed. Figure 4 shows the APDUs for secure messaging for a case in which the length of the secure messaging data field is between 256 and 510 bytes, requiring the data to be fragmented across two APDUs. The APDUs are constructed in the same manner as when fragmentation is not required, except that the CLA byte for the first APDU is '1C', the first APDU contains the first 255 bytes of the secure messaging data field, and the second APDU contains the remaining bytes of the secure messaging data field. The PIV Card Application provides a two-byte response of '90 00' for the first APDU. After receiving the second APDU the PIV Card Application reconstructs and processes the entire command.

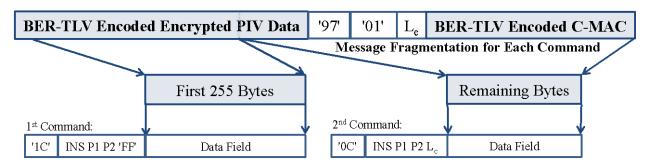


Figure 4. Chained Command under Secure Messaging

## 4.2.5 Response Integrity

The Response MAC (R-MAC) shall be generated by applying CMAC [SP800-38B] to the data field and status bytes of the response using the  $SK_{RMAC}$  session key. An R-MAC shall be generated for each response that corresponds to a command that was sent to the card using secure messaging.

The data to be MACed,  $M_{R-MAC}$ , shall be constructed by concatenating the following:

- 1. The 16-byte MAC chaining value (MCV). For the first response sent after successful completion of the key establishment protocol the MCV consists of 16 bytes of '00'. For each subsequent response the MCV is the 16-byte MAC value computed for the previous response.
- 2. The data field (if present), which is the BER-TLV encoded encrypted message.
- 3. The status words, SW1 and SW2, encapsulated in BER-TLV format with tag '99'.

Let  $T_{R-MAC} = \text{CMAC}(\text{SK}_{\text{RMAC}}, M_{R-MAC})$  as described in [SP800-38B]. The BER-TLV encoded R-MAC for the response shall be the 8 most significant bytes of  $T_{R-MAC}$  encapsulated in BER-TLV format with tag '8E'. The entire 16-byte value  $T_{R-MAC}$  will be the MCV for the next response.

Figure 5 below illustrates how the R-MAC is generated for the response.

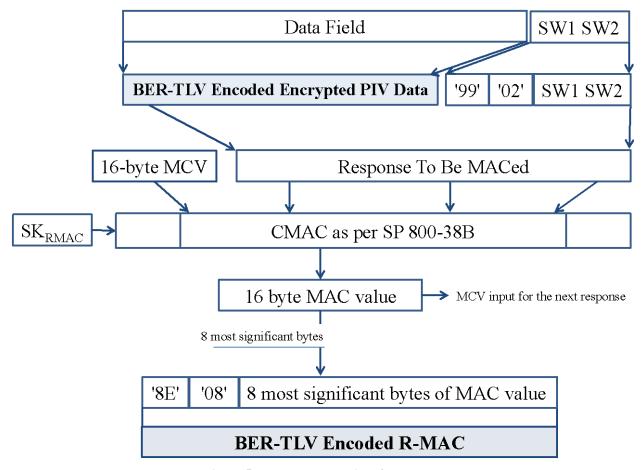


Figure 5. PIV Data Integrity of Response

## 4.2.6 Response with PIV Secure Messaging

For secure messaging, the secure messaging data field that is sent by the PIV Card Application shall be constructed as the concatenation of the following: the BER-TLV encoded encrypted message (when present); the 4-byte BER-TLV encoded the status words, as described in Section 4.2.5; and the 10-byte BER-TLV encoded R-MAC of the response, as described in Section 4.2.5.

Figure 6 illustrates a response under secure messaging for the case in which response chaining is not required. The APDU consists of the secure messaging data field and the 2-byte SW protocol ('90 00'), which indicates that the PIV Card Application successfully verified the C-MAC on the command and decrypted the data field in the command (if present). If the PIV Card Application was unable to verify the C-MAC on the command or decrypt the data field in the command, then it shall return a 2-byte error response, as described in Section 4.2.7.

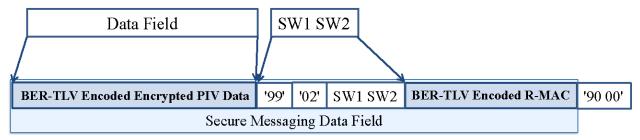


Figure 6. Single Response under Secure Messaging

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If the secure messaging data field to be transported is larger than 256 bytes, response chaining<sup>9</sup> will be needed. Figure 7 shows the APDUs for secure messaging that are sent by the PIV Card Application for a case in which the length of the secure messaging data field is between 513 and 768 bytes, requiring the data to be fragmented across three APDUs. After the first response an APDU of '00 C0 00 00 00' would be sent to request the second response, and after the second response an APDU of '00 C0 00 00 xx' would be sent to request the third response.

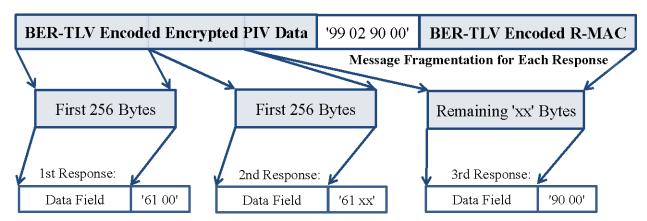


Figure 7. Chained Response under Secure Messaging

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## 4.2.7 Error Handling

The SW protocol is the status byte of the overall secure messaging command and response processing. It indicates if the secure messaging was performed successfully. If the processing was successful, it shall be '90 00'; otherwise, it shall be as follows:

- + '68 82' Secure messaging not supported
- + '69 82' Security condition not satisfied<sup>10</sup>
- + '69 87' Expected secure messaging data objects are missing
- + '69 88' Secure messaging data objects are incorrect

If the command processing was unsuccessful, the card shall return one of the above errors without performing further secure messaging.

<sup>&</sup>lt;sup>9</sup> The response chaining is accomplished by issuing several GET RESPONSE commands to the card.

<sup>&</sup>lt;sup>10</sup> Status word '69 82' is used when secure messaging is requested, but no session keys have been established.

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905	4.3 Sessi	on Key Destruction
906 907		teys established after successful execution of the key establishment protocol in Section 4.1 zed in the following circumstances:
908	+	the card is reset;
909	+	an error occurs in secure messaging; or
910	+	new session keys are requested by the client application by sending a GENERAL
911		AUTHENTICATE command to the card to perform the key establishment protocol using
912		the PIV Secure Messaging key.
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## 915 A

## Appendix A—Examples of the Use of the GENERAL AUTHENTICATE Command

## A.1 Authentication of the PIV Card Application Administrator

- 917 The PIV Card Application Administrator is authenticated by the PIV Card Application using a
- challenge/response protocol. A challenge retrieved from the PIV Card Application is encrypted by the
- 919 client application and returned to the PIV Card Application associated with key reference '9B', the key
- 920 reference of the PIV Card Application Administration key. The PIV Card Application decrypts the
- response using this reference data and the algorithm associated with the key reference (for example, 3
- 922 Key Triple DES ECB, algorithm identifier '00'). If this decrypted value matches the previously
- provided challenge, then the security status indicator of the PIV Card Application Administration key is
- 924 set to TRUE within the PIV Card Application.

Table 16 shows the GENERAL AUTHENTICATE card commands sent to the PIV Card Application to realize this particular challenge/response protocol.

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#### Table 16. Authentication of PIV Card Application Administrator

Command	Response	Comment
'00 87 00 9B 04 7C 02 81		Client application requests a
00'		challenge from the PIV Card
		Application.
	'7C 0A 81 08 01 02 03 04	Challenge ('01 02 03 04 05 06 07
	05 06 07 08 90 00'	08') returned to client application
		by the PIV Card Application.
'00 87 00 9B 0C 7C 0A 82		Client application returns the
08 88 77 66 55 44 33 22 11'		encryption of the challenge ('88
		77 66 55 44 33 22 11')
		referencing algorithm '00' and
		key reference '9B'. [SP800-78,
		Tables 6-1 and 6-2]
	'90 00'	PIV Card Application indicates
		successful authentication of PIV
		Card Application Administrator
		after decrypting '88 77 66 55 44
		33 22 11' using the referenced
		algorithm and key and getting '01
		02 03 04 05 06 07 08'.

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## A.2 Mutual Authentication of Client Application and Card Application

The PIV Card Application Administrator and the PIV Card Application authenticate each other using a challenge/response protocol. A witness retrieved from the PIV Card Application is decrypted by the client application and returned to the PIV Card Application associated with key reference '9B', the key reference of the PIV Card Application Administration key. The command including the decrypted witness also includes a challenge for the PIV Card Application. The PIV Card Application verifies that the decrypted witness matches the value that it encrypted to create the witness. If it does, then the security status indicator of the PIV Card Application Administration key is set to TRUE within the PIV Card Application, and the PIV Card Application encrypts the challenge that it received from the client

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application and returns the result. The witness and challenge are encrypted/decrypted using the same the key and algorithm. Table 18 shows the GENERAL AUTHENTICATE card commands sent to the PIV Card Application to realize mutual authentication using 3 Key Triple DES – ECB (algorithm identifier '00').

Table 17. Mutual Authentication of Client Application and PIV Card Application

Command	Response	Comment
'00 87 00 9B 04 7C 02 80		Client application requests a
00'		witness from the PIV Card
		Application.
	'7C 0A 80 08 88 77 66 55	PIV Card Application returns a
	44 33 22 11 90 00'	witness that is created by
		generating 8 bytes of random
		data ('01 02 03 04 05 06 07 08')
		and encrypting it using the
		referenced key ('9B') and
		algorithm ('00'). [SP800-78,
		Tables 6-1 and 6-2]
'00 87 00 9B 18 7C 16 80		Client application returns the
08 01 02 03 04 05 06 07 08		decrypted witness ('01 02 03 04
81 08 09 0A 0B 0C 0D 0E		05 06 07 08') referencing
0F 10 82 00'		algorithm '00' and key reference
		'9B'. Client application requests
		encryption of challenge data ('09
		0A 0B 0C 0D 0E 0F 10') from
		the card using the same key.
	'7C 0A 82 08 11 FF EE DD	PIV Card Application
	CC BB AA 99 90 00'	authenticates the client
		application by verifying the
		decrypted witness. PIV Card
		Application indicates successful
		authentication of PIV Card
		Application Administrator and
		sends back the encrypted
		challenge ('11 FF EE DD CC BB
		AA 99'). Client application
		authenticates the PIV Card
		Application by decrypting the
		encrypted challenge and getting
		('09 0A 0B 0C 0D 0E 0F 10').

## A.3 Authentication of PIV Cardholder

The PIV cardholder is authenticated by first retrieving and validating either the X.509 Certificate for PIV Authentication or the X.509 Certificate for Card Authentication. Assuming the certificate is valid, the client application requests the PIV Card Application to sign a challenge using the private key associated with this certificate (i.e., key reference '9A' or '9E') and the appropriate algorithm (e.g., algorithm identifier '07'), which can be determined from the certificate as described in Part 1, Appendix C.1. The

response from the card is verified using the public key in the certificate. If the signature verifies, then the PIV cardholder is authenticated.

Table 17 shows the GENERAL AUTHENTICATE card commands sent to the PIV Card Application to realize the cardholder authentication when the X.509 Certificate for PIV Authentication includes a 2048-bit RSA public key. It is assumed that the cardholder PIN or OCC data has been successfully verified prior to sending the GENERAL AUTHENTICATE command.

Table 18. Validation of the PIV Card Application Using GENERAL AUTHENTICATE

Command	Response	Comment
'10 87 07 9A FF 7C 82 01 06 82		Client application sends a
00 81 82 01 00 00 01 FF FF FF		challenge to the PIV Card
FF FF FF FF FF 00 9D F4		Application indicating the
6E 09 E7 D6 19 18 53 1E 6E 1C		reference data associated with
66 87 C4 3E CF FF 7D 53 47 BD		key reference '9A' is to be used
2E 93 19' ("" represents 208		with algorithm '07'. [SP800-78,
bytes of challenge data)		Tables 6-1 and 6-2] The
bytes of chancinge data)		challenge data, which in this
		example is encoded as specified
		for TLS version 1.1 client
		authentication, is '00 01 FF 18
		BC A7'. Bit 5 of CLA byte is set
		to one indicating command
		chaining is needed. L <sub>e</sub> is absent
		indicating no data is expected.
	'90 00'	PIV Card Application indicates it
		received the command
		successfully.
'00 87 07 9A 0B 94 53 76 FE A7		Client application sends
91 72 14 18 BC A7 00'		remaining data with the second
		and last command of the chain.
		L <sub>e</sub> is '00' to indicate that the
		expected length of the response
		data field is 256 bytes.
	'7C 82 01 04 82 82 01 00 29 69	PIV Card Application returns the
	44 3B 49 AC 5B 70 63 51 A1 5B	result of signing the challenge
	B5 AD F7 0B 7D A6 4C 6C	using the indicated key reference
	AA 62 40 C5 FA A8 7E A2 2B	data and algorithm ('29 69 44 3B
	DC 92 18 56 8B CE F4 69 14 D9	49 AC'). The last two bytes
	83 61 08' ("" represents 208	'61 08' indicate 8 more bytes are
	bytes of response data)	available to read from the card.
'00 C0 00 00 08		The GET RESPONSE command
		is used to request remaining 8
		bytes.
	'30 1B 11 06 AE E2 F1 2E 90 00	PIV Card Application sends the
		remaining 8 bytes.

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## A.4 Signature Generation with the Digital Signature Key

- The GENERAL AUTHENTICATE command can be used to generate signatures. The pre-signature hash
- and padding (if applicable) is computed off card. The PIV Card Application receives the hashed value of
- the original message, applies the private signature key (key reference '9C'), and returns the resulting
- signature to the client application.
- Listed below are the card commands sent to the PIV Card Application to generate a signature. It is
- assumed that the cardholder PIN or OCC data has been successfully verified prior to sending the
- 965 GENERAL AUTHENTICATE command.

## A.4.1 RSA

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- This example illustrates signature generation using RSA 2048 (i.e., algorithm identifier '07'). Command
- chaining is used in the first command since the padded hash value sent to the card for signature generation
- is bigger than the length of the data field.

## Command 1: (GENERAL AUTHENTICATE – first chain):

CLA	'10' indicating command chaining
INS	'87'
P1	'07'
P2	'9C'
Lc	Length of data field
Data Field	'7C' – L1 { '82' '00' '81' L2 {first part of the PKCS #1 v1.5 or PSS padded message hash value }}
Le	Absent (no response expected)

## 971 972 **Response 1:**

Data Field	Absent
SW1-SW2	'90 00' (Status word)

#### Command 2: (GENERAL AUTHENTICATE – last chain):

CLA	'00' indicates last command of the chain
INS	'87'
P1	'07'
P2	'9C'
Lc	Length of data field
Data Field	{second and last part of the PKCS #1 v1.5 or PSS padded message hash value}
Le	Length of expected response

### **Response 2:**

Data Field	'7C' - L1 {'82' L2 {first part of signature} }
SW1-SW2	'61 xx' where xx indicates the number of bytes remaining to send by the PIV Card Application

## 977 Command 3: (GET RESPONSE APDU):

CLA	'00'
INS	'C0'
P1	'00'
P2	'00'
L <sub>e</sub>	xx Length of remaining response as indicated by previous SW1-SW2

## Response 3:

Data Field	{second and last part of signature}
SW1-SW2	'90 00' (Status word)

## A.4.2 ECDSA

The following example illustrates signature generation with ECDSA using ECC: Curve P-256 (i.e., algorithm identifier '11'). Command chaining is not used in this example, as the hash value fits into the data field of the command. Padding does not apply to ECDSA.

## **Command – GENERAL AUTHENTICATE**

CLA	'00'
INS	'87'
P1	'11'
P2	'9C'
L <sub>c</sub>	Length of data field
Data Field	'7C' - L1 { '82' '00' '81' L2 {hash value of message}}
L <sub>e</sub>	Length of expected response

## Response:

Data Field	'7C' – L1 {'82' L2 (r,s)} where  • (r,s) is DER encoded with the following ASN.1 structure:  Ecdsa-Sig-Value ::= SEQUENCE {  r INTEGER,
	s INTEGER }  L1 is the length of tag '82' TLV structure  L2 is the length of the DER encoded Ecdsa-Sig-Value
SW1-SW2	structure  '90 00' (Status word)

## A.5 Key Establishment Schemes with the PIV Key Management Key

FIPS 201 specifies a public key pair and associated X.509 Certificate for Key Management. The key management key (KMK) is further defined in SP 800-78, which defines two distinct key establishment schemes for the KMK:

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- 995 1) RSA key transport and
- 996 2) Elliptic Curve Diffie-Hellman (ECDH) key agreement.
- The use of the KMK for RSA key transport and ECDH key agreement is discussed in Appendices A.5.1 and A.5.2, respectively.

## A.5.1 RSA Key Transport

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- 1000 In general, RSA transport keys are used to establish symmetric keys, where a sender encrypts a symmetric
- key with the receiver's public key and sends the encrypted key to the receiver. The receiver decrypts the
- encrypted key with the corresponding private key. The decrypted symmetric key subsequently is used by
- both parties to protect further communication between them. Many types of security protocols employ
- the RSA key transport technique. S/MIME for secure email is one of the many protocols employing RSA
- transport keys to distribute symmetric keys between entities.

## A.5.1.1 RSA Key Transport with the PIV KMK

- As specified in SP 800-78, the on-card private KMK can be an RSA transport key that complies with
- 1008 [PKCS1]. In the scenario described above, a sender encrypts a symmetric key with the KMK's public
- 1009 RSA transport key. The role of the on-card KMK private RSA transport key is to decrypt the sender's
- symmetric key on behalf of the cardholder and provide it to the client application cryptographic module.

## 1011 A.5.1.1.1 The GENERAL AUTHENTICATE Command

- 1012 Listed below are the card commands sent to the PIV Card to decrypt the symmetric key. It is assumed
- that the cardholder's PIN or OCC data has been successfully verified prior to sending the GENERAL
- 1014 AUTHENTICATE command to the card.

## 1015 Command 1 – GENERAL AUTHENTICATE (first chain)

CLA	'10' indicates command chaining
INS	'87'
P1	'07'
P2	'9D'
Lc	Length of data field
Data Field	'7C' – L1 {'82' '00' '81' L2 {first part of C}} where C is the ciphertext to be decrypted, as defined in [PKCS1, Sections 7.1.2 and 7.2.2]
Le	Absent (no response expected)

#### 1017 **Response 1:**

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Data Field	Absent
SW1-SW2	'90 00' (Status word)

## 1020 Command 2 – GENERAL AUTHENTICATE (last chain)

CLA	'00' indicates last command of the chain		
INS	'87'		
P1	'07'		
P2	'9D'		
L <sub>c</sub>	Length of data field		
Data Field	{second and last part of ciphertext to be decrypted C }}		
Le	Length of expected response		

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## Response 2:

Data Field	'7C' – L1 {'82' L2 {first part of encoded message EM}} where EM is as defined in [PKCS1, Sections 7.1.2 and 7.2.2]
SW1-SW2	'61 xx' where x indicates the number of bytes remaining to send

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## **Command 3: GET RESPONSE APDU:**

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CLA	'00'	
INS	'C0'	
P1	'00'	
P2	'00'	
Le	xx Length of remaining response as indicated by previous SW1-SW2	

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## Response 3:

Data Field	{second and last part of encoded message EM}
SW1-SW2	'90 00' (Status word)

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## A.5.2 Elliptic Curve Cryptography Diffie-Hellman

An ECDH key agreement scheme does not send an encrypted symmetric key to the participating entities.

Instead, the two entities involved in the key agreement scheme compute a shared secret by combining
their ECC private key(s) with the other party's public key(s). The resulting shared secret (Z) serves as an
input to a key derivation function (KDF), which each entity independently invokes to derive a common
secret key. The secret key may be used as a session key or may be used to encrypt a session key.

## 1036 A.5.2.1 ECDH with the PIV KMK

The PIV Card supports ECDH key agreement by performing the elliptic curve cryptography cofactor
Diffie-Hellman (ECC CDH) primitive [SP800-56A, Section 5.7.1.2] using its ECC KMK private key and
an ECC public key that is provided as input to the GENERAL AUTHENTICATE command. All other
procedures required to complete the key agreement are performed by the cardholder's client application
and its associated cryptographic module.

#### A.5.2.1.1 The GENERAL AUTHENTICATE Command

The sequence of commands to perform the ECC CDH primitive from [SP800-56A, Section 5.7.1.2] with the private ECC KMK is illustrated below for ECC: Curve P-256:

### **Command – GENERAL AUTHENTICATE**

CLA	'00'	
INS	'87'	
P1	'11'	
P2	'9D'	
L <sub>c</sub>	Length of data field	
Data Field	'7C' – L1 {'82' '00' '85' L2 { '04'    X    Y}} , where  • '04'    X    Y is the other party's public key, a point on Curve P-256, encoded without the use of point compression as described in [SECG, Section 2.3.3].  • The length of each coordinate (X and Y) is 32 bytes and  • The value of L2 is 65 bytes	
L <sub>e</sub>	Length of expected response	

### Response:

Data Field	'7C' - L1 {'82' L2 {shared secret Z}} where
	<ul> <li>Z is the X coordinate of point P as defined in [SP800-56A, Section 5.7.1.2]</li> </ul>
	L2 is 32 bytes
SW1-SW2	'90 00' (Status word)

## A.5.2.2 PIV KMK Specific ECDH Key Agreement Schemes

SP 800-56A describes five different ECDH key agreement schemes that a client application cryptographic module may implement. These schemes differ in 1) the number of keys (1 or 2) and 2) the type of keys (ephemeral or static) used by each party. Since the PIV Card only computes the ECC CDH primitive using its static private key, the client application cryptographic module only employs the PIV Card in implementing an ECDH key agreement scheme when the scheme involves the use of the cardholder's static key pair. The ECDH key agreement schemes that involve the use of at least one party's static key pair, and thus may involve the use of the PIV Card are:

+ C(2, 2) – Each party has a static key pair and generates an ephemeral key pair [SP800-56A, Section 6.1.1]

In this scheme, the information sent between the client application and the PIV Card is the same when acting as the initiator or the responder; the other party's static public key is sent to the PIV Card, and a static shared secret is returned by the PIV Card in plaintext. Note that an ephemeral key pair is generated by the client application, and the private key of that key pair is combined with the other party's ephemeral public key to produce an ephemeral shared secret.

+ C(1, 2) – The initiator has a static key pair and generates an ephemeral key pair, while the responder has a static key pair [SP800-56A, Section 6.2.1]

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When the cardholder is acting as the initiator, the other party's static public key is sent to the PIV Card, and a static shared secret is returned in plaintext by the PIV Card. Note that in this case, an ephemeral key pair is generated by the client application's cryptographic module, and the corresponding ephemeral private key is combined with the other party's static public key to produce a second shared secret.

When the cardholder is acting as the responder, two public keys are sent by the client application to the PIV Card (the other party's static and ephemeral public keys), and two shared secrets are returned in plaintext (the static shared secret and the ephemeral shared secret). Note that two GENERAL AUTHENTICATE commands are required to provide the two shared secrets to the client application's cryptographic module.

+ C(1, 1) – The initiator generates only an ephemeral key pair, while the responder has only a static key pair [SP800-56A, Section 6.2.2]

In this scheme, the PIV Card is only employed by the client application if the cardholder is acting as the responder. In this case, the other party's ephemeral public key is sent to the PIV Card, and the shared secret is returned by the PIV Card in plaintext.

+ C(0, 2) – Both the initiator and responder use only static key pairs [SP800-56A, Section 6.3]

In the C(0, 2) scheme, the information sent between the client application's cryptographic module and the PIV Card is the same when acting as the initiator or the responder; the other party's static public key is sent to the PIV Card, and the static shared secret is returned in plaintext. Note that for this scheme, the client application's cryptographic module also generates a nonce when acting as the initiator of the scheme.

The C(2, 0) scheme does not involve the use of static keys and so the PIV Card would not be involved in the implementation of this scheme.

## A.6 Authentication of the PIV Cardholder Over the Virtual Contact Interface

- If the PIV Card supports secure messaging and the pairing code, then all non-card-management operations of the PIV Card Application may be performed over the contactless interface. In order to perform an operation that would otherwise be restricted to the contact interface, the key establishment protocol in Section 4.1 needs to be performed to establish session keys for secure messaging, and then the pairing code needs to be submitted over secure messaging in order to establish a virtual contact interface.
- This appendix shows an example of the establishment of a VCI and its use to perform cardholder
- authentication using the PIV Authentication key. First, the GENERAL AUTHENTICATE command is
- used to perform the key establishment protocol, and then the VERIFY command is used to submit the
- pairing code and establish the VCI. At this point the GET DATA command is used to read the X.509
- 1102 Certificate for PIV Authentication. Then the GENERAL AUTHENTICATE command is used to perform
- a challenge/response with the PIV Authentication key after the PIN is submitted using the VERIFY
- 1104 command.

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Command	Response	Comment
00 87 27 03 4E 81 4A 10 00 00		The GENERAL
00 00 00 00 00 00 04 X Y 82 00		AUTHENTICATE command is
		used to perform the key
		establishment protocol, as
		specified in Section 4.1.8, where
		cipher suite CS2 is being used,
		$ID_{sH}$ is all zeros, and X and Y are
		the coordinates of Q <sub>eH</sub> . X and Y
		are 32 bytes each.
	82 LL 10 N <sub>ICC</sub>	The response for the key
	AuthCryptogram <sub>ICC</sub> EncGUID	establishment protocol, as
	C <sub>ICC</sub> *	specified in Section 4.1.8, where
		N <sub>ICC</sub> , AuthCryptogram <sub>ICC</sub> , and
		EncGUID are 16 bytes each, and
		C <sub>ICC</sub> * is as specified in Sections
		4.1.3 and 4.1.5.
00 00 00 00 00 00 04 X Y 82 00	AuthCryptogram <sub>ICC</sub> EncGUID	AUTHENTICATE command is used to perform the key establishment protocol, as specified in Section 4.1.8, where cipher suite CS2 is being used, ID <sub>sH</sub> is all zeros, and X and Y are the coordinates of Q <sub>eH</sub> . X and Y are 32 bytes each.  The response for the key establishment protocol, as specified in Section 4.1.8, where N <sub>ICC</sub> , AuthCryptogram <sub>ICC</sub> , and EncGUID are 16 bytes each, and C <sub>ICC</sub> * is as specified in Sections 4.1.3 and 4.1.5.

After the client application verifies  $C_{\rm ICC}$  and the authentication cryptogram and validates the content signing certificate needed to verify the signature on  $C_{\rm ICC}$ , the PIV Card has been authenticated and session keys for secure messaging have been established ( $SK_{\rm ENC}$ ,  $SK_{\rm MAC}$ , and  $SK_{\rm RMAC}$ ).

0C 20 00 98 1D 87 11 01 ENC <sub>C1</sub>		The VERIFY command is used
$8E\ 08\ T_8(T_{C-MAC,1})$		over secure messaging to submit
		the pairing code to the card.
	99 02 90 00 8E 08 T <sub>8</sub> (T <sub>R-MAC,1</sub> )	The card responds that the
	90 00	command has been successfully
		executed, and that the VCI has
		been established.

Once the VCI has been established, the GET DATA command may be used to retrieve the X.509 Certificate for PIV Authentication.

0C CB 3F FF 20 87 11 01 ENC <sub>C2</sub>	The GET DATA command is
97 01 00 8E 08 T <sub>8</sub> (T <sub>C-MAC,2</sub> )	used to request the X.509
	Certificate for PIV
	Authentication. The command is
	submitted over VCI.

# Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface

Command	Response	Comment
	87 82 05 91 01 bytes 1 – 251 of	The response includes the tag,
	$ENC_{R2} > 61\ 00$	length, and padding indicator
	LIVERZE OF OO	bytes of the BER-TLV encoded
		encrypted response data along
		with the first 251 bytes of the
		encrypted response, and an
		indicator that at least 256 bytes
		of additional data is available.
		The padding indicator is '01' to
		indicate that padding was
		required.
0C C0 00 00 00		Request the next 256 bytes of the
00 00 00 00		response.
	   	Return the next 256 bytes of the
	00	response.
0C C0 00 00 A3		Request the final 163 bytes of the
		response.
	   description of the control of the	Return the final 163 bytes of the
	99 02 90 00 8E 08 T <sub>8</sub> (T <sub>R-MAC,2</sub> )	response, including the BER-
	90 00	TLV encoded status words for
		the command and the BER-TLV
		encoded R-MAC.
	ad could be used to submit the PIV C s example, for illustrative purposes of	
	order to retrieve the current value of t	
the PIV Card Application PIV.	radi to formeve the current value of	and reary counter associated with
For the command,		
0C 20 00 80 0A 8E 08		The VERIFY command is used
$T_8(T_{C-MAC,3})$		to retrieve the number of further
- 8\ - C-MAC,3/		retries allowed for the PIV Card
		Application PIN.
	99 02 63 C3 8E 08 T <sub>8</sub> (T <sub>R-MAC,3</sub> )	The PIV Card Application
	90 00	indicates that 3 further retries are
		allowed ('63 C3').
The VERIFY command is used to	submit the PIV Card Application PII	·
	ne PIN value, the command and response	* *
the VERIFY command to submit t	_	
	sult of encrypting the PIN value alon	g with the padding bytes using an
	00 00 00 00 00 00 00 00 00 00 03'),	
	ne encryption counter used to general	·
	nmand since no encryption was perfo	
For the response, T <sub>R-MAC,4</sub> is comp		•
0C 20 00 80 1D 87 11 01 ENC <sub>C3</sub>		The VERIFY command is used
8E 08 T <sub>8</sub> (T <sub>C-MAC.4</sub> )		to submit the PIV Card
		Application PIN to the card.
	99 02 90 00 8E 08 T <sub>8</sub> (T <sub>R-MAC.4</sub> )	The card responds that the
	90 00	command has been successfully
		executed.
	l	L

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# Interfaces for Personal Identity Verification – Part 2: PIV Card Application Card Command Interface

Command	Response	Comment
Now that a virtual contact interface	e has been established and the PIV-C	ard Application PIN has been
	be performed over the contactless in	* *
	ed to perform a challenge/response v	
	sult of encrypting the challenge along	· · · · · · · · · · · · · · · · · · ·
IV of AES(SK <sub>ENC</sub> , '00 00 00 00 00	00 00 00 00 00 00 00 00 00 00 04'),	and T <sub>C-MAC,5</sub> is computed using
$T_{C-MAC,4}$ as the MCV. The challeng	ge to be encrypted is '7C 82 01 06 82	00 81 82 01 00 00 01 FF FF
BC A7' from the example in Table	18.	
	lt of encrypting the response using a	
	04'), and $T_{R-MAC,5}$ is computed using	
*	01 04 82 82 01 00 29 69 44 3B E	2 F1 2E' from the example in
Table 18.		
1C 87 07 9A FF 87 82 01 11 01		The GENERAL
$\langle \text{bytes } 1 - 250 \text{ of ENC}_{C5} \rangle$		AUTHENTICATE command is
		used to send a challenge to the
		PIV Card. This command
		includes the first part of the
		challenge.
	90 00	PIV Card Application indicates
		that it received the first part of
00.07.07.04.22.1251		the command successfully.
0C 87 07 9A 23 SENCE - 07 01 00 0F 00		The remaining challenge data is
272 of ENC <sub>C5</sub> > 97 01 00 8E 08		sent, including the BER-TLV
$T_8(T_{C-MAC,5})$		encoded L <sub>e</sub> and the BER-TLV
	97.92.01.17.02.4	encoded C-MAC.
	87 82 01 17 02 <bytes 1="" 251="" of<="" td="" –=""><td>PIV Card Application sends first part of the result of signing the</td></bytes>	PIV Card Application sends first part of the result of signing the
	$ENC_{R5} > 61 1B$	challenge. The padding indicator
		is '02' to indicate that no padding
		was required.
0C C0 00 00 1B		The remaining portion of
00 00 00 10		response is requested.
	   	PIV Card Application sends final
	02 90 00 8E 08 T <sub>8</sub> (T <sub>R-MAC.5</sub> ) 90	portion of the result of signing
	00	the challenge, along with the
		BER-TLV encoded status words
		and R-MAC.
	<u> </u>	

1108	Appendix B—Tern	ns, Acronyms, and Notation
1109	B.1 Terms	
1110	Application Identifier	A globally unique identifier of a card application as defined in ISO/IEC 7816-4.
1111 1112 1113	Algorithm Identifier	A PIV algorithm identifier is a one-byte identifier that specifies a cryptographic algorithm and key size. For symmetric cryptographic operations, the algorithm identifier also specifies a mode of operation (i.e., ECB).
1114 1115	Authenticable Entity	An entity that can successfully participate in an authentication protocol with a card application.
1116	BER-TLV Data Object	A data object coded according to ISO/IEC 8825-2.
1117	Card	An integrated circuit card.
1118 1119	Card Application	A set of data objects and card commands that can be selected using an application identifier.
1120 1121	Card Verifiable Certificate	A certificate stored on the card that includes a public key, the signature of a certification authority, and further information needed to verify the certificate.
1122 1123	Data Object	An item of information seen at the card command interface for which is specified a name, a description of logical content, a format, and a coding.
1124 1125 1126	Key Reference	A PIV key reference is a one-byte identifier that specifies a cryptographic key according to its PIV Key Type. The identifier is part of cryptographic material used in a cryptographic protocol such as an authentication or a signing protocol.
1127 1128	MAC Chaining Value	MAC Chaining Value is a 16-byte value that is input to the CMAC function. It is used to detect communication errors in duplicate or missing commands.
1129	Object Identifier	A globally unique identifier of a data object as defined in ISO/IEC 8824-2.
1130 1131 1132 1133	Reference Data	Cryptographic material used in the performance of a cryptographic protocol such as an authentication or a signing protocol. The reference data length is the maximum length of a password or PIN. For algorithms, the reference data length is the length of a key.
1134 1135	Status Word	Two bytes returned by an integrated circuit card after processing any command that signify the success of or errors encountered during said processing.
1136 1137	Template	A (constructed) BER-TLV data object whose value field contains specific BER-TLV data objects.

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1139	B.2	Acronyms
1140	AES	Advanced Encryption Standard
1141	AID	Application Identifier
1142	APDU	Application Protocol Data Unit
1143	API	Application Programming Interface
1144	APT	Application Property Template
1145	ASCII	American Standard Code for Information Interchange
1146	ASN.1	Abstract Syntax Notation One
1147	BER	Basic Encoding Rules
1148	CLA	Class (first) byte of a card command
1149	CMAC	Cipher-based Message Authentication Code
1150	C-MAC	Command Message Authentication Code
1151	CVC	Card Verifiable Certificate
1152	DER	Distinguished Encoding Rules
1153	DES	Data Encryption Standard
1154	ECB	Electronic Codebook
1155	ECC	Elliptic Curve Cryptography
1156	<b>ECDSA</b>	Elliptic Curve Digital Signature Algorithm
1157	ECDH	Elliptic Curve Diffie-Hellman
1158	EC CDH	Elliptic Curve Cryptography Cofactor Diffie-Hellman
1159		
1160	FIPS	Federal Information Processing Standards
1161	FISMA	Federal Information Security Management Act
1162	HSPD	Homeland Security Presidential Directive
1163	ICC	Integrated Circuit Card
1164	IEC	International Electrotechnical Commission
1165	IETF	Internet Engineering Task Force
1166	INS	Instruction (second) byte of a card command
1167	INCITS	InterNational Committee for Information Technology Standards
1168	ISO	International Organization for Standardization
1169	ITL	Information Technology Laboratory
1170	KDF	Key Derivation Function
1171	LSB	Least Significant Bit
1172	MAC	Message Authentication Code
1173	MSB	Most Significant Bit
1174	MCV	MAC Chaining Value
1175	NIST	National Institute of Standards and Technology
1176	OCC	On-Card Biometric Comparison
1177	OID	Object Identifier

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1178	OMB	Office of Management and Budget
1179	P1	First parameter of a card command
1180	P2	Second parameter of a card command
1181	PKCS	Public-Key Cryptography Standards
1182	PIN	Personal Identification Number
1183	PIV	Personal Identity Verification
1184	PIX	Proprietary Identifier extension
1185	PUK	PIN Unblocking Key
1186	RFU	Reserved for Future Use
1187	RID	Registered application provider Identifier
1188	R-MAC	Response Message Authentication Code
1189	RSA	Rivest, Shamir, Adleman
1190	SM	Secure Messaging
1191	S/MIME	Secure/Multipurpose Internet Mail Extensions
1192	SP	Special Publication
1193	SW1	First byte of a two-byte status word
1194	SW2	Second byte of a two-byte status word
1195	TLS	Transport Layer Security
1196	TLV	Tag-Length-Value
1197	VCI	Virtual Contact Interface
1198	B.3	Notation
1199	The sixteen h	<b>Notation</b> nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C,
1199 1200	The sixteen h	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits
1199 1200 1201	The sixteen h D, E, and F. are represente	Notation  exadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation
1199 1200 1201 1202	The sixteen h D, E, and F. are represent marks, for ex	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits
1199 1200 1201	The sixteen h D, E, and F. are represente	Notation  exadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation
1199 1200 1201 1202	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '01' diso be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the
1199 1200 1201 1202 1203	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least signification	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the last bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB.
1199 1200 1201 1202 1203	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least signification	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '01' diso be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the
1199 1200 1201 1202 1203 1204 1205	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least signification Thus, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and F. are represente marks, for example of the sixteen h D, E, and E, an	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the last bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB.
1199 1200 1201 1202 1203 1204 1205 1206	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least significa Thus, for exa All bytes spe	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '01' diso be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the ant bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB. Imple, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0.
1199 1200 1201 1202 1203 1204 1205 1206	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least signification. Thus, for exa All bytes spe	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the ant bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB. Imple, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0.  cified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.
1199 1200 1201 1202 1203 1204 1205 1206 1207 1208 1209	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least significa Thus, for exa All bytes spe All lengths sl The expression	Notation  Decadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits and in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation tample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the last bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB. Imple, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0.  Decified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  The last characters of two hexadecimal digits of the two hexadecimal digits of two hexadecimal digits of the two hexadecimal digits of the two hexadecimal digits of two hexadecimal digits of two hexadecimal digits of the two hexadecimal d
1199 1200 1201 1202 1203 1204 1205 1206 1207	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least significa Thus, for exa All bytes spe All lengths sl The expression The symbol	nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the eart bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB. Imple, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0. cified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.
1199 1200 1201 1202 1203 1204 1205 1206 1207 1208 1209 1210 1211	The sixteen h D, E, and F. are represente marks, for ex '16'.  A byte can al least significa Thus, for exa All bytes spe All lengths sl The expression The symbol   then X    Y is	Notation  Decadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits and in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the last bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB. ample, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0.  Decified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  The last and the least significant bit, b1, is 0.  Decified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  The last and the least significant bit, b1, is 0.  Decified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  The last and the least significant bit, b1, is 0.  Decified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  The last and the least significant bit, b1, is 0.  Decified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.
1199 1200 1201 1202 1203 1204 1205 1206 1207 1208 1209 1210 1211 1212	The sixteen h D, E, and F. are represented marks, for eximal least signification. Thus, for example All bytes specified and least significations. The expression of the expression of the transfer of the transfer of the transfer of the transfer of the symbol   then X    Y is the transfer of the transfer	Notation  exadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '00' '01' also be represented by bits b8 to b1, where b8 is the most significant bit (MSB) and b1 is the ant bit (LSB) of the byte. In textual or graphic representations, the leftmost bit is the MSB. ample, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0.  cified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  hall be measured in number of bytes unless otherwise noted.  on X & Y is a bitwise AND operation between bytes X and Y.  means concatenation of byte strings. For example, if X is '00 01 02' and Y is '03 04 05', '00 01 02 03 04 05'.  in templates are described as being mandatory (M), optional (O), or conditional (C).
1199 1200 1201 1202 1203 1204 1205 1206 1207 1208 1209 1210 1211 1212 1213	The sixteen h D, E, and F. are represented marks, for eximal least significations, for example and the symbol of them X    Y is the symbol of the	Notation  Decention  D
1199 1200 1201 1202 1203 1204 1205 1206 1207 1208 1209 1210 1211 1212	The sixteen h D, E, and F. are represented marks, for eximal least significations, for example and the symbol of them X    Y is the symbol of the	Notation  nexadecimal digits shall be denoted using the alphanumeric characters 0, 1, 2,, 9, A, B, C, A byte consists of two hexadecimal digits, for example, '2D'. The two hexadecimal digits ed in quotations '2D' or as 0x2D. A sequence of bytes may be enclosed in single quotation cample 'A0 00 00 01 16', rather than given as a sequence of individual bytes, 'A0' '00' '00' '01' on the byte. In textual or graphic representations, the leftmost bit is the MSB. ample, the most significant bit, b8, of '80' is 1 and the least significant bit, b1, is 0.  cified as RFU shall be set to '00' and all bits specified as RFU use shall be set to 0.  hall be measured in number of bytes unless otherwise noted.  on X & Y is a bitwise AND operation between bytes X and Y.  means concatenation of byte strings. For example, if X is '00 01 02' and Y is '03 04 05', '00 01 02 03 04 05'.  in templates are described as being mandatory (M), optional (O), or conditional (C).  means the data object shall appear in the template. 'Optional' means the data object may template. In the case of 'Conditional' data objects, the conditions under which they are

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1216 1217	In other tables the M/O/C column identifies properties of the PIV Card Application that shall be present (M), may be present (O), or are conditionally required to be present (C).
1218	BER-TLV data object tags are represented as byte sequences as described above. Thus, for example,
1219	0x4F is the interindustry data object tag for an application identifier and 0x7F60 is the interindustry data
1220	object tag for the biometric information template.

1245 1246

1221	
1222	Appendix C—References
1223	[ANSI504-1] Generic Identity Command Set – Part 1: Card Application Command Set.
1224 1225	[FIPS201] Federal Information Processing Standard 201-2, <i>Personal Identity Verification (PIV) of Federal Employees and Contractors</i> . (See <a href="http://csrc.nist.gov">http://csrc.nist.gov</a> )
1226 1227	[ISO7816] ISO/IEC 7816 (Parts 4, 5, 6, 8, and 9), Information technology — Identification cards — Integrated circuit(s) cards with contacts.
1228 1229	[ISO8824] ISO/IEC 8824-2:2002, Information technology Abstract Syntax Notation One (ASN.1): Information object specification.
1230 1231 1232	[ISO8825] ISO/IEC 8825-1:2002, Information technology — ASN.1 encoding rules: Specification of Basic Encoding Rules (BER), Canonical Encoding Rules (CER) and Distinguished Encoding Rules (DER).
1233 1234 1235	[PKCS1] Jakob Jonsson and Burt Kaliski, "Public-Key Cryptography Standards (PKCS) #1: RSA Cryptography Specifications Version 2.1", RFC 3447, February 2003. (See <a href="http://tools.ietf.org/html/rfc3447">http://tools.ietf.org/html/rfc3447</a> )
1236 1237	[SECG] Standards for Efficient Cryptography Group (SECG), "SEC 1: Elliptic Curve Cryptography", Version 1.0, September 2000.
1238 1239	[SP800-38B] NIST Special Publication 800-38B, <i>Recommendation for Block Cipher Modes of Operation: The CMAC Mode for Authentication</i> , May 2005. (See <a href="http://csrc.nist.gov">http://csrc.nist.gov</a> )
1240 1241	[SP800-56A] NIST Special Publication 800-56A, Recommendation for Pair-Wise Key Establishment Schemes Using Discrete Logarithm Cryptography (Revised), March 2007. (See <a href="http://csrc.nist.gov">http://csrc.nist.gov</a> )
1242 1243	[SP800-76] Draft NIST Special Publication 800-76-2, <i>Biometric Data Specification for Personal Identity Verification</i> , July 2012. (See <a href="http://csrc.nist.gov">http://csrc.nist.gov</a> )

[SP800-78] Draft NIST Special Publication 800-78-4, Cryptographic *Algorithms and Key Sizes for Personal Identity Verification*. (See <a href="http://csrc.nist.gov">http://csrc.nist.gov</a>)